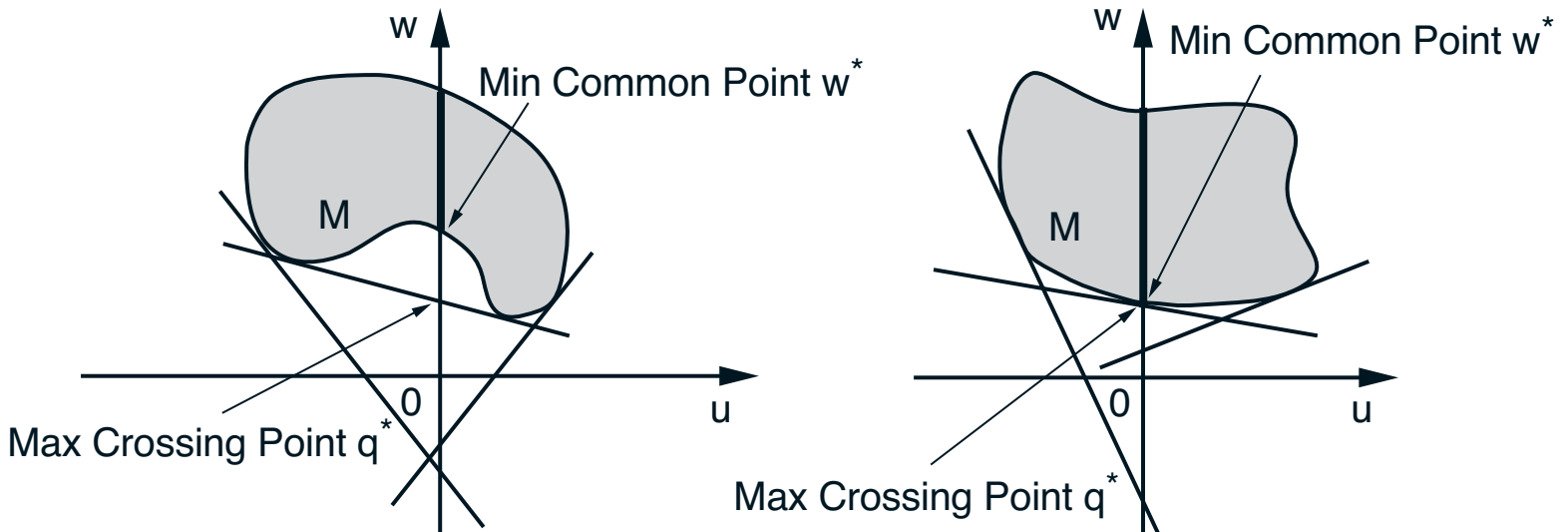


# LECTURE 8

## LECTURE OUTLINE

- Min Common / Max Crossing problems
- Weak duality
- Strong duality
- Existence of optimal solutions
- Minimax problems



# WEAK DUALITY

- Optimal value of the min common problem:

$$w^* = \inf_{(0,w) \in M} w.$$

- Max crossing problem: Focus on hyperplanes with normals  $(\mu, 1)$  whose crossing point  $\xi$  satisfies

$$\xi \leq w + \mu'u, \quad \forall (u, w) \in M.$$

Max crossing problem is to maximize  $\xi$  subject to  $\xi \leq \inf_{(u,w) \in M} \{w + \mu'u\}$ ,  $\mu \in \mathbb{R}^n$ , or

$$\text{maximize } q(\mu) \triangleq \inf_{(u,w) \in M} \{w + \mu'u\}$$

subject to  $\mu \in \mathbb{R}^n$ .

- For all  $(u, w) \in M$  and  $\mu \in \mathbb{R}^n$ ,

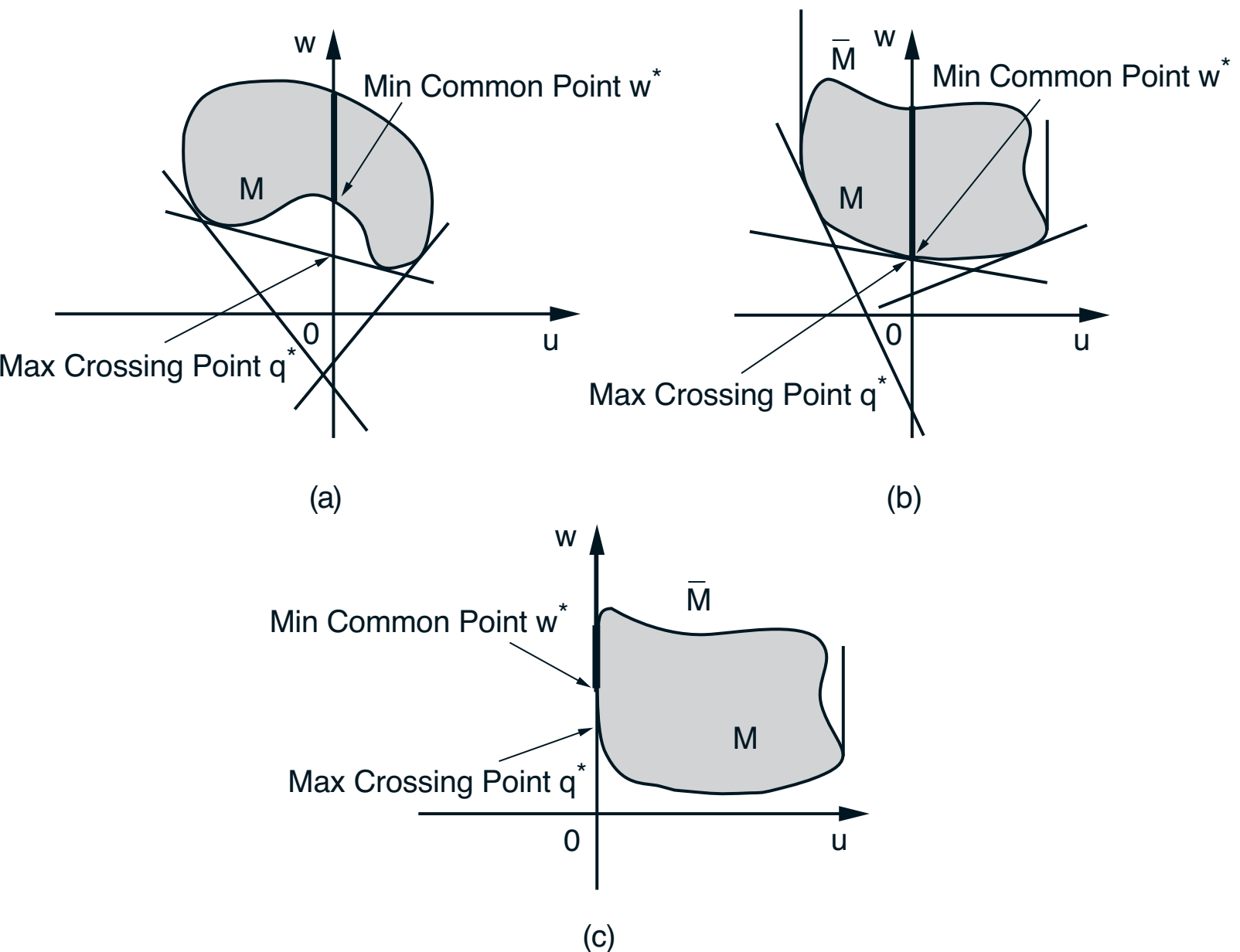
$$q(\mu) = \inf_{(u,w) \in M} \{w + \mu'u\} \leq \inf_{(0,w) \in M} w = w^*,$$

so by taking the supremum of the left-hand side over  $\mu \in \mathbb{R}^n$ , we obtain  $q^* \leq w^*$ .

- Note that  $q$  is concave and upper-semicontinuous.

# STRONG DUALITY

- Question: Under what conditions do we have  $q^* = w^*$  and the supremum in the max crossing problem is attained?



## DUALITY THEOREMS

- Assume that  $w^* < \infty$  and that the set

$$\overline{M} = \left\{ (u, w) \mid \text{there exists } \overline{w} \text{ with } \overline{w} \leq w \text{ and } (u, \overline{w}) \in M \right\}$$

is convex.

- *Min Common/Max Crossing Theorem I:* We have  $q^* = w^*$  if and only if for every sequence  $\{(u_k, w_k)\} \subset M$  with  $u_k \rightarrow 0$ , there holds  $w^* \leq \liminf_{k \rightarrow \infty} w_k$ .

- *Min Common/Max Crossing Theorem II:* Assume in addition that  $-\infty < w^*$  and that the set

$$D = \left\{ u \mid \text{there exists } w \in \mathfrak{R} \text{ with } (u, w) \in \overline{M} \right\}$$

contains the origin in its relative interior. Then  $q^* = w^*$  and there exists a vector  $\mu \in \mathfrak{R}^n$  such that  $q(\mu) = q^*$ . If  $D$  contains the origin in its interior, the set of all  $\mu \in \mathfrak{R}^n$  such that  $q(\mu) = q^*$  is compact.

- *Min Common/Max Crossing Theorem III:* Involves polyhedral assumptions, and will be developed later.

## PROOF OF THEOREM I

• Assume that for every sequence  $\{(u_k, w_k)\} \subset M$  with  $u_k \rightarrow 0$ , there holds  $w^* \leq \liminf_{k \rightarrow \infty} w_k$ . If  $w^* = -\infty$ , then  $q^* = -\infty$ , by weak duality, so assume that  $-\infty < w^*$ . Steps of the proof:

- (1)  $\overline{M}$  does not contain any vertical lines.
- (2)  $(0, w^* - \epsilon) \notin \text{cl}(\overline{M})$  for any  $\epsilon > 0$ .
- (3) There exists a nonvertical hyperplane strictly separating  $(0, w^* - \epsilon)$  and  $\overline{M}$ . This hyperplane crosses the  $(n + 1)$ st axis at a vector  $(0, \xi)$  with  $w^* - \epsilon \leq \xi \leq w^*$ , so  $w^* - \epsilon \leq q^* \leq w^*$ . Since  $\epsilon$  can be arbitrarily small, it follows that  $q^* = w^*$ .

Conversely, assume that  $q^* = w^*$ . Let  $\{(u_k, w_k)\} \subset M$  be such that  $u_k \rightarrow 0$ . Then,

$$q(\mu) = \inf_{(u, w) \in M} \{w + \mu' u\} \leq w_k + \mu' u_k, \quad \forall k, \forall \mu \in \mathfrak{R}^n.$$

Taking the limit as  $k \rightarrow \infty$ , we obtain  $q(\mu) \leq \liminf_{k \rightarrow \infty} w_k$ , for all  $\mu \in \mathfrak{R}^n$ , implying that

$$w^* = q^* = \sup_{\mu \in \mathfrak{R}^n} q(\mu) \leq \liminf_{k \rightarrow \infty} w_k.$$

## PROOF OF THEOREM II

• Note that  $(0, w^*)$  is not a relative interior point of  $\overline{M}$ . Therefore, by the Proper Separation Theorem, there exists a hyperplane that passes through  $(0, w^*)$ , contains  $\overline{M}$  in one of its closed halfspaces, but does not fully contain  $\overline{M}$ , i.e., there exists  $(\mu, \beta)$  such that

$$\beta w^* \leq \mu'u + \beta w, \quad \forall (u, w) \in \overline{M},$$

$$\beta w^* < \sup_{(u, w) \in \overline{M}} \{\mu'u + \beta w\}.$$

Since for any  $(\bar{u}, \bar{w}) \in M$ , the set  $\overline{M}$  contains the halfline  $\{(\bar{u}, w) \mid \bar{w} \leq w\}$ , it follows that  $\beta \geq 0$ . If  $\beta = 0$ , then  $0 \leq \mu'u$  for all  $u \in D$ . Since  $0 \in \text{ri}(D)$  by assumption, we must have  $\mu'u = 0$  for all  $u \in D$  a contradiction. Therefore,  $\beta > 0$ , and we can assume that  $\beta = 1$ . It follows that

$$w^* \leq \inf_{(u, w) \in \overline{M}} \{\mu'u + w\} = q(\mu) \leq q^*.$$

Since the inequality  $q^* \leq w^*$  holds always, we must have  $q(\mu) = q^* = w^*$ .

## MINIMAX PROBLEMS

Given  $\phi : X \times Z \mapsto \mathbb{R}$ , where  $X \subset \mathbb{R}^n$ ,  $Z \subset \mathbb{R}^m$   
consider

$$\text{minimize } \sup_{z \in Z} \phi(x, z)$$

$$\text{subject to } x \in X$$

and

$$\text{maximize } \inf_{x \in X} \phi(x, z)$$

$$\text{subject to } z \in Z.$$

- Some important contexts:
  - Worst-case design. Special case: Minimize over  $x \in X$

$$\max \{ f_1(x), \dots, f_m(x) \}$$

- Duality theory and zero sum game theory (see next two slides)
- We will study minimax problems using the min common/max crossing framework

# MAIN DUALITY STRUCTURE

- For the problem

minimize  $f(x)$

subject to  $x \in X$ ,  $g_j(x) \leq 0$ ,  $j = 1, \dots, r$

introduce the Lagrangian function

$$L(x, \mu) = f(x) + \sum_{j=1}^r \mu_j g_j(x)$$

- Primal problem (equivalent to the original)

$$\min_{x \in X} \sup_{\mu \geq 0} L(x, \mu) = \begin{cases} f(x) & \text{if } g(x) \leq 0, \\ \infty & \text{otherwise,} \end{cases}$$

- Dual problem

$$\max_{\mu \geq 0} \inf_{x \in X} L(x, \mu)$$

- Key duality question: Is it true that

$$\sup_{\mu \geq 0} \inf_{x \in \mathfrak{R}^n} L(x, \mu) = \inf_{x \in \mathfrak{R}^n} \sup_{\mu \geq 0} L(x, \mu)$$

# ZERO SUM GAMES

- Two players: 1st chooses  $i \in \{1, \dots, n\}$ , 2nd chooses  $j \in \{1, \dots, m\}$ .
- If moves  $i$  and  $j$  are selected, the 1st player gives  $a_{ij}$  to the 2nd.
- Mixed strategies are allowed: The two players select probability distributions

$$x = (x_1, \dots, x_n), \quad z = (z_1, \dots, z_m)$$

over their possible moves.

- Probability of  $(i, j)$  is  $x_i z_j$ , so the expected amount to be paid by the 1st player

$$x'Az = \sum_{i,j} a_{ij} x_i z_j$$

where  $A$  is the  $n \times m$  matrix with elements  $a_{ij}$ .

- Each player optimizes his choice against the worst possible selection by the other player. So
  - 1st player minimizes  $\max_z x'Az$
  - 2nd player maximizes  $\min_x x'Az$

# MINIMAX INEQUALITY

- We always have

$$\sup_{z \in Z} \inf_{x \in X} \phi(x, z) \leq \inf_{x \in X} \sup_{z \in Z} \phi(x, z)$$

[for every  $\bar{z} \in Z$ , write

$$\inf_{x \in X} \phi(x, \bar{z}) \leq \inf_{x \in X} \sup_{z \in Z} \phi(x, z)$$

and take the sup over  $\bar{z} \in Z$  of the left-hand side].

- This is called the *minimax inequality*. When it holds as an equation, it is called the *minimax equality*.
- The minimax equality need not hold in general.
- When the minimax equality holds, it often leads to interesting interpretations and algorithms.
- The minimax inequality is often the basis for interesting bounding procedures.