

Solutions to In-Class Problems — Week 2, Mon

Problem 1. For each of the logical formulas, indicate whether or not it is true when the domain of discourse is \mathbb{N} (the natural numbers $0, 1, 2, \dots$), \mathbb{Z} (the integers), \mathbb{Q} (the rationals), \mathbb{R} (the real numbers), and \mathbb{C} (the complex numbers).

$$\begin{array}{l} \exists x \quad (x^2 = 2) \\ \forall x \exists y \quad (x^2 = y) \\ \forall y \exists x \quad (x^2 = y) \\ \forall x \neq 0 \exists y \quad (xy = 1) \\ \exists x \exists y \quad (x + 2y = 2) \wedge (2x + 4y = 5) \end{array}$$

Solution.

<i>Statement</i>	\mathbb{N}	\mathbb{Z}	\mathbb{Q}	\mathbb{R}	\mathbb{C}
$\exists x (x^2 = 2)$	<i>f</i>	<i>f</i>	<i>f</i>	<i>t</i> ($x = \sqrt{2}$)	<i>t</i>
$\forall x \exists y (x^2 = y)$	<i>t</i>	<i>t</i>	<i>t</i>	<i>t</i> ($y = x^2$)	<i>t</i>
$\forall y \exists x (x^2 = y)$	<i>f</i>	<i>f</i>	<i>f</i>	<i>f</i> (take $y < 0$)	<i>t</i>
$\forall x \neq 0 \exists y (xy = 1)$	<i>f</i>	<i>f</i>	<i>t</i>	<i>t</i> ($y = 1/x$)	<i>t</i>
$\exists x \exists y (x + 2y = 2) \wedge (2x + 4y = 5)$	<i>f</i>	<i>f</i>	<i>f</i>	<i>f</i>	<i>f</i>



Problem 2. A media tycoon has an idea for an all-news television network called LNN: The Logic News Network. Each segment will begin with a definition of the domain of discourse and a few predicates. The day's happenings can then be communicated concisely in logic notation. For example, a broadcast might begin as follows:

“THIS IS LNN. The domain of discourse is {Bill, Monica, Ken, Linda, Betty}. Let $D(x)$ be a predicate that is true if x is deceitful. Let $L(x, y)$ be a predicate that is true if x likes y . Let $G(x, y)$ be a predicate that is true if x gave gifts to y .”

Complete the broadcast by translating the following statements into logic notation.

(a) If neither Monica nor Linda is deceitful, then Bill and Monica like each other.

Solution.

$$(\neg(D(\text{Monica}) \vee D(\text{Linda}))) \longrightarrow (L(\text{Bill}, \text{Monica}) \wedge L(\text{Monica}, \text{Bill}))$$

■

(b) Everyone except for Ken likes Betty, and no one except Linda likes Ken.

Solution.

$$\begin{aligned} \forall x (x = \text{Ken} \wedge \neg L(x, \text{Betty})) \vee (x \neq \text{Ken} \wedge L(x, \text{Betty})) \wedge \\ \forall x (x = \text{Linda} \wedge L(x, \text{Ken})) \vee (x \neq \text{Linda} \wedge \neg L(x, \text{Ken})) \end{aligned}$$

■

(c) If Ken is not deceitful, then Bill gave gifts to Monica, and Monica gave gifts to someone.

Solution.

$$\neg D(\text{Ken}) \longrightarrow (G(\text{Bill}, \text{Monica}) \wedge \exists x G(\text{Monica}, x))$$

■

(d) Everyone likes someone and dislikes someone else.

Solution.

$$\forall x \exists y \exists z (y \neq z) \wedge L(x, y) \wedge \neg L(x, z)$$

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(e) How could you express “Everyone except for Ken likes Betty” using just propositional connectives *without* using any quantifiers (\forall, \exists)? Can you generalize to explain how *any* logical formula over this domain of discourse can be expressed without quantifiers? How big would the formula in the previous part be if it was expressed this way?

Solution.

$$L(\text{Bill}, \text{Betty}) \wedge L(\text{Monica}, \text{Betty}) \wedge L(\text{Linda}, \text{Betty}) \wedge L(\text{Betty}, \text{Betty})$$

In general, quantifiers can be eliminated by treating $\forall x P(x)$ as an abbreviation for

$$P(\text{Bill}) \wedge P(\text{Monica}) \wedge P(\text{Ken}) \wedge P(\text{Linda}) \wedge P(\text{Betty}),$$

and $\exists x P(x)$ as an abbreviation for

$$P(\text{Bill}) \vee P(\text{Monica}) \vee P(\text{Ken}) \vee P(\text{Linda}) \vee P(\text{Betty}).$$

Expanded this way, the three-quantifier formula of the previous part would expand by a factor of $5 \times 5 \times 5 = 125$. So using quantifiers can pay off even when they are not strictly necessary. ■

Problem 3. Prove by induction that

$$\sum_{i=0}^n (2i+1)^2 = \frac{(n+1)(2n+1)(2n+3)}{3}$$

Solution. See Rosen, solution to §3.2, Exercise 9. ■