6.207/14.15: Networks Lecture 11: Giant Component, Generalized Random Graphs

Outline

- Emergence and size of a giant component in Erdös-Renyi graphs
- An application: contagion and diffusion
- Generalized random graph models
- Graphs with prescribed degrees configuration model
- Emergence of a giant component in the configuration model

Reading:

- Newman, Sections 12.1-12.5, 12.7-12.8.
- Newman, Sections 13.2 (skip 13.2.2), 13.3,13.4.

Giant Component

- We have shown that when $p(n) \ll \frac{\log(n)}{n}$, the Erdös-Renyi graph is disconnected with high probability.
- In cases for which the network is not connected, the component structure is of interest.
- We have argued that in this regime the expected number of isolated nodes goes to infinity. This suggests that the Erdös-Renyi graph should have an arbitrarily large number of components.
- We will next argue that the threshold $p(n) = \frac{\lambda}{n}$ plays an important role in the component structure of the graph.
 - -~ For $\lambda < 1,$ all components of the graph are "small".
 - For $\lambda > 1$, the graph has a (unique) giant component, i.e., a component that contains a constant fraction of the nodes.

- We will analyze the component structure in the vicinity of $p(n) = \frac{\lambda}{n}$ using a branching process approximation.
- We assume $p(n) = \frac{\lambda}{n}$.
- $B(n, \frac{\lambda}{n})$: binomial random variable with parameters $n, \frac{\lambda}{n}$.
- Consider starting from node 1 and exploring the graph.



(a) Erdos-Renyi graph process.



(b) Branching Process Approx.

- $\circ~$ We first consider the case when $\lambda < 1.$
- Let Z_k^G and Z_k^B denote the number of individuals at stage k for the graph process and the branching process approximation, respectively.
- In view of the "overcounting" feature of the branching process, we have

 $Z_k^G \leq Z_k^B$ for all k.

• From branching process analysis (see Lecture 3 notes), we have

$$\mathbb{E}[Z_k^B] = \lambda^k,$$

(since the expected number of children is given by $n \times \frac{\lambda}{n} = \lambda$).

- Let S_1 denote the number of nodes in the Erdös-Renyi graph connected to node 1, i.e., the size of the component which contains node 1.
- Then, we have

$$\mathbb{E}[S_1] = \sum_k \mathbb{E}[Z_k^G] \le \sum_k \mathbb{E}[Z_k^B] = \sum_k \lambda^k = \frac{1}{1-\lambda}.$$

 $\circ~$ The preceding result suggests that for $\lambda < 1,$ the sizes of the components are "small".

Theorem

Let $p(n) = \frac{\lambda}{n}$ and assume that $\lambda < 1$. For all (sufficiently large) a > 0, we have

$$\mathbb{P}\Big(\max_{1\leq i\leq n}|S_i|\geq a\log(n)\Big)\to 0 \quad as \ n\to\infty.$$

Here $|S_i|$ is the size of the component that contains node *i*.

- This result states that for $\lambda < 1$, all components are small [in particular they are of size $O(\log(n))$].
- Proof is beyond the scope of this course.

- $\circ~$ We next consider the case when $\lambda>1.$
- We claim that $Z_k^G \approx Z_k^B$ when $\lambda^k \leq O(\sqrt{n})$.
- The expected number of conflicts at stage k + 1 satisfies

 $\mathbb{E}[\text{number of conflicts at stage } k+1] \approx np^2 \mathbb{E}[Z_k^2] = n \frac{\lambda^2}{r^2} \mathbb{E}[Z_k^2].$



• We assume for large *n* that Z_k is a Poisson random variable and therefore $var(Z_k) = \lambda^k$. This implies that

$$\mathbb{E}[Z_k^2] = \operatorname{var}(Z_k) + \mathbb{E}[Z_k]^2 = \lambda^k + \lambda^{2k} \approx \lambda^{2k}.$$

• Combining the preceding two relations, we see that the conflicts become non-negligible only after $\lambda^k \approx \sqrt{n}$.

- Hence, there exists some c > 0 such that $\mathbb{P}(\text{there exists a component with size } \geq c\sqrt{n} \text{ nodes}) \rightarrow 1 \text{ as}$ $n \rightarrow \infty$.
- Moreover, between any two components of size \sqrt{n} , the probability of having a link is given by

$$\mathbb{P}(ext{there exists at least one link}) = 1 - (1 - rac{\lambda}{n})^n pprox 1 - e^{-\lambda}$$
,

i.e., it is a positive constant independent of n.

• This argument can be used to see that components of size $\leq \sqrt{n}$ connect to each other, forming a connected component of size qn for some q > 0, a giant component.

Size of the Giant Component

- Form an Erdös-Renyi graph with n-1 nodes with link formation probability $p(n) = \frac{\lambda}{n}$, $\lambda > 1$.
- Now add a last node, and connect this node to the rest of the graph with probability p(n).
- Let q be the fraction of nodes in the giant component of the n 1 node network. We can assume that for large n, q is also the fraction of nodes in the giant component of the n-node network.
- The probability that node *n* is not in the giant component is given by $\mathbb{P}(\text{node } n \text{ not in the giant component}) = 1 - q \equiv \rho.$

• The probability that node *n* is not in the giant component is equal to the probability that none of its neighbors is in the giant component, yielding

$$\rho = \sum_{k=0}^{n-1} p_k \rho^k \equiv \Phi(\rho).$$

• Like before, this equation has a fixed point $\rho^* \in (0, 1)$.

An Application: Contagion and Diffusion

- Consider a society of *n* individuals.
- A randomly chosen individual is infected with a contagious virus.
- Assume that the network of interactions in the society is described by an Erdös-Renyi graph with link probability p.
- Assume that any individual is immune with a probability π .
- We would like to find the expected size of the epidemic as a fraction of the whole society.
- The spread of disease can be modeled as:
 - Generate an Erdös-Renyi graph with *n* nodes and link probability *p*.
 - Delete πn of the nodes uniformly at random.
 - Identify the component that the initially infected individual lies in.
- We can equivalently examine a graph with $(1 \pi)n$ nodes with link probability p.

An Application: Contagion and Diffusion

- We consider 3 cases:
- $p(1-\pi)n < 1$:

 $\mathbb{E}[\text{size of epidemic as a fraction of the society}] \leq \frac{\log((1-\pi)n)}{n} \approx 0.$

• $1 < p(1-\pi)n < \log((1-\pi)n)$:

$$\begin{split} \mathbb{E}[\text{size of epidemic as a fraction of the society}] \\ &= \frac{qq(1-\pi)n + (1-q)\log((1-\pi)n))}{n} \approx q^2(1-\pi), \end{split}$$

where q denotes the fraction of nodes in the giant component of the graph with $(1 - \pi)n$ nodes, i.e., $q = 1 - e^{-q(1 - \pi)np}$.

 $\circ \ p > \frac{\log((1-\pi)n)}{(1-\pi)n}:$

 $\mathbb{E}[\text{size of epidemic as a fraction of the society}] = (1 - \pi).$

Configuration Model—1

- We have seen that the Erdös-Renyi model has a Poisson degree distribution, which falls off very fast.
- Our next goal is to generate random networks with a "given degree distribution".
- One of the most widely method used for this purpose is the configuration model developed by Bender and Canfield in 1978.
- The configuration model is specified in terms of a degree sequence, i.e., for a network of *n* nodes, we have a desired degree sequence (k_1, \ldots, k_n) , which specifies the degree k_i of node *i*, for $i = 1, \ldots, n$.
 - Given a degree distribution p_k , we can generate the degree sequence for *n* nodes by sampling the degrees independently from the distribution p_k , i.e., $k_i \sim p_k$.
 - A law of large numbers argument establishes that the frequency of degrees $p_k^{(n)}$ converges to the degree distribution p_k as n goes to infinity.

Configuration Model—2

- Given the degree k_i for node *i* for all i = 1, ..., n, we create a random network with these degrees as follows:
- We give each node *i*, k_i "stubs" sticking out of it, which are ends of edges-to-be (there are a total of $\sum_i k_i = 2m$ stubs, where *m* is the number of edges).
- We choose two stubs uniformly at random and create an edge between the corresponding nodes.
- We choose another pair from the remaining 2m 2 stubs, connect those and continue until all the stubs are used up.
- Remarks:
 - This process generates each possible matching of stubs with equal probability.
 - The sum of degrees needs to be even (or else an entry will be left out at the end).
 - It is possible to have self-edges and multiedges.

Distribution of the Degree of a Neighboring Node—1

- We will use a branching process approximation to study the giant component in the configuration model.
- For this we need to understand the distribution of the degree of a neighboring node, i.e., given some node *i* with degree *d_i*, consider a neighbor *j*. What is the degree distribution of node *j*?



- Naive intuition: Same distribution as node *i*.
- Example: Consider a graph with 4 nodes and links $\{1,2\}$, $\{2,3\}$, $\{3,4\}$.
 - We have $p_1 = p_2 = 1/2$. Pick a link at random, then randomly pick an end of it, there is a 2/3 chance of finding a node with degree 2 and 1/3 chance of finding a node with degree 1.
 - Higher degree nodes are involved in a higher percentage of the links.

Distribution of the Degree of a Neighboring Node—2

- The degree of a node we reach by following a randomly chosen edge is not given by p_k .
- In the configuration model, an edge emerging from a node has equal chance of terminating at any of the stubs.
- Since there are 2m stubs in total, the probability of this edge ending at any particular node of degree k is k/2m.
- Since the total number of nodes with degree k is given by np_k , the probability of the edge attaching to a node with degree k is given by

$$\frac{k}{2m}np_k = \frac{kp_k}{\langle k \rangle},$$

where $\langle k \rangle$ is the expected degree in the network and the equality follows from the relation $2m = n \langle k \rangle$.

Distribution of the Degree of a Neighboring Node—3

- Intuitively, there are k edges that arrive at a node of degree k, we are k times as likely to arrive at that node than another node that has degree 1.
- Thus, the degree distribution of the neighboring node \tilde{p}_k is proportional to kp_k , $kp_k = kp_k$

$$ilde{p}_k = rac{kp_k}{\sum_j jp_j} = rac{kp_k}{\langle k \rangle}.$$



Emergence of a Giant Component in the Configuration Model—1

- We will use a branching process approximation to analyze the emergence of the giant component.
 - We ignore self loops (can be shown to have small probability) and conflicts (do not matter until the graph grows to a substantial size).
- Note that we have

$$\mathcal{U} = \tilde{\mathbb{E}}[\text{number of children}] = \tilde{\mathbb{E}}[k-1]$$
$$= \sum_{k} k \tilde{p}_{k} - 1$$
$$= \sum_{k} \frac{k^{2} p_{k}}{\langle k \rangle} - 1$$
$$= \frac{\langle k^{2} \rangle}{\langle k \rangle} - 1.$$

Emergence of a Giant Component in the Configuration Model—2

 Using the branching process analysis, this yields the following threshold for the emergence of the giant component:

Subcritical: $\mu < 1$, or equivalently

$$rac{\langle k^2
angle}{\langle k
angle} < 2 \quad \Leftrightarrow \quad \langle k(k-2)
angle < 0.$$

Supercritical: $\mu > 1$, or equivalently

$$\langle k(k-2) \rangle > 0.$$

• In the case of an Erdös-Renyi graph, we have $\langle k^2 \rangle = \langle k \rangle + \langle k \rangle^2$, and so the giant component emerges when

$$\langle k \rangle^2 > \langle k \rangle \quad \Leftrightarrow \quad \langle k \rangle > 1.$$

• Since $\langle k \rangle = (n-1)p$ in the Erdös-Renyi graph, this indeed yields the threshold function $t(n) = \frac{1}{n}$ for the emergence of the giant component.

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