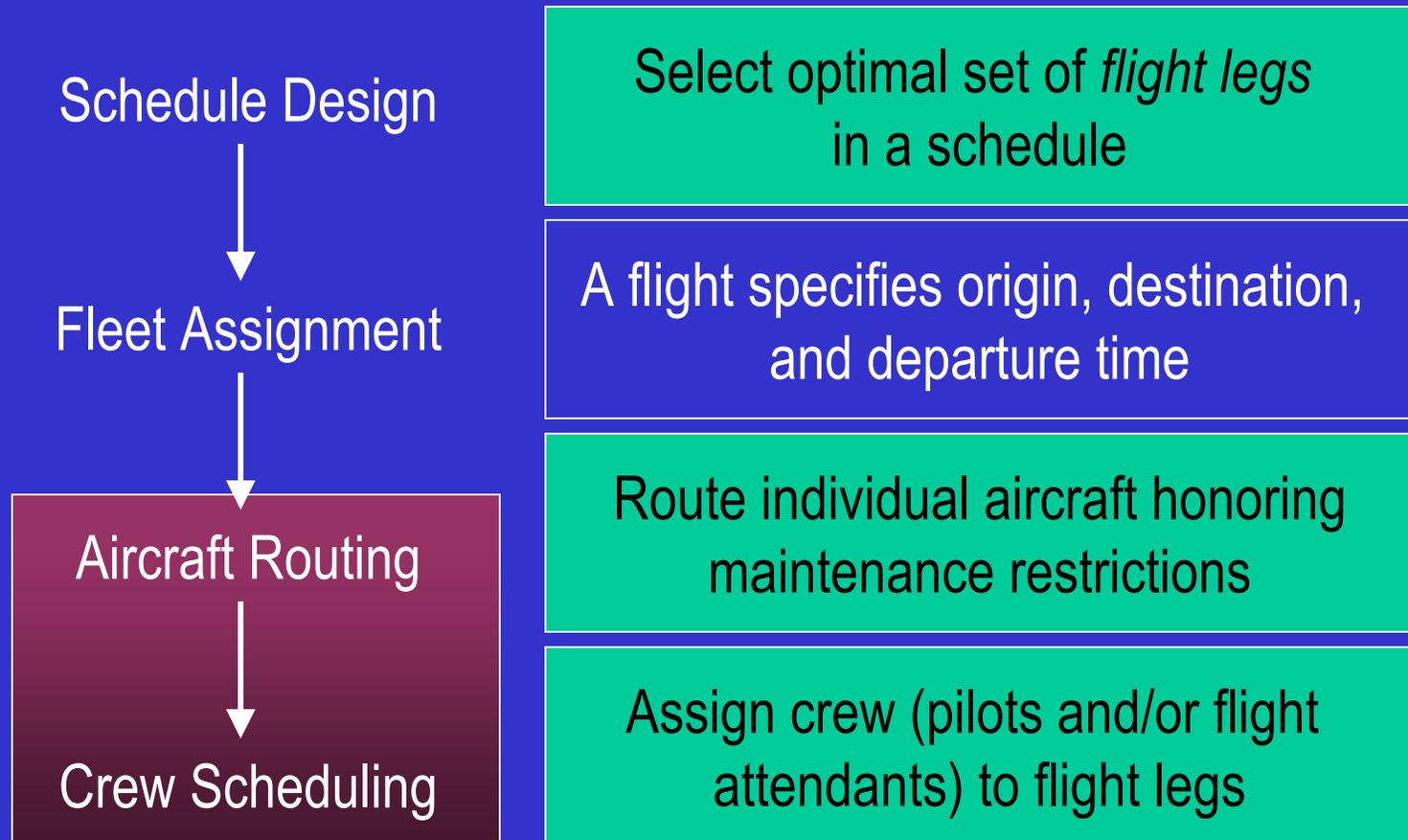


# 1.224J/ ESD.204J: Airline Crew Scheduling

## Outline

- Crew Scheduling
  - Problem Definition
    - Costs
  - Set Partitioning Model and Solution
- Enhanced Crew Scheduling
  - Link to Maintenance Routing Problem
  - Enhanced Model and Solution Approach

# Airline Schedule Planning



# Crew Scheduling: Some Background

- This problem has been studied by operations researchers for at least 4 decades
- Most major U.S. airlines use crew pairing optimizers for the cockpit crews
  - Crew costs are the airlines' second largest operating expense
  - Even small improvements in efficiency can have large financial benefits

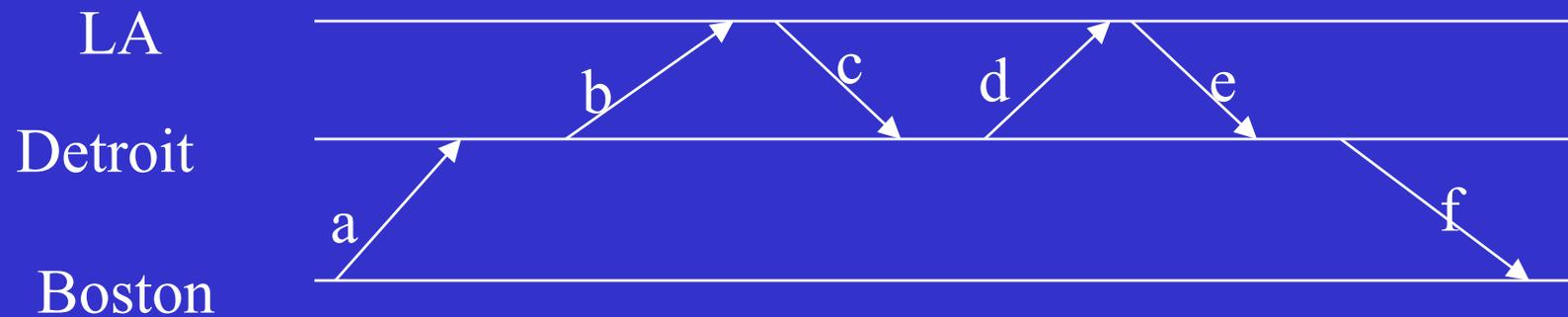
# Airline Crew Scheduling

- 2-stage process:
  - Crew Pairing Optimization
    - Construct minimum cost work schedules, called pairings, spanning several days
  - Bidline Generation/ Rostering
    - Construct monthly work schedules from the pairings generated in the first stage
      - Bidlines
      - Individualized schedules
    - Objective to balance workload, maximize number of crew requests granted, etc.

# Some Definitions

- A *crewbase* is the home station, or domicile, of a crew
- A *crew pairing* is a sequence of *flights* that can be flown by a single crew:
  - beginning and ending at a crewbase
  - spanning one or more days
  - satisfying FAA rules and collective bargaining agreements, such as:
    - maximum flying time in a day
    - minimum rest requirements
    - minimum connection time between two flights

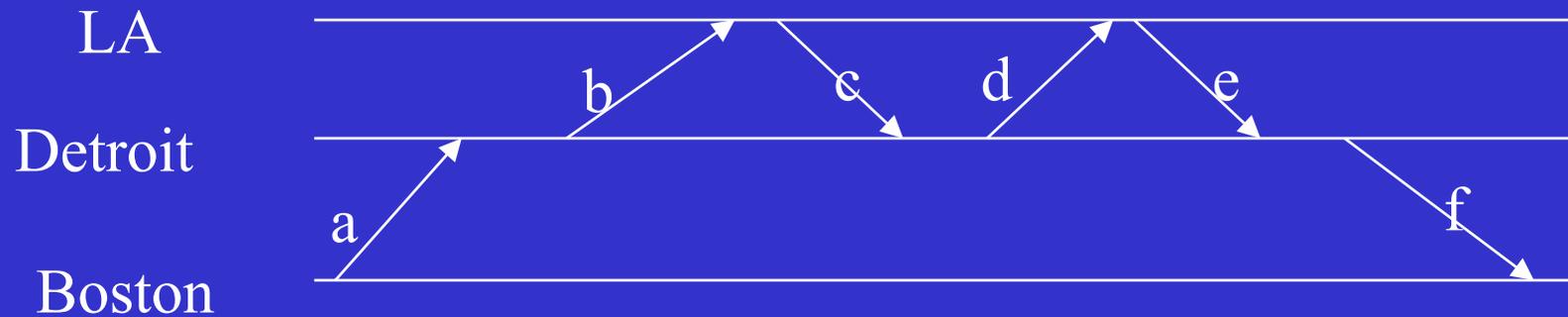
# Example: A Crew Pairing



# Some More Definitions

- A *duty period* (or *duty*) is a sequence of flight legs comprising a day of work for a crew
  - Alternative pairing definition: a *crew pairing* is a sequence of *duties* separated by rests
- A *crew schedule* is a sequence of *pairings* separated by time-off, satisfying numerous restrictions from regulatory agencies and collective bargaining agreements

# Example: Duty Periods



$$\text{Pairing} = \text{DP1}(a,b,c) + \text{rest} + \text{DP2}(d,e) + \text{rest} + \text{DP3}(f)$$



# Crew Pairing Problem (CP)

- Assign crews to flights such that every flight is covered, costs are minimized and labor rules are satisfied:
  - Maximum flying time in a day
  - Minimum rest requirements
  - Minimum connection time

# Crew Pairing Costs

- Duty costs: *Maximum* of 3 elements:
  - $f1$ \*flying time cost
  - $f2$ \*elapsed time cost
  - $f3$ \*minimum guarantee
- Pairing costs: *Maximum* of 3 elements:
  - $f1$ \*duty cost
  - $f2$ \*time-away-from-base
  - $f3$ \*minimum guarantee

# Set Partitioning Model for CP: Variable Definition and Constraints

- A variable is a *pairing*
  - Binary variables: =1 if pairing is assigned to a crew; = 0 if pairing not flown
- Set partitioning constraints require each flight to be covered exactly once
- Number of possible pairings (variables) grows exponentially with the number of flights

# An Example

Flights:

A B C D E F G H

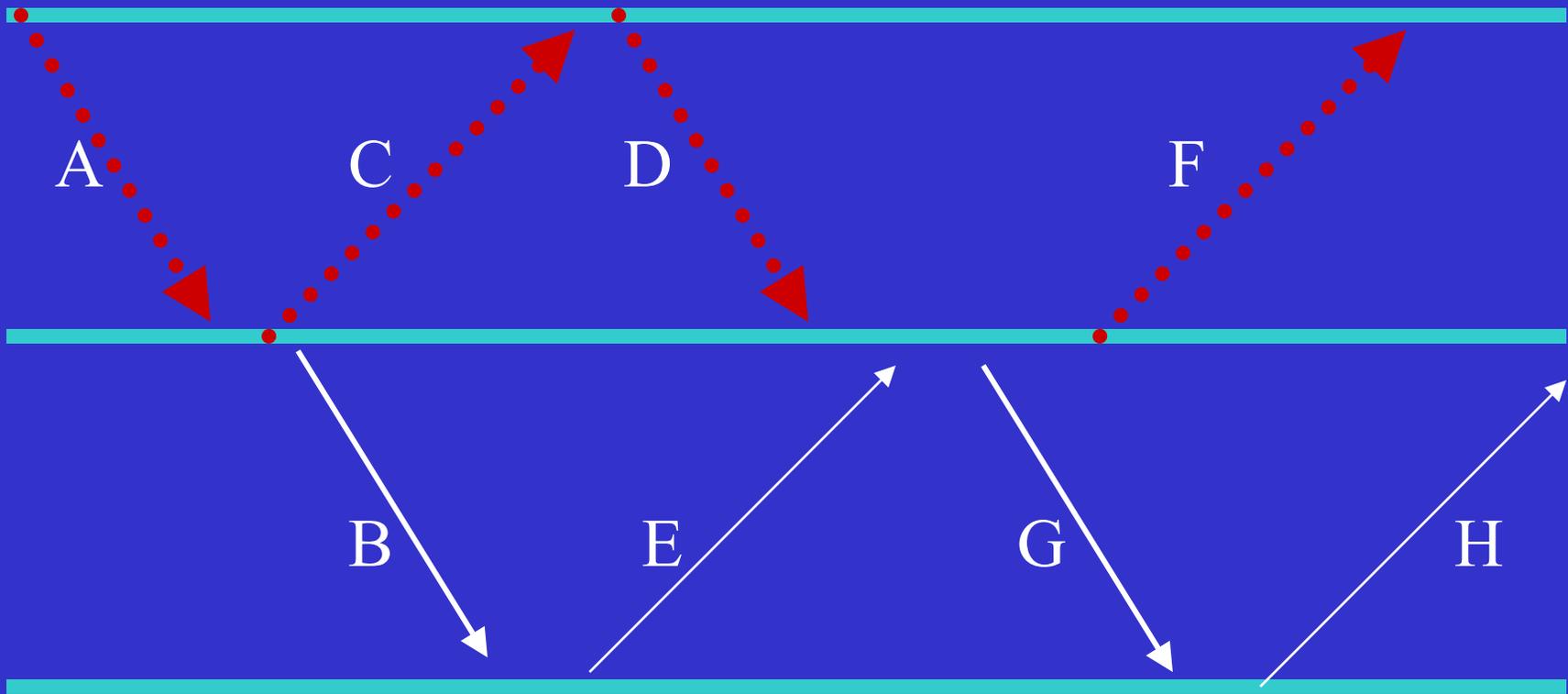
Potential pairings:

- A-C-D-F  $(y_1)$ : \$1
- A-B-E-F  $(y_2)$ : \$2
- C-D-G-H  $(y_3)$ : \$4
- B-E-G-H  $(y_4)$ : \$6

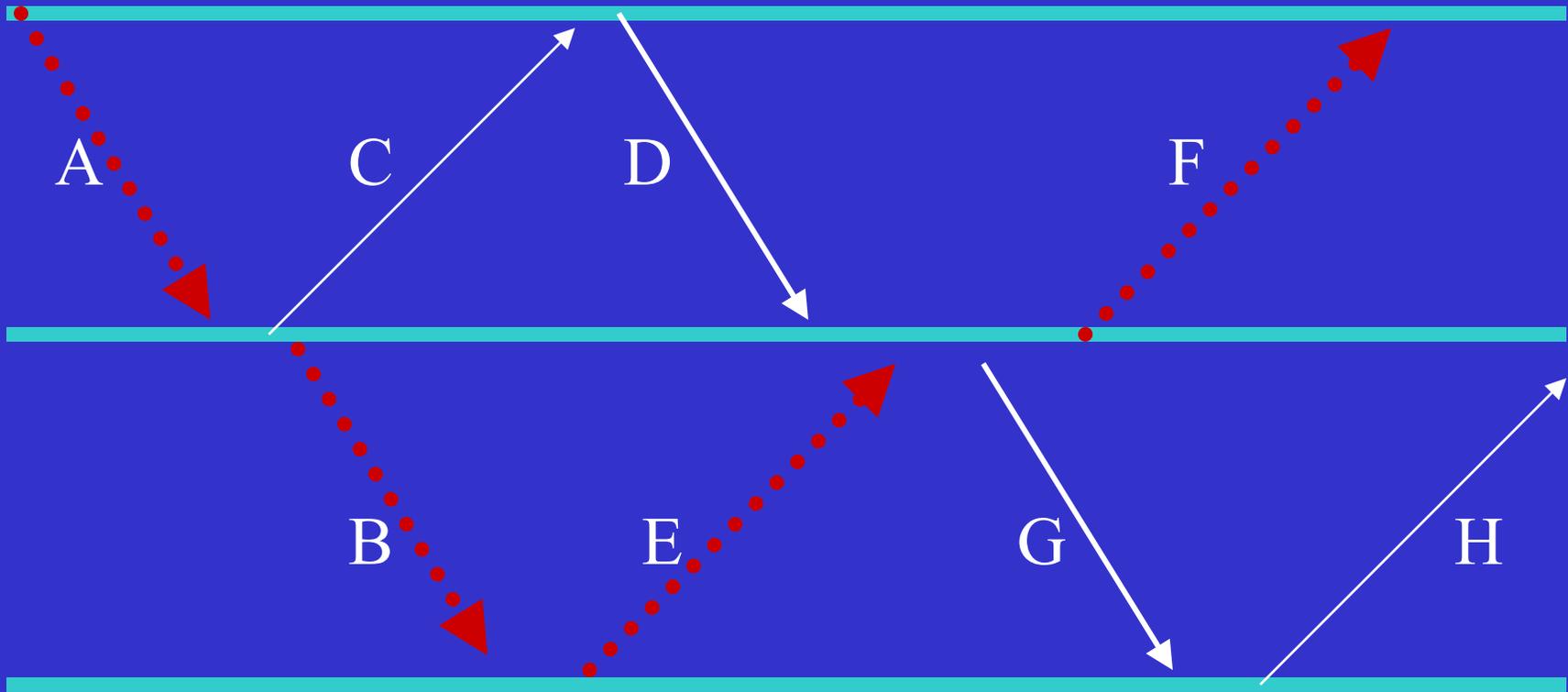
Crew pairing solutions:

- $x_1 \Rightarrow$  pairings 1, 4: \$7
- $x_2 \Rightarrow$  pairings 2, 3: \$6

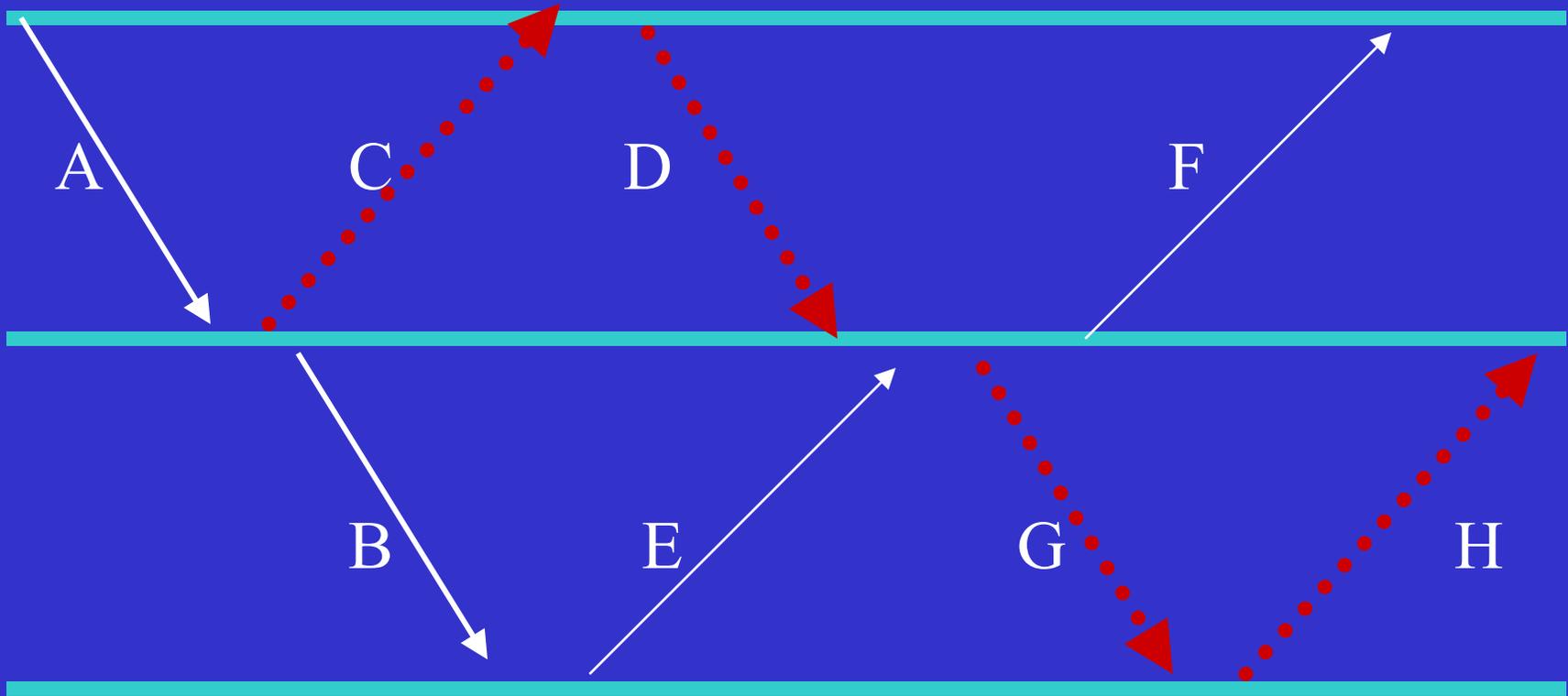
# Pairing 1



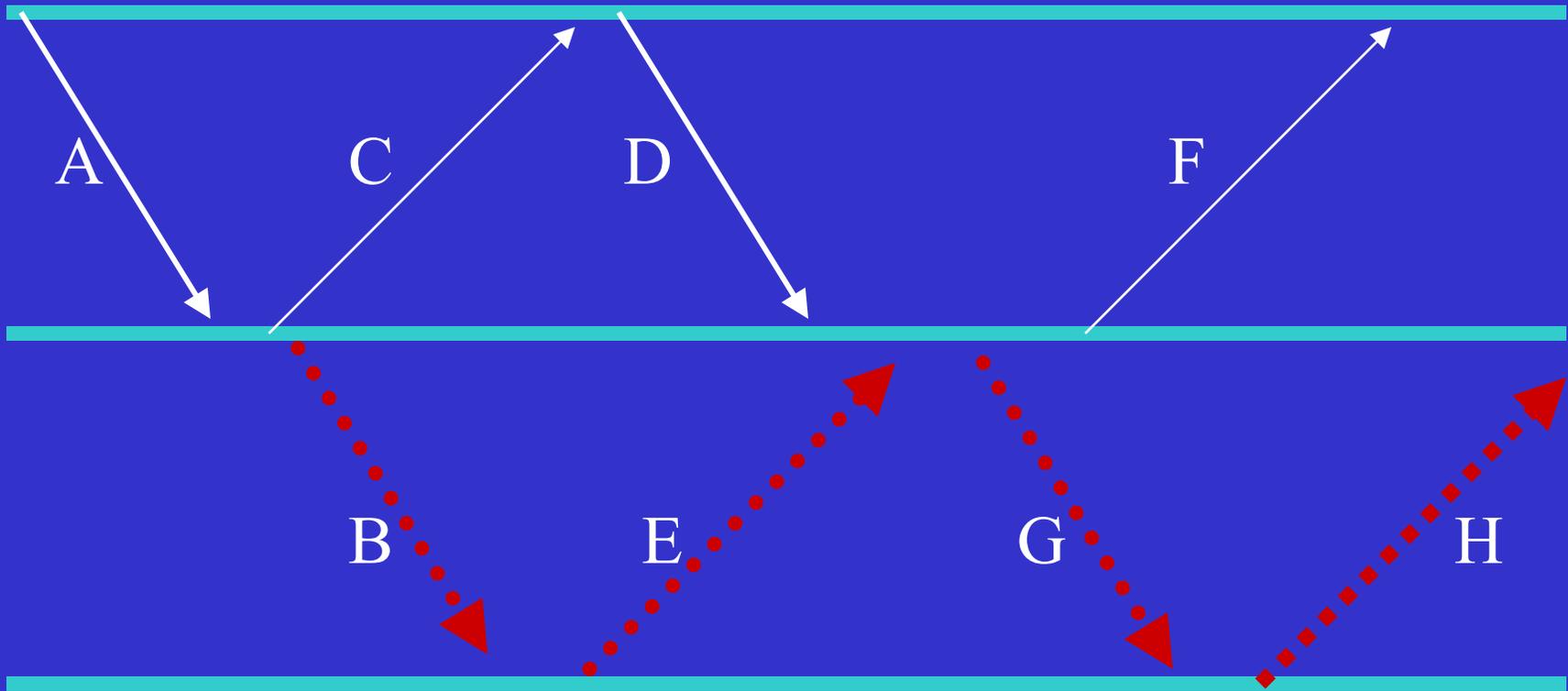
# Pairing 2



# Pairing 3



# Pairing 4



# Notation

- $P^k$  is the set of feasible pairings for fleet family  $k$
- $F^k$  is the set of daily flights assigned to fleet family  $k$
- $\delta_{fp}$  equals 1 if flight  $f$  is in pairing  $p$ , else 0
- $c_p$  is the cost of pairing  $p$
- $y_p$  is 1 if pairing  $p$  is in the solution, else 0

# Formulation

$$\min \sum_{p \in P^k} c_p y_p$$

*st*

$$\sum_{p \in P^k} \delta_{fp} y_p = 1 \quad \forall f \in F^k$$

$$y_p \in \{0,1\} \quad \forall p \in P^k$$

# Example

Flight  
Cover  
Constraints:

$y_1$	+	$y_2$		=	1	A		
		$y_2$	+		$y_4$	= 1	B	
$y_1$	+			$y_3$		= 1	C	
$y_1$	+			$y_3$		= 1	D	
		$y_2$	+		$y_4$	= 1	E	
$y_1$	+	$y_2$				= 1	F	
				$y_3$	+	$y_4$	= 1	G
				$y_3$	+	$y_4$	= 1	H
$y_1$						∈ 0,1	1	
		$y_2$				∈ 0,1	2	
				$y_3$		∈ 0,1	3	
					$y_4$	∈ 0,1	4	

Binary  
Pairing  
Restrictions:

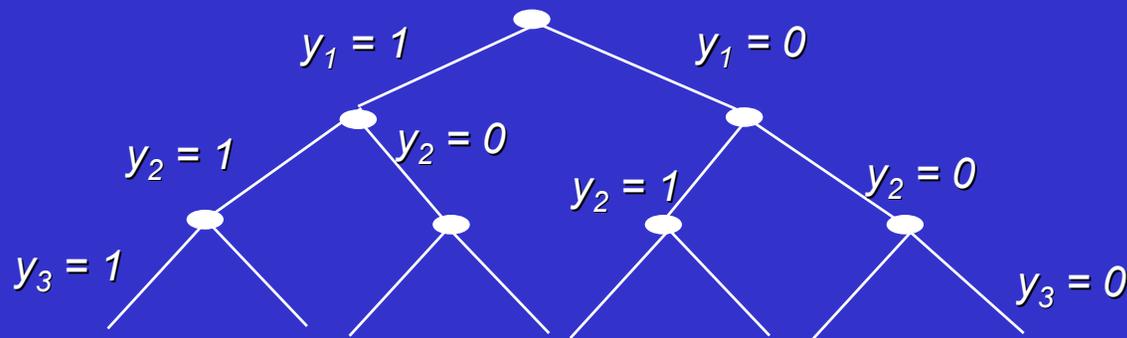
# Set Partitioning Model: Advantages and Disadvantages

- Advantages:
  - Easy to model complex work rules
    - Very few constraints
    - Linear objective function and constraints
- Disadvantages:
  - Huge number of variables- number of variables grows exponentially with the number of flights

# Problem Size

- A typical US airline (with a hub-and-spoke network) has millions or billions of potential pairings
  - Example
    - 150 flights                      90,000 pairings
    - 250 flights                      6,000,000 pairings
- Need a specialized approach to consider problems of this size

# Branch-and-Price: Branch-and-Bound for Large-Scale Integer Programs



*All possible solutions at leaf nodes of tree ( $2^n$  solutions, where  $n$  is the number of variables)*

# Column Generation

Millions/Billions of Variables

constraints

Initial

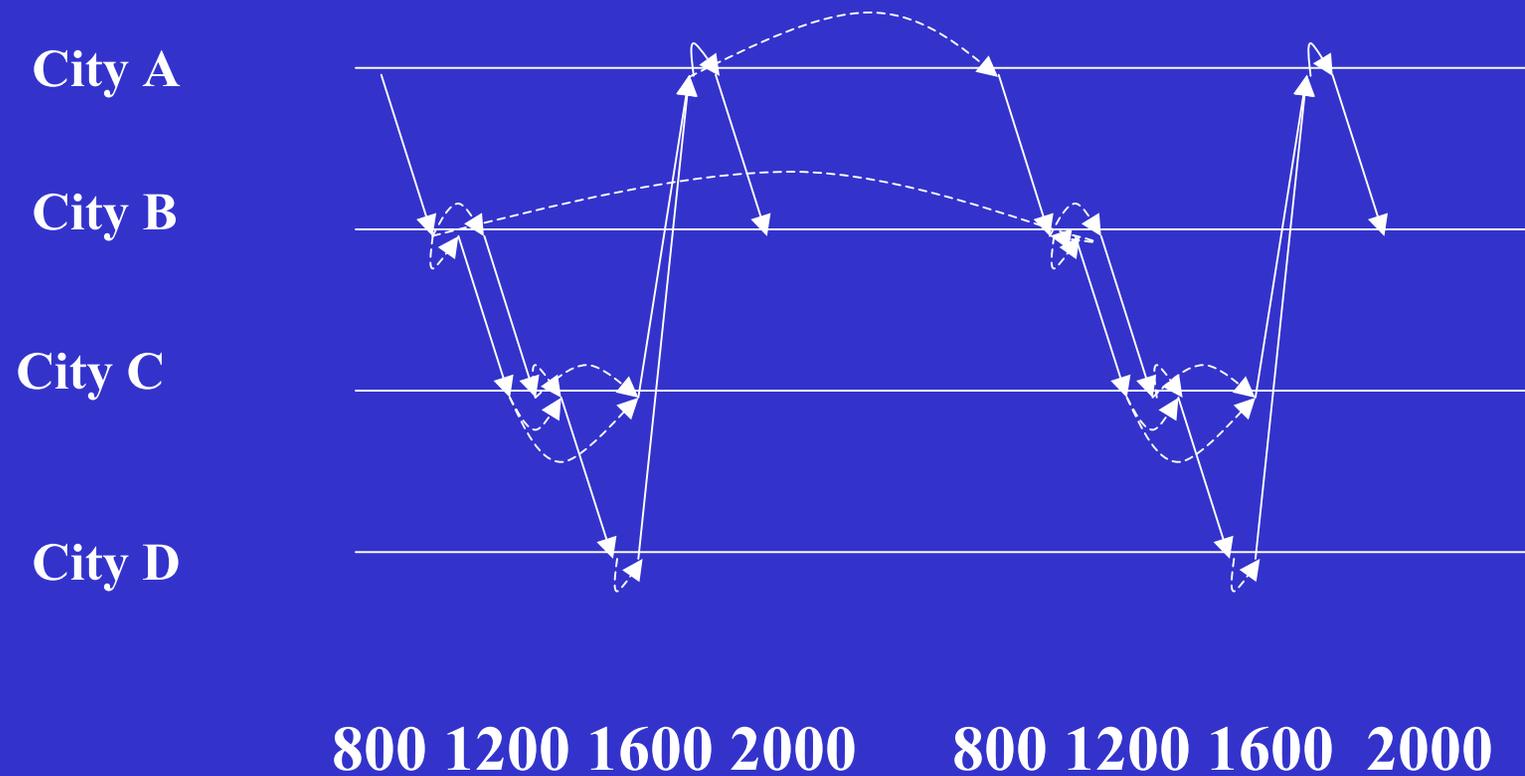
Added

Never Considered

# LP Solution: Column Generation

- Step 1: Solve Restricted Master Problem
- Step 2: Solve Pricing Problem (generate columns with negative reduced cost)
- Step 3: If columns generated, return to Step 1; otherwise STOP

# Network Representation



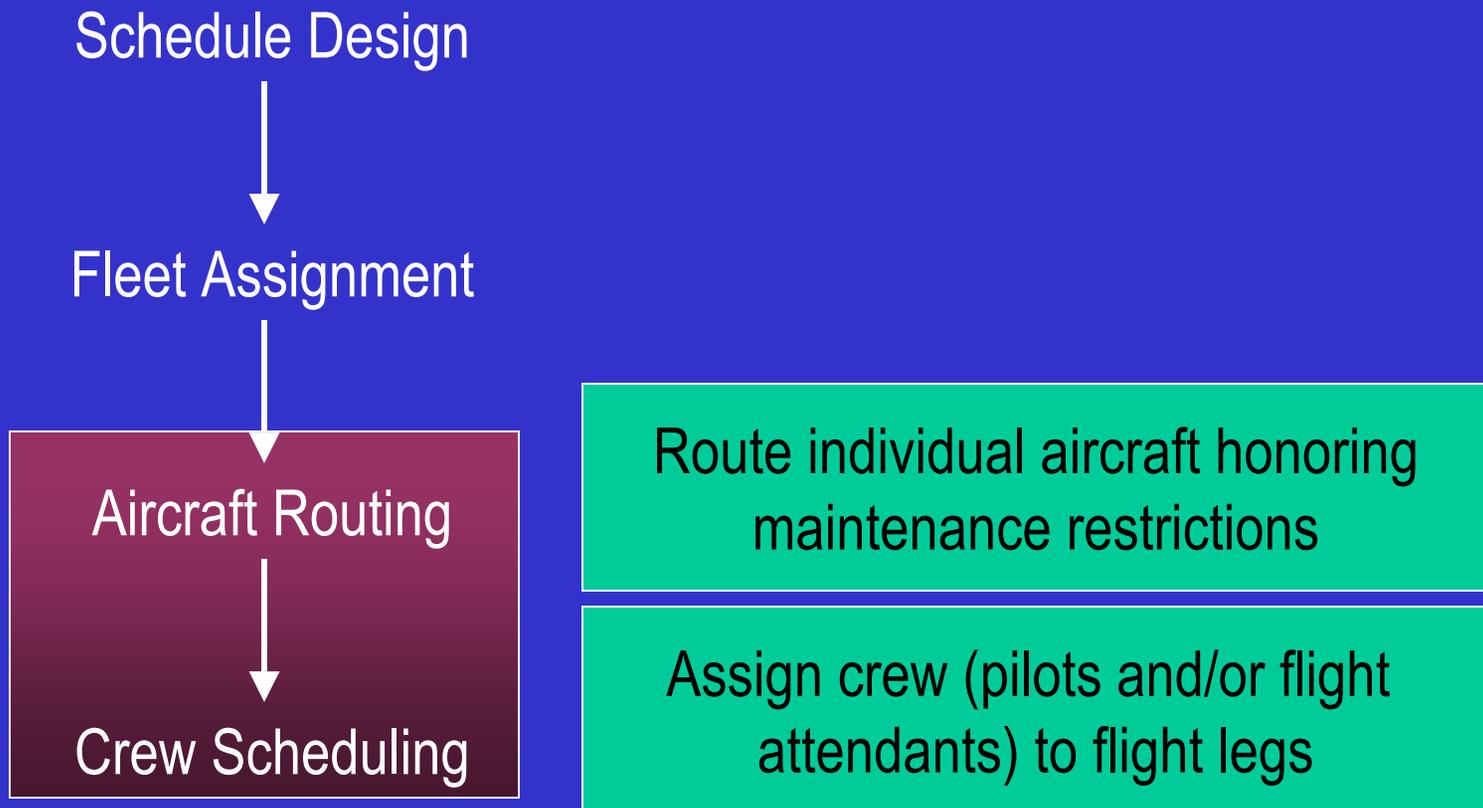
# Branch-and-Bound with Too Many Variables

- Branch-and-Price
  - Branch-and-bound with bounding provided by LP solutions
  - CP has too many variables to consider all of them
    - Solve linear programming relaxation using column generation

# A Twist...

- Crew scheduling is critical to the airline industry
  - Second largest operating expense
  - Small improvement in solution quality has significant financial impact
- For decades, researchers have worked on finding better crew scheduling *algorithms*
- Our approach is to instead improve solution quality by *expanding the feasible set of solutions*

# Airline Schedule Planning



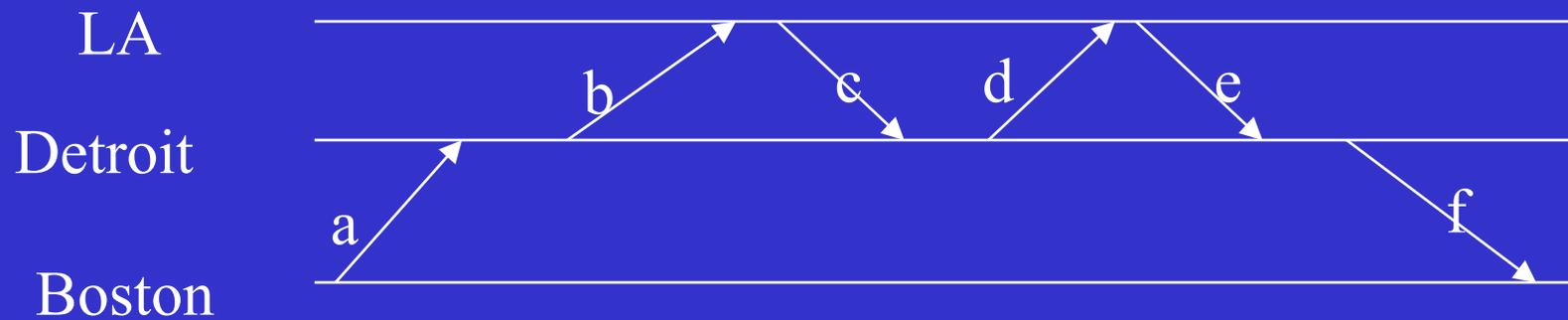
# Aircraft Maintenance Routing: Problem Definition

- Given:
  - Flight Schedule for a single fleet
    - Each flight covered exactly once by fleet
  - Number of Aircraft by Equipment Type
    - Can't assign more aircraft than are available
  - Turn Times at each Station
  - FAA Maintenance Requirements
    - Maintenance required every 60 hours of flying
    - Airlines maintain aircraft every 40-45 hours of flying with the maximum time between checks restricted to three to four calendar days

# Aircraft Maintenance Routing: Objective

- Find:
  - Feasible assignment of individual aircraft to scheduled flights
    - Each flight is covered exactly once
    - Maintenance requirements are satisfied
    - Conservation of flow (balance) of aircraft is achieved
    - The number of aircraft used does not exceed the number available

# Example: Maintenance Station in Boston



# String Model: Variable Definition

- A string is a sequence of flights beginning and ending at a maintenance station with maintenance following the last flight in the sequence
  - Departure time of the string is the departure time of the first flight in the sequence
  - Arrival time of the string is the arrival time of the last flight in the sequence + maintenance time

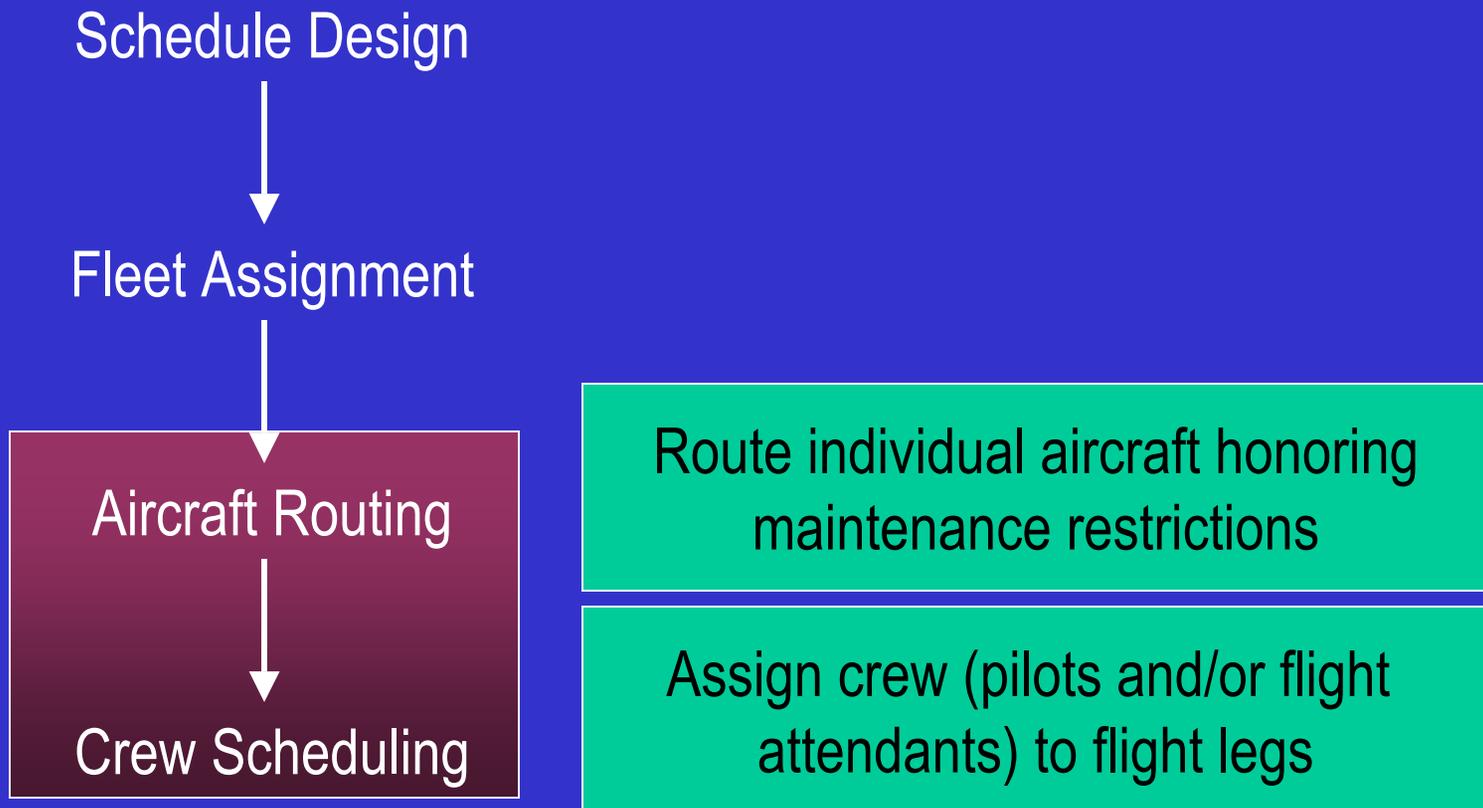
# String Model: Constraints

- Maintenance constraints
  - Satisfied by variable definition
- Cover constraints
  - Each flight must be assigned to exactly one string
- Balance constraints
  - Needed only at maintenance stations
- Fleet size constraints
  - The number of assigned aircraft cannot exceed the number of aircraft in the fleet

# Model Solution

- Complex constraints can be handled easily
- Model size
  - Huge number of variables
- Solution approach: branch-and-price
  - Generate string variables only as-needed

# Airline Schedule Planning



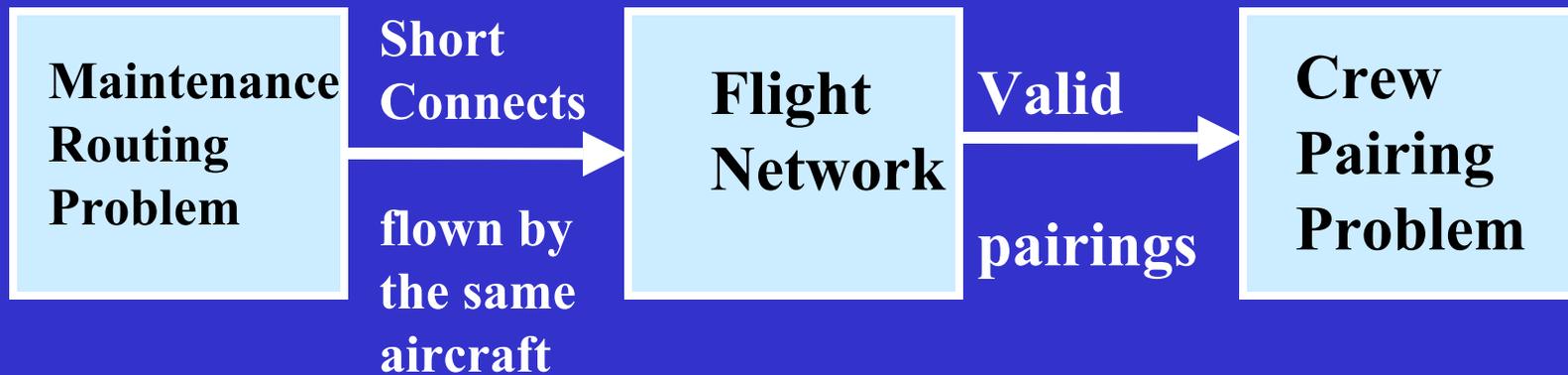
# Integrate?

- Crew scheduling options are limited by maintenance routing decisions made earlier in the airline planning process
- Solving maintenance routing and crew scheduling simultaneously yields a large and challenging problem
- Idea is to improve crew scheduling by incorporating relevant maintenance routing decisions

# Maintenance Routing and its Link to the Crew Pairing Problem

- The Maintenance Routing Problem (MR) - find *feasible* routing of aircraft ensuring adequate aircraft maintenance opportunities and flight coverage
- Crews need enough time between two sequential flights to travel through the terminal -- *minimum connect time*
- If both flights are covered by the same aircraft, connection time can be reduced
- A *short connect* is a connection that is crew-feasible only if both flights are assigned to the same aircraft

# Sequential Approach



# Klabjan, Johnson, and Nemhauser

- Solve the crew pairing problem first, including all short connects in the crew pairing network
- Given the crew solution, require all short connects included in it to be part of the maintenance solution, which is solved second
- For “good” instances, this yields the optimal solution to the integrated problem (and many problems are “good”)
- For “bad” instances, this leads to maintenance infeasibility

# Cordeau, Stojković, Soumis, Desrosiers

- Directly integrate crew and maintenance routing models
- Basic maintenance routing and crew pairing variables and constraints, plus linking constraints
- Benders decomposition approach using a heuristic branching strategy
- For non-hub-and-spoke networks, positive computational results

# Our Approach

- Generate different solutions to the maintenance routing problem
- Allow the crew pairing model to choose the maintenance routing solution with the most useful set of short connects

# The Example Again

Flights:

A B C D E F G H

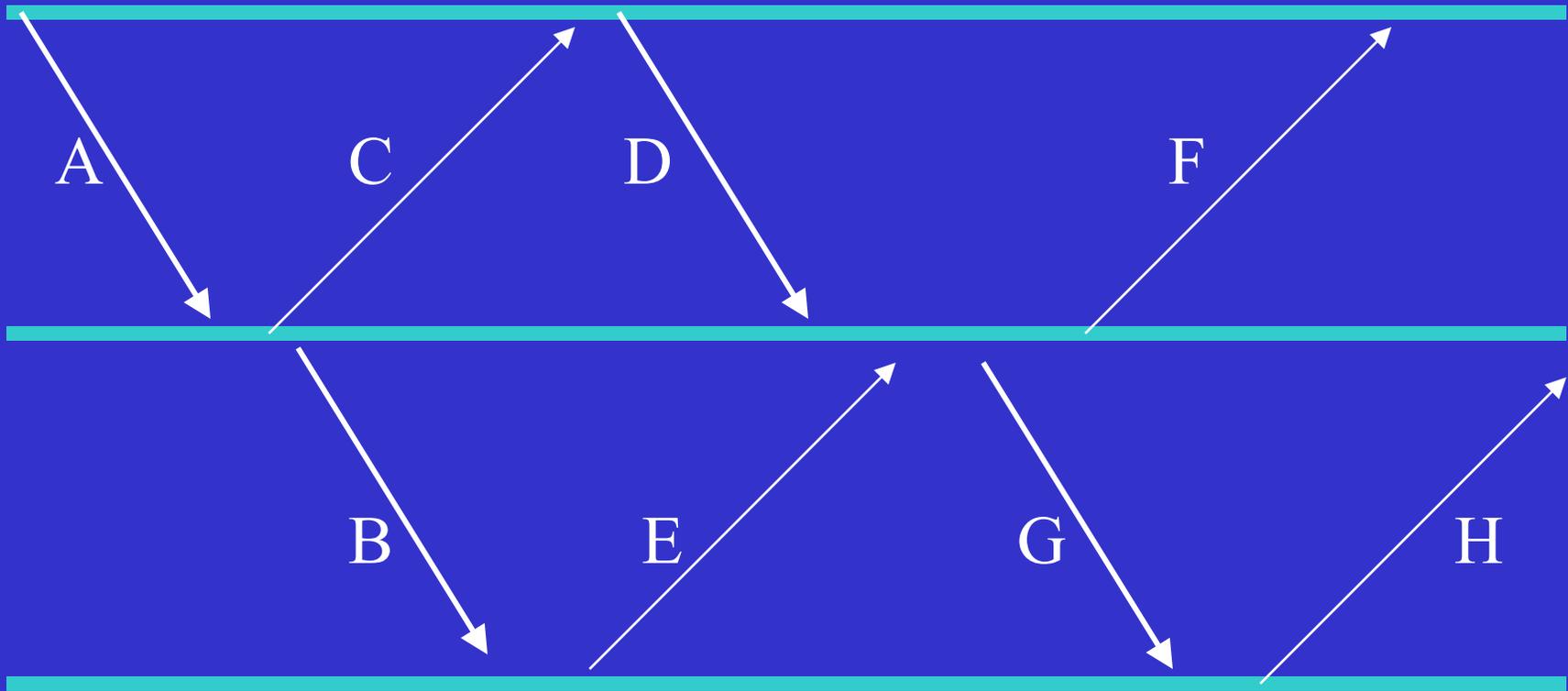
All Possible Short  
Connects:

A-B A-C E-G

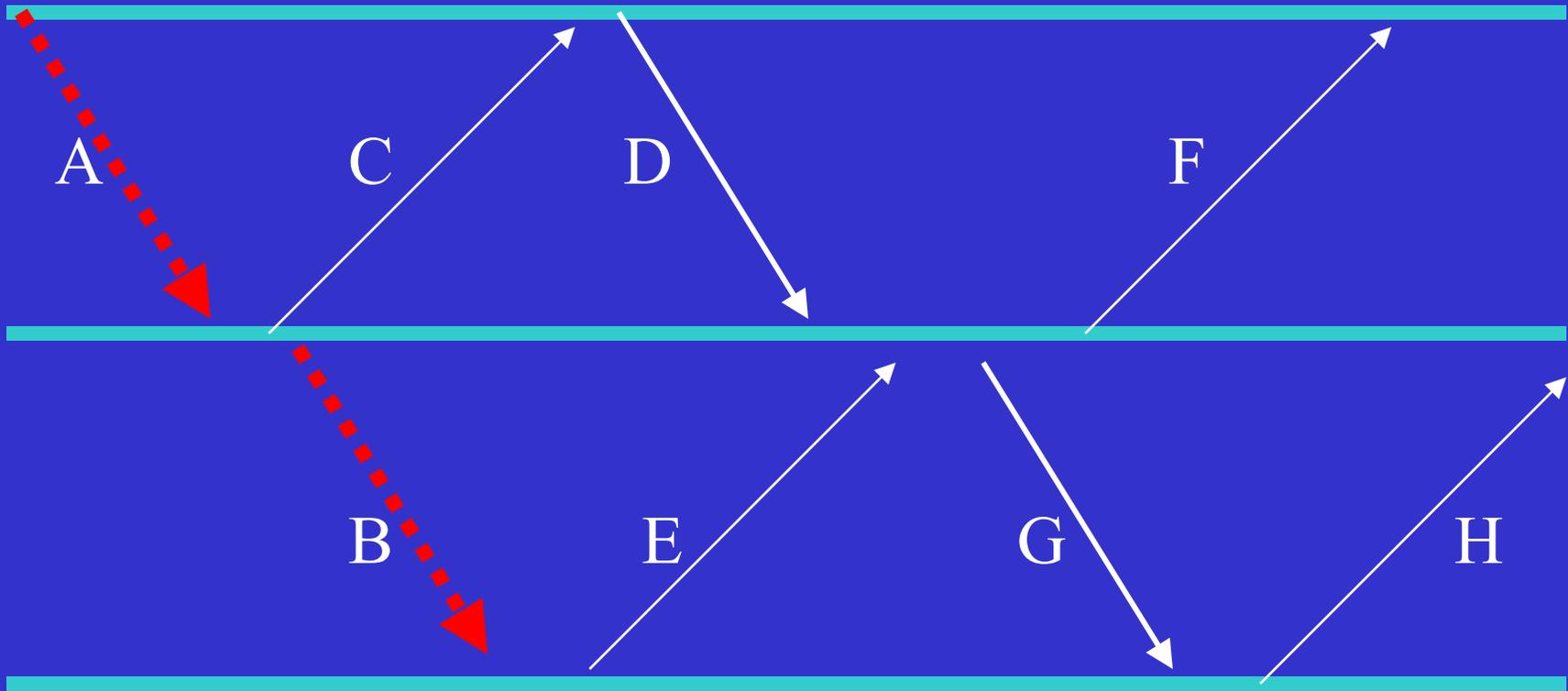
- MR solution ( $x_1$ ) assigns the same aircraft to short connects A-C and E-G
- MR solution ( $x_2$ ) assigns the same aircraft to short connect A-B

- Potential pairings:
  - A-C-D-F ( $y_1$ ): \$1
  - A-B-E-F ( $y_2$ ): \$2
  - C-D-G-H ( $y_3$ ): \$4
  - B-E-G-H ( $y_4$ ): \$6
- Crew pairing solutions:
  - $x_1 \Rightarrow$  pairings 1, 4: \$7
  - $x_2 \Rightarrow$  pairings 2, 3: \$6

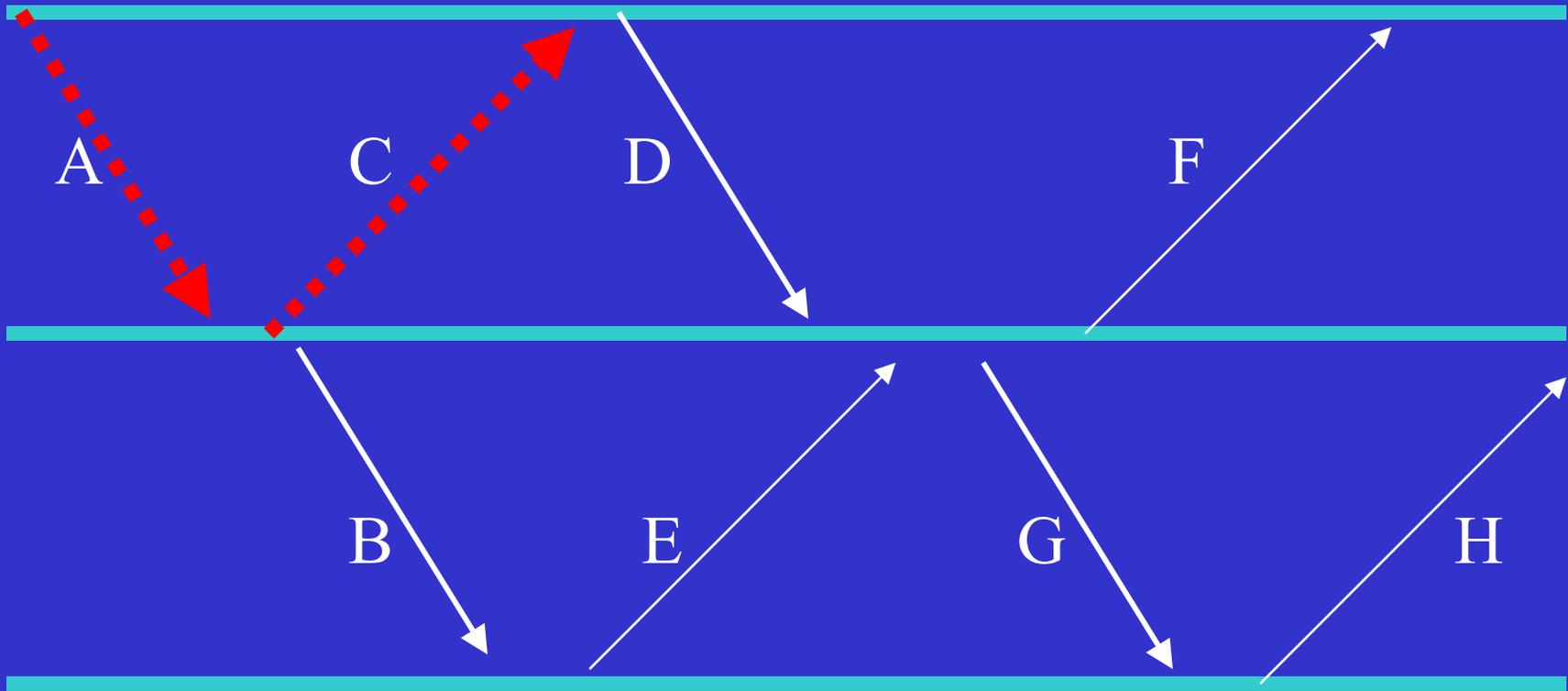
# Flights



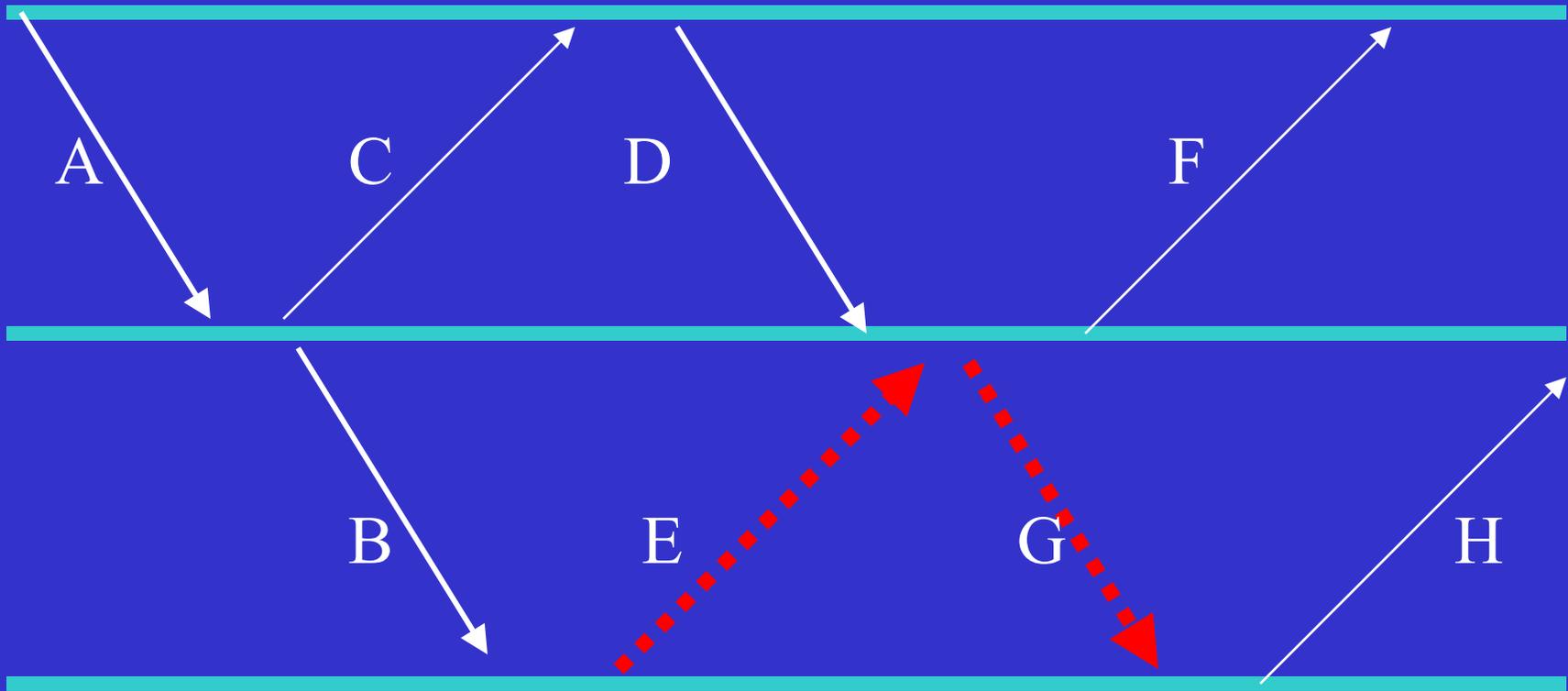
# Short Connect



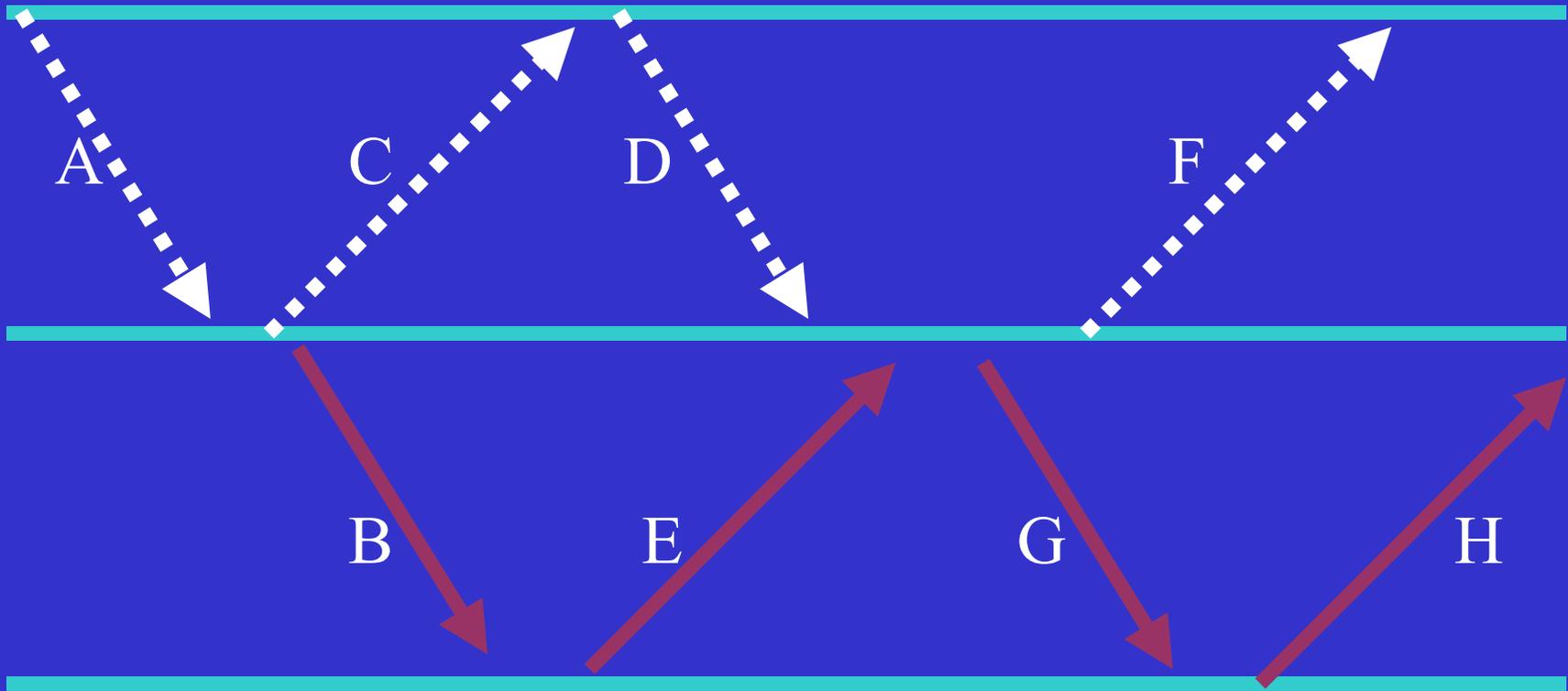
# Short Connect



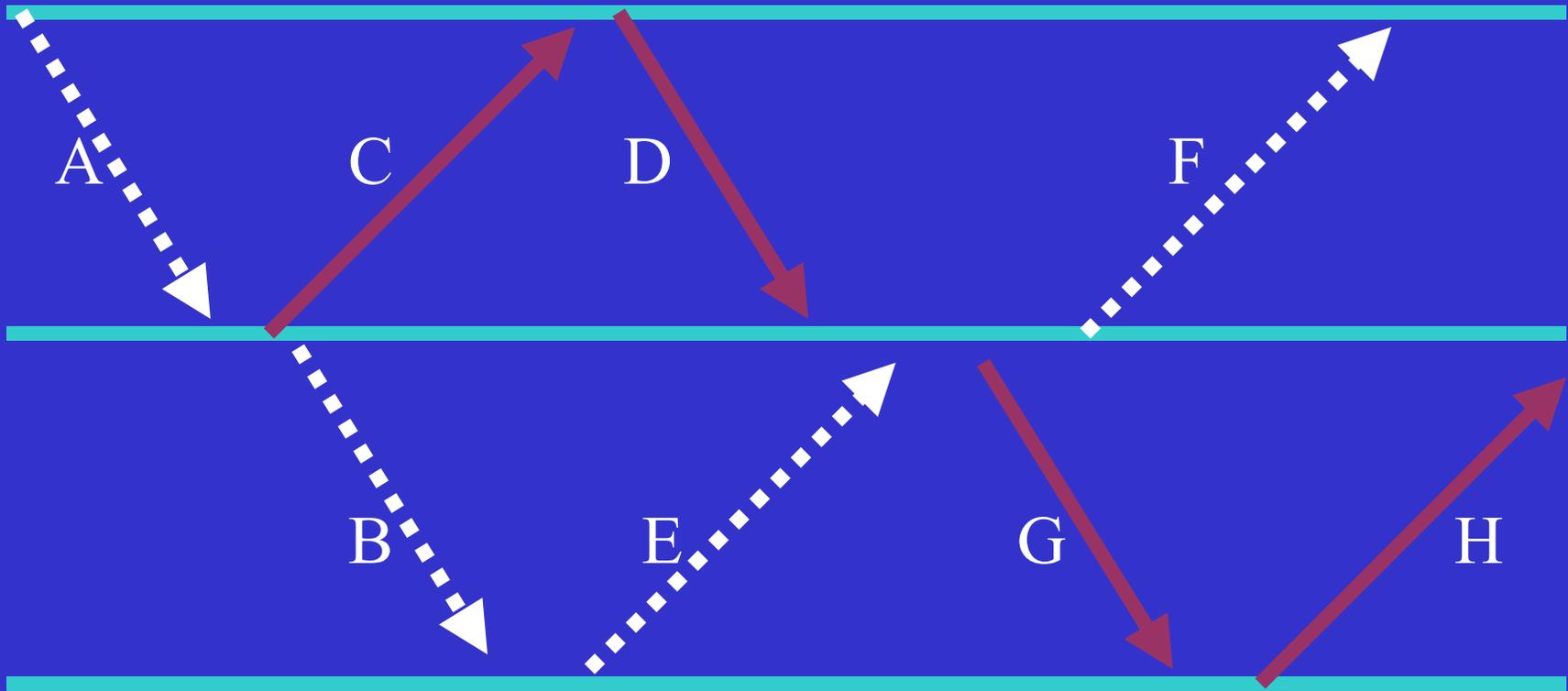
# Short Connect



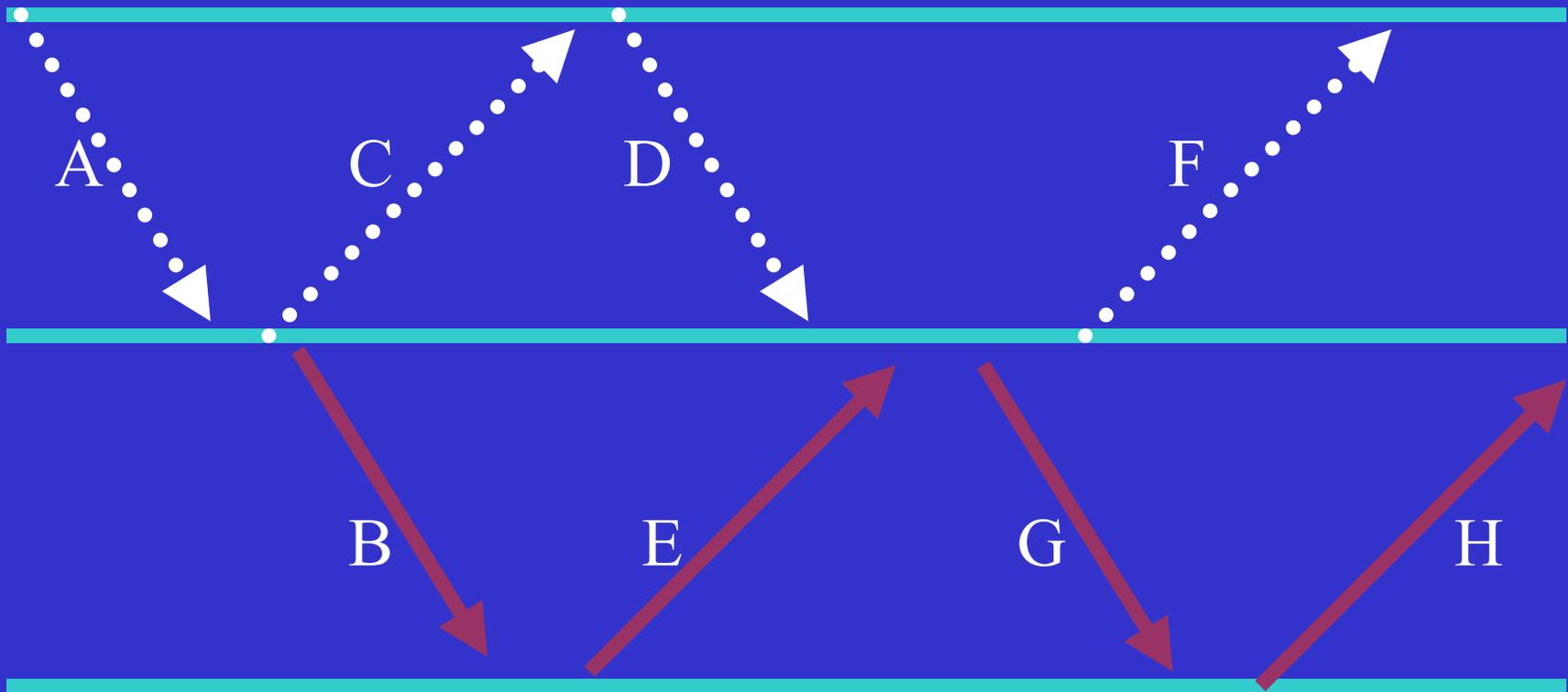
# Maintenance Solution 1



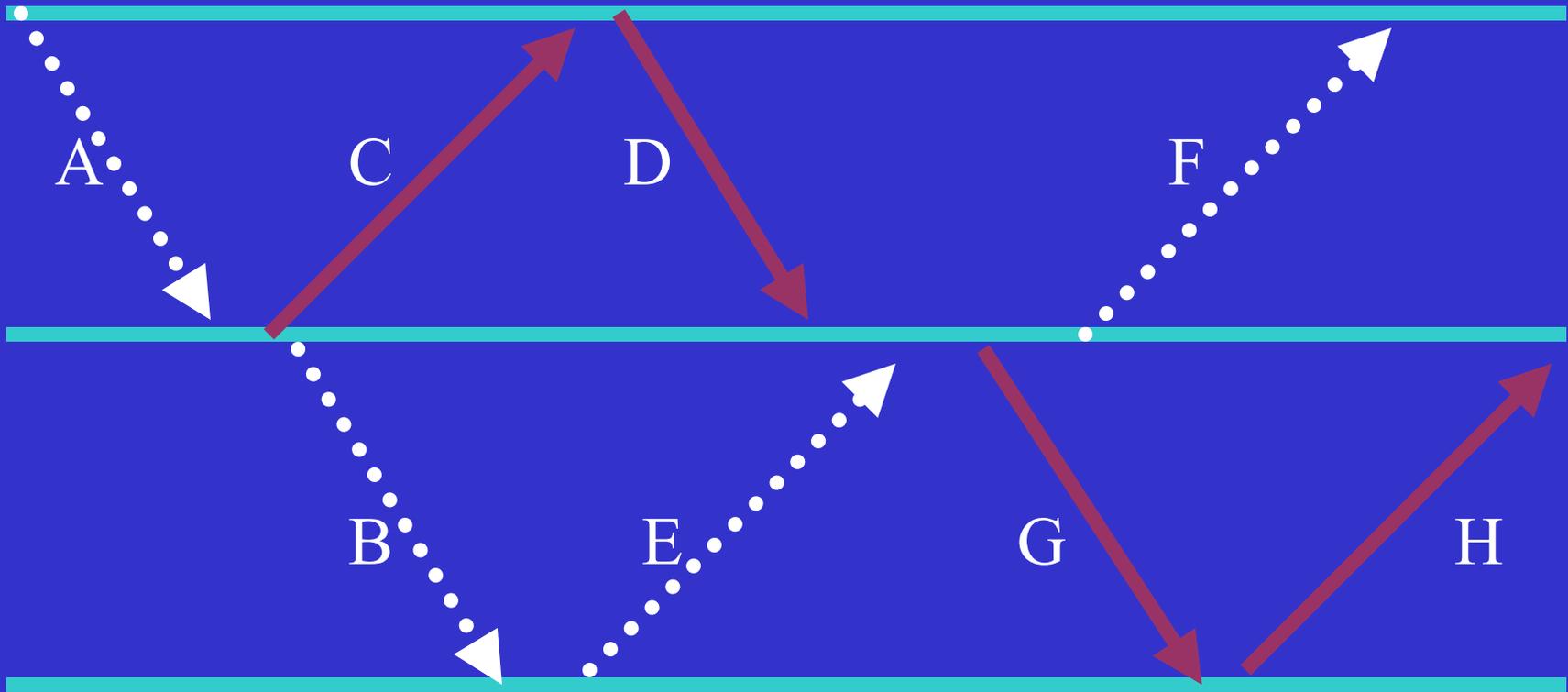
# Maintenance Solution 2



If MR Solution 1 (A-C, E-G) =>  
Optimal: Pairings 1, 4 -- \$7



If MR Solution 2 (A-B) =>  
Optimal: Pairings 2, 3 -- \$6



# Approach

- In the sequential approach, given a maintenance routing solution, the crew pairing problem is solved
- We allow the crew scheduler to choose from a collection of maintenance routing solutions
  - Select the one containing the set of short connects that allows the minimum cost crew pairing solution
- Problem: We don't want to solve one crew pairing problem for each maintenance routing solution
- Solution: Extended Crew Pairing Model (ECP)

# The Extended Crew Pairing Model (ECP)

- Simultaneously select a cost minimizing set of crew pairings and a corresponding feasible maintenance routing solution from a given set of maintenance routing solutions
- Add constraints that allow pairings with a short connect to be selected only if the chosen maintenance solution assigns the same aircraft to that short connect

# The Example Again

Flights:

A B C D E F G H

Short Connects:

A-B A-C E-G

- MR solution ( $x_1$ )  
uses short connects  
A-C and E-G
- MR solution ( $x_2$ )  
uses short connect  
A-B

- Potential pairings:
  - A-C-D-F ( $y_1$ ): \$1
  - A-B-E-F ( $y_2$ ): \$2
  - C-D-G-H ( $y_3$ ): \$4
  - B-E-G-H ( $y_4$ ): \$6
- Crew pairing solutions:
  - $x_1 \Rightarrow$  pairings 1, 4: \$7
  - $x_2 \Rightarrow$  pairings 2, 3: \$6

# Matrix Representation for the Example

	$x_1$	$x_2$	$y_1$	$y_2$	$y_3$	$y_4$	rhs	
Flights:	0	0	1	1	0	0	= 1	A
	0	0	0	1	0	1	= 1	B
	0	0	1	0	1	0	= 1	C
	0	0	1	0	1	0	= 1	D
	0	0	0	1	0	1	= 1	E
	0	0	1	1	0	0	= 1	F
	0	0	0	0	1	1	= 1	G
	0	0	0	0	1	1	= 1	H
Short Connects:	0	1	0	-1	0	0	$\geq 0$	A-B
	1	0	-1	0	0	0	$\geq 0$	A-C
	1	0	0	0	0	-1	$\geq 0$	E-G
Convexity:	1	1	0	0	0	0	= 1	Conv.

# Notation

- $P^k$  is the set of feasible pairings for fleet family  $k$
- $F^k$  is the set of flights assigned to fleet family  $k$
- $T^k$  is the set of short connects for the flights assigned to fleet family  $k$
- $S^k$  is the set of feasible MR solutions for the flights assigned to fleet family  $k$

## Notation, cont.

- $\delta_{fp}$  is 1 if flight  $f$  is included in pairing  $p$ , else 0
- $\alpha_{ts}$  is 1 if MR solution  $s$  includes short connect  $t$ , else 0
- $\beta_{tp}$  is 1 if short connect  $t$  is contained in pairing  $p$ , else 0
- $c_p$  is the cost of pairing  $p$

## Notation, cont.

- $x_s$  is a binary decision variable with value 1 if MR solution  $s$  is chosen, else 0
- $y_p$  is a binary decision variable with value 1 if pairing  $p$  is chosen, else 0

# ECP Formulation

$$\min \sum_{p \in P^k} c_p y_p$$

*st*

Flights: 
$$\sum_{p \in P^k} \delta_{fp} y_p = 1 \quad \forall f \in F^k$$

Short Connects: 
$$\sum_{s \in S^k} \alpha_{ts} x_s - \sum_{p \in P^k} \beta_{tp} y_p \geq 0 \quad \forall t \in T^k$$

Convexity: 
$$\sum_{s \in S^k} x_s = 1$$

$$x_s, y_p \in \{0,1\} \quad \forall s, p$$

# ECP Enhancements

- By exploiting dominance relationships, can dramatically reduce the **number** of MR columns considered in finding an optimal ECP solution
  - MR1 containing short connects AB, CD, GH *dominates* MR2 containing short connect AB
    - Do not need to include MR2 in ECP
- Theoretical bounds and computational observations
  - Example: 61 flights => >> 25,000 MR solutions => 4 non-dominated MR solutions (bounded by 35)
  - Can find these 4 non-dominated MR solutions by solving 4 MR problems

# ECP Enhancements, cont.

- Proof: Can relax the **integrality** of MR columns and still achieve integer solutions:
  - Same number of binary variables as original CP

$$\begin{array}{l}
 \min \sum_p c_p y_p \\
 st \\
 \sum_{p:f \in p} y_p = 1 \quad \forall f \\
 \sum_{s:t \in s} z_s - \sum_{p:t \in P} y_p \geq 0 \quad \forall t \\
 \sum_s z_s = 1 \\
 \boxed{z_s \in \{0,1\}} \quad \forall s \\
 y_p \in \{0,1\} \quad \forall p
 \end{array}
 \equiv
 \begin{array}{l}
 \min \sum_p c_p y_p \\
 st \\
 \sum_{p:f \in p} y_p = 1 \quad \forall f \\
 \sum_{s:t \in s} z_s - \sum_{p:t \in P} y_p \geq 0 \quad \forall t \\
 \sum_s z_s = 1 \\
 \boxed{z_s \geq 0} \quad \forall s \\
 y_p \in \{0,1\} \quad \forall p
 \end{array}$$

LP relaxation of ECP is tighter than LP relaxation of a basic integrated approach

# Computational Experiment

- **Problem A:**

Lower bound: 31,396.10

ECP with 16 MR columns: 31,396.10

Optimality gap: 0%

- **Problem B:**

Lower bound: 25,076.60

ECP with 20 MR columns: 25,498.60

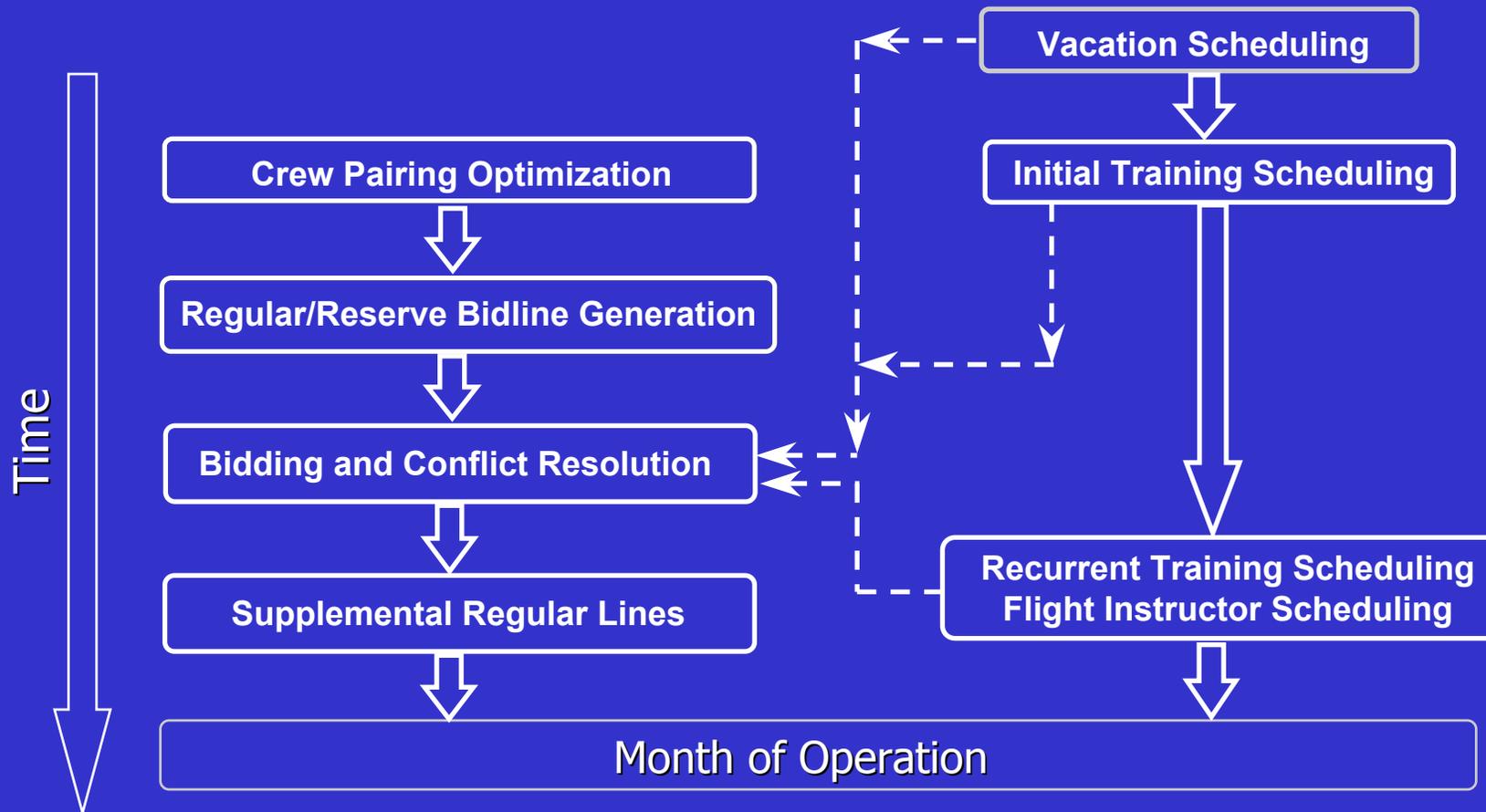
Optimality gap: 1.7%

# Airline Crew Scheduling Successes

- Excess crew costs in the planning process has been driven to 0-3%
  - AA was 8-10% 15 yrs ago: now 0-2%
  - Each 1% is worth about \$10 million/yr
    - 1997 had 9,000 pilots costing \$1.2 billion
  - Larger schedules and complex rules

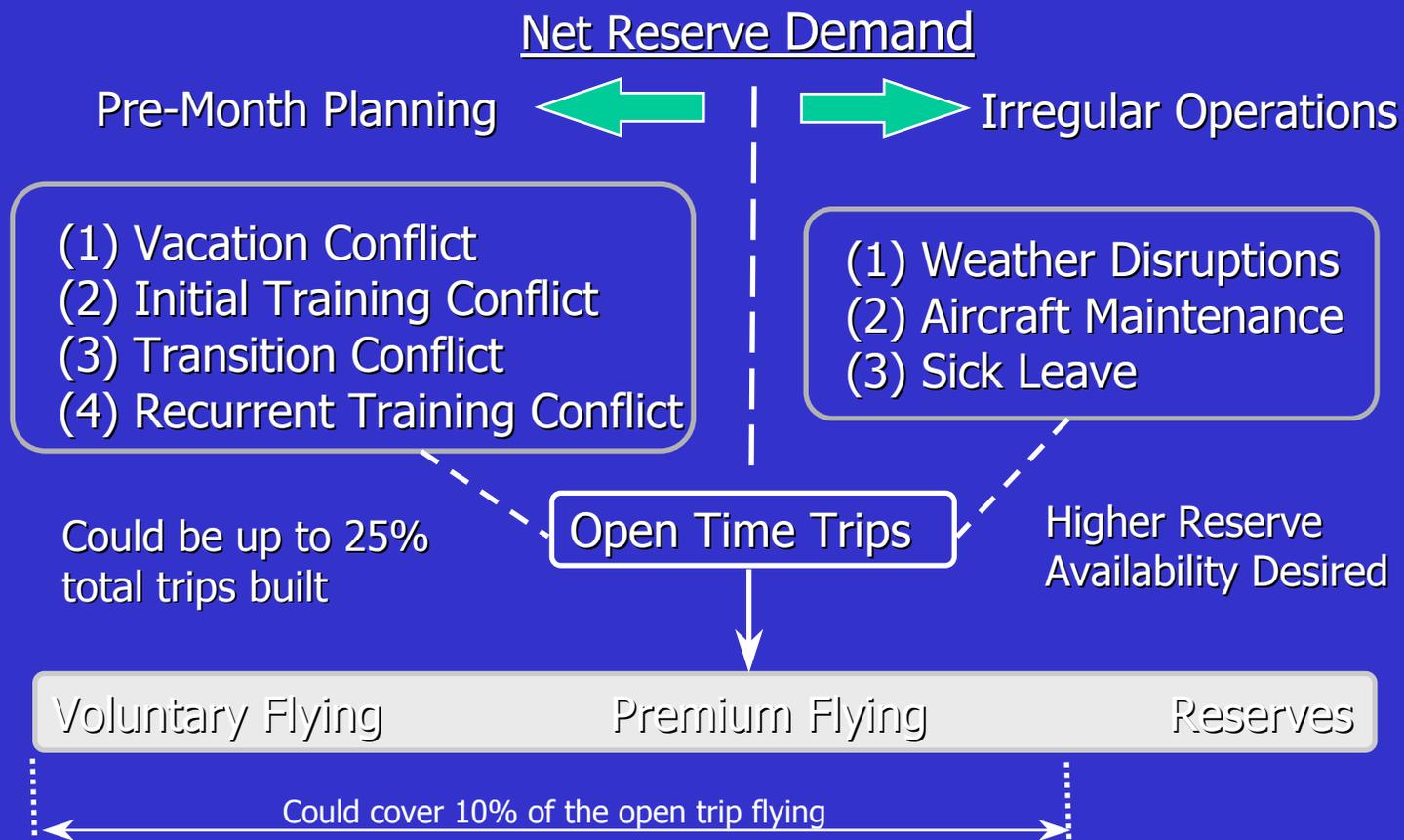
# Crew Pre-Month Planning

Ellis Johnson, Georgia Tech



# Reserve Demand

Ellis Johnson, Georgia Tech



# Crew: The Rest of The Story

Ellis Johnson, Georgia Tech

- Manpower planning, conflicts, over-time flying, and reserves
- In US airlines as high as 30% of the pilots may be on reserve bid lines
  - Actual flying is about 50% of usual
  - Of that flying, more than half is to cover conflicts and as little as 1/3 is to cover disruptions

# Conclusions

- Crew scheduling is critical to airline profitability
  - Making maintenance routing decisions independently increases costs
- A model that fully integrates MR and CP can be inflexible and difficult to solve
  - ECP exploits that only some maintenance routing information is relevant and uses dominance to reduce the size of the problem
- More work to be done... especially post-pairing optimization