#### **Spatially Distributed Queues**

M/G/1 2 Servers N servers Approximations

### Why Spatial Queues?

Herein Construction Systems **#Organ donation queues** Warehouses **Supply chains #Cell phone systems #**People waiting to be evacuated in a hurricane





#### M/G/1



#### **M/G/1**

#Ambulance always returns home with each service; standard M/G/1 applies
#But suppose we have an emergency repair vehicle that travels directly from one customer to the next?.....

 $\overline{S_1}, \sigma_{S_1}^2 =$  expected value and variance, respectively, of the 1st service time in a busy period  $\overline{S_2}, \sigma_{S_2}^2 =$  expected value and variance, respectively, of the 2nd & all succeeding service times in a busy period  $\lambda \overline{S_2} < 1$ 

 $\rho = 1 - P_0$  = fraction of time server is busy

$$\rho = \frac{\lambda \overline{S_1}}{1 - \lambda (\overline{S_2} - \overline{S_1})}$$

$$L = \rho + \frac{\lambda^2}{1 - \lambda(\overline{S_2} - \overline{S_1})} \left[ \frac{\sigma_{S_1}^2 + \overline{S_1}^2 + \lambda \{\overline{S_1}(\sigma_{S_2}^2 + \overline{S_2}^2) - \overline{S_2}(\sigma_{S_1}^2 + \overline{S_1}^2)}{2(1 - \lambda \overline{S_2})} \right]$$

Little's Law: Buy one, get three others for free!

 $L = \lambda W$ 

$$L_q = \lambda W_q$$

See the book, Eqs. (5.0) - (5.5)

### Boes this new more general M/G/1 model apply exactly to the ambulance problem?

₩Why or why not?

#### Two-Server "Hypercube" Queueing Model

Distinguishable servers
 Different workloads (due to geography)
 Can appear with or without queueing
 With -- usually FCFS

Without -- usually means a backup contract service is in place

![](_page_10_Picture_0.jpeg)

![](_page_11_Picture_0.jpeg)

![](_page_12_Picture_0.jpeg)

![](_page_13_Picture_0.jpeg)

![](_page_14_Picture_0.jpeg)

![](_page_15_Picture_0.jpeg)

![](_page_16_Picture_0.jpeg)

![](_page_17_Picture_0.jpeg)

#### Poisson Arrivals from any sub-region A

![](_page_18_Figure_1.jpeg)

#### **Service Discipline**

#### %1st Dispatch Preference to 'primary server'

## Content of the server, if available

## Contension: See Section 2018 Section 2018

![](_page_20_Picture_0.jpeg)

![](_page_21_Figure_0.jpeg)

#### Balance of Flow Equations, Loss System

![](_page_22_Figure_1.jpeg)

![](_page_23_Figure_0.jpeg)

![](_page_24_Figure_0.jpeg)

#### Balance of Flow Equations, Loss System

$$P_{00} (\lambda_1 + \lambda_2) = P_{01} \mu + P_{10} \mu$$

$$P_{01} \left(\lambda + \mu\right) = P_{11} \mu + P_{00} \lambda_1$$

Etc.

$$P_{00} + P_{10} + P_{01} + P_{11} = 1$$

#### **Workload and Imbalances**

$$\begin{aligned} & \Re \rho_1 = W_1 = P_{01} + P_{11} \\ & \Re \rho_2 = W_2 = P_{10} + P_{11} \\ & \Re \text{ Workload Imbalance} = \Delta W = |W_1 - W_2| \end{aligned}$$

#### To Obtain Travel Times, We Must Have Server Response Patterns

 $\Re f_{nj}$  = fraction of dispatches that are server *n* to response area *j* 

 $\mathcal{H}_{n}(\mathbf{C}) = \text{average time for server } n \text{ to travel}$ to a customer in region  $\mathbf{C}$ 

ℋT(A) = average system-wide travel time, assuming that server 1's primary response area is region A.

#### Average System-Wide Travel Time

### $T(\mathbf{A}) = f_{11}T_1(\mathbf{A}) + f_{22}T_2(\mathbf{B}-\mathbf{A})$ $+ f_{12}T_1(\mathbf{B}-\mathbf{A}) + f_{21}T_2(\mathbf{A})$

#### Average System-Wide Travel Time

 $T(\mathbf{A}) = f_{11}T_1(\mathbf{A}) + f_{22}T_2(\mathbf{B}-\mathbf{A}) + f_{12}T_1(\mathbf{B}-\mathbf{A}) + f_{21}T_2(\mathbf{A})$ Queueing

#### Average System-Wide Travel Time

 $T(\mathbf{A}) = f_{11} T_1(\mathbf{A}) + f_{22} T_2(\mathbf{B}-\mathbf{A}) + f_{12} T_1(\mathbf{B}-\mathbf{A}) + f_{21} T_2(\mathbf{A})$ 

![](_page_30_Picture_2.jpeg)

#### How do we obtain the $f_{ni}$ 's?

Consider a long time interval T

 $f_{12}$ =(# requests that assign unit 1 to area 2)/ (total # requests answered)

**#**Total # requests answered =  $(1-P_{11})\lambda T$ 

**#**Average # requests that are "server 1 to area 2" is  $\lambda_2 TP_{10}$ . Why?

**Solution** Herefore 
$$f_{12} = (\lambda_2 T P_{10} / [1 - P_{11}]\lambda T) = {\lambda_2/(1 - P_{11})\lambda} P_{10}$$

# How do we generalize this to *N* servers?

### New York City EMS Hypercube

![](_page_33_Figure_1.jpeg)

#### New York City EMS Hypercube

![](_page_34_Figure_1.jpeg)

#### Rectangular City Example

![](_page_35_Figure_1.jpeg)

![](_page_36_Figure_0.jpeg)

### **Optimal Districting**

"Dispatch the closest available server' is often not optimal, where 'optimal' implies minimizing mean travel time

### May not be good for reducing workload imbalance either

₩With numerical example in book, the optimal boundary line is shifted to the right by 10/126 miles.

![](_page_38_Figure_0.jpeg)

#### **Boundary Line Comparison**

#Equal travel time boundary line  $ightarrow T(A_{w=1/2})=0.46246$   $ightarrow \Delta W = 0.05236$ #Optimal boundary line  $ightarrow T(A_{w^*})=0.46166$  $ightarrow \Delta W = 0.04405$ 

#### Two server Loss Model: Boundary Line Result

### **%To minimize mean city-wide mean travel time:**

\* The optimal partitioning consists of a set of points within the region that is a constant travel time s<sub>0</sub> closer to facility 1 than to facility 2. (Carter, Chaiken, Ignall, 1972)

#Does our rectangular city example work
for this?

#### So: Optimal Partitioning

 $\alpha \equiv \lambda/2\mu$ 

 $\mu_1 = \mu_2$ 

 $S_0 = [2\alpha/(2\alpha+1)]\{T_2(B) - T_1(B)\}$ 

![](_page_42_Figure_0.jpeg)

![](_page_43_Picture_0.jpeg)

![](_page_44_Figure_0.jpeg)

![](_page_45_Figure_0.jpeg)

![](_page_46_Picture_0.jpeg)

## And what about the corner case?

![](_page_48_Figure_0.jpeg)

#### The S<sub>0</sub> result is general

%Works for discrete grid
%One way streets
%General transportation network

∺Rick Jarvis, in an MIT Ph.D. thesis, generalized this to N servers