

Heuristics



- What are they?
- Why use them?
- Different kinds of Heuristics
- Examples of strengths and weaknesses

What is a Heuristic?



- Any method that is not mathematically proven to find an optimal solution.
- Does not guarantee an optimal solution
- Generally give answers
 - ▶ faster than exact methods
 - ▶ or with less coding effort

Why Use Heuristics



■ Obvious reasons

- ▶ Faster answers, less work

■ Less Obvious reasons

- ▶ Data is approximate or uncertain
- ▶ Model is approximate
- ▶ Gaps between recommendation and implementation
 - Why get exact answers to an inexact problem?

Why Not Solve Exactly

- Why not solve inexact problem exactly
 - ▶ Exact optimum may be very sensitive to the data
 - if the data are off a little, the exact answer could be off by a lot
 - ▶ Exact optimum may be more sensitive to disruption
 - Exact optimum may fail dramatically when faced with a local disruption

We always use heuristics



- Optimizing component behavior rarely optimizes system behavior
- Example:
 - ▶ Purchasing minimizing piece cost
 - ▶ Transportation minimizing freight cost
 - ▶ Manufacturing minimizing changeover costs
 - ▶ ...

Fundamental Heuristics

- Hierarchical approach to systems
 - ▶ Strategic: Setting Direction
 - ▶ Tactical: Plotting a course
 - ▶ Operational: Staying on course
- Hierarchical solution strategies
 - ▶ Ford Finished Vehicles
 - Allocate dealers to ramps
 - Route vehicles to ramps
 - ▶ Airlines
 - Assign body types to flight legs
 - Build routes for the planes
 - Build routes for the crews

Basic Forms



■ Construction heuristics

- ▶ Build a solution by making a sequence of good (local) decisions

■ Improvement heuristics

- ▶ Improve an existing solution by making a sequence of good (local) modifications to an existing solution

Examples



■ Traveling Salesman Problem

- ▶ Visit each city once and return to the starting point at minimum cost
- ▶ Applications:
 - Routing
 - Robotic assembly
 - Production scheduling
 - ...

Construction Heuristics

- Unalloyed Greed: Nearest Neighbor
 - ▶ Start somewhere
 - ▶ Go next to the closest as yet unvisited city
 - ▶ Nearest Neighbor Example
(<http://itp.nat.uni-magdeburg.de/~mertens/TSP/>)
 - ▶ Fast, easy to implement, but
 - ▶ Myopic behavior leads to problems

Construction Heuristics

- **Tempered Greed: Farthest Insertion**
 - ▶ Start with a circumnavigating tour
 - ▶ Adjust the tour in the best way to include the worst excluded city
 - ▶ **Farthest Insertion Example**
(<http://itp.nat.uni-magdeburg.de/~mertens/TSP/>)
 - ▶ Fast, easy to implement,
 - ▶ Tends to do better

Idea



- Construction heuristics are fast
- Temper myopic behavior with global perspective
- Dangers
 - ▶ It's not always easy to find a feasible solution!
 - ▶ Obvious improvements possible

Improvement Heuristics



- Start with a feasible solution and iteratively improve it.
- Generally hard to spot improvements to the resulting solutions
- Generally fast, but this depends
- Not always best to begin with a good solution

Improvement Heuristics

- 2-Opt: Break the tour and glue it back in a better way
 - ▶ Start with a tour
 - ▶ Break two edges
 - ▶ Put the pieces together in the best way
 - ▶ 2-Opt Example

(<http://itp.nat.uni-magdeburg.de/~mertens/TSP>)

Advanced Improvement



- Finding an improvement might be a (smaller, faster, easier) optimization problem:
 - ▶ Machine Scheduling:
 - Start with a schedule
 - Remove one job and reinsert it into the schedule in the best way
 - Repeat

Local Optima



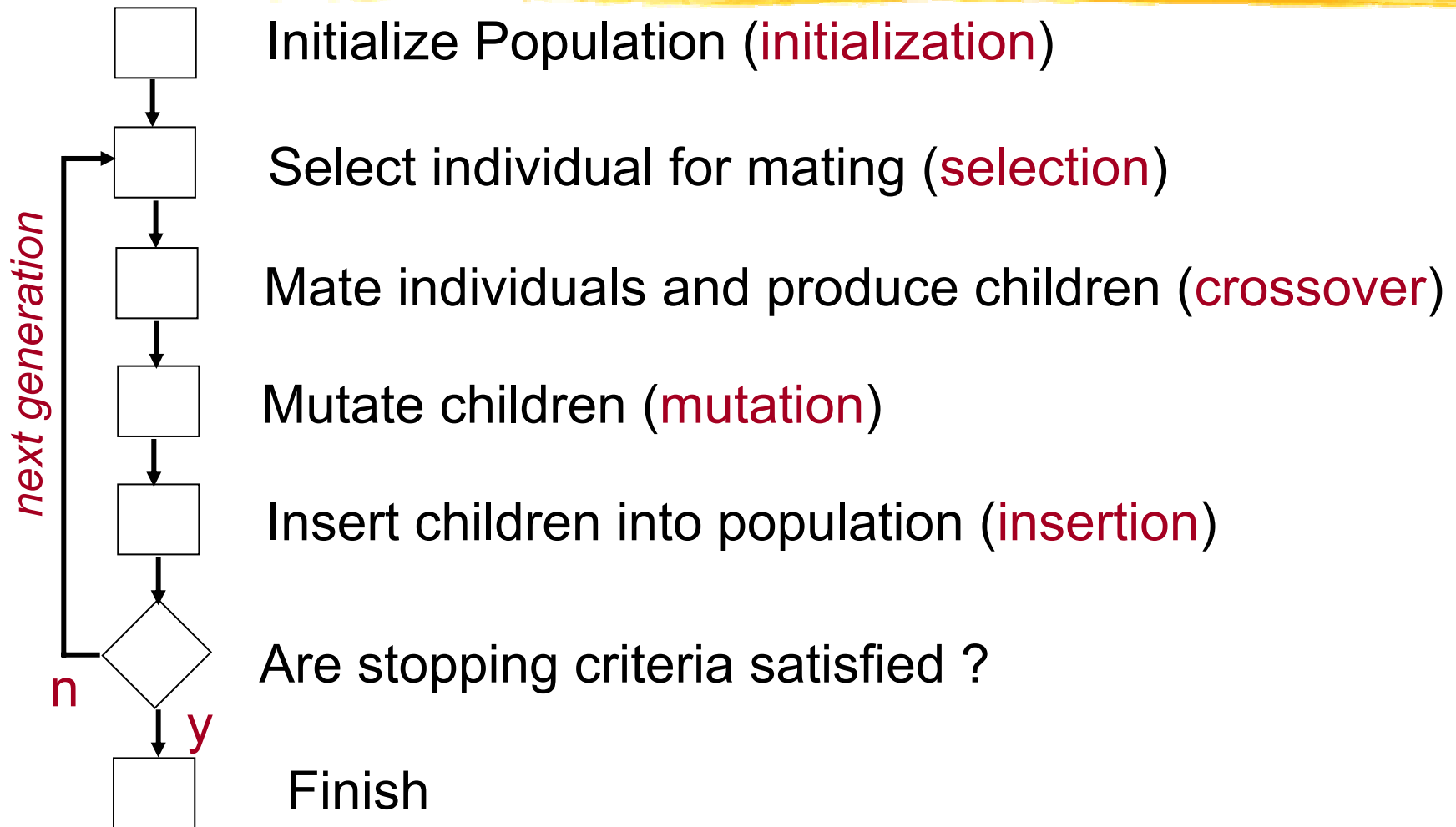
- Construction and Improvement Heuristics find local optima.
- Increasing the neighborhood to look over avoids some of these but increases the time required
 - ▶ 3-Opt
 - ▶ K-Opt
 - ▶ Reschedule 2 jobs at a time, ...

Avoiding Local Optima



- Randomized methods
- Simulated annealing
 - ▶ move to a worse solution with some probability
 - ▶ this probability declines with time
- Genetic Algorithms
 - ▶ Using a population of feasible solutions

Genetic Algorithms



Genetic Algorithms



■ GA Example

(<http://cs.felk.cvut.cz/%7Exobitko/ga/index.html>)

- Quickly becoming an off-the-shelf tool
- Key to success for heuristics – quick implementation
- An implementation included with Premium Solver from Frontline Systems, Inc. (Included with text).

Eliminate the Problem!



- Self organizing mechanisms

- Examples:

- ▶ Honey bee forager allocation

- No one is in charge. The process works

- ▶ Bucket brigades

- ▶ http://www.isye.gatech.edu/people/faculty/John_Bartholdi/bucket-brigades.html

Honey Bees



- Allocate Foragers to flower patches
- Quality and locations change
- Survival depends on good answers
- No
 - ▶ Satellites
 - ▶ Cell phones
 - ▶ Supercomputers
 - ▶ Centralized control mechanism

Like the Ants



- Waggle dance advertises site
- Faster bees dance more often
- Assume:
 - ▶ More bees gather more quickly in total
 - ▶ But more slowly on average
- Converge to equal time allocation
 - ▶ Every active site takes the same time to collect from

This Answer is Good



- In the worst case, with satellites, cell phones and supercomputers, bees could do at most twice as well
- In the likely case (specific reasonable models of cost) it's more like 5% better
- The analytical approach of setting derivatives to 0 can be arbitrarily bad

Performance Guarantees



- Some heuristics do come with guarantees
- The answer from the heuristic is within $X\%$ of the best answer
- Value: reassuring
- Disadvantage: Require assumptions or simplification

A Heuristic with guaranteed performance for The TSP

■ A Spanning Tree

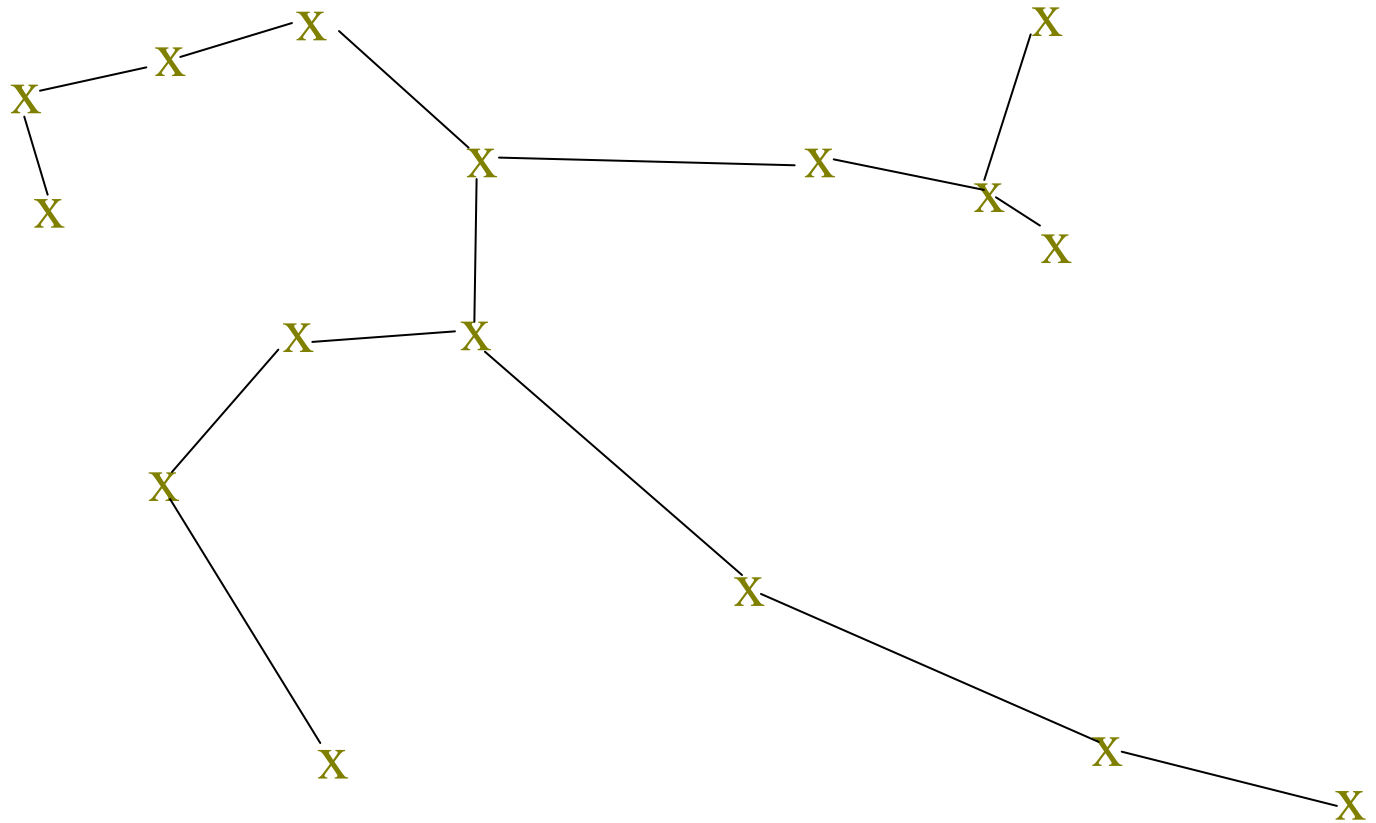
▶ Use the Greedy Algorithm

- Add edges in increasing order of length
- Discard any that create a cycle

■ Is a Lower bound on the TSP

- ▶ Drop one edge from the TSP and you have a spanning tree
- ▶ It must be at least as long as the minimum spanning tree

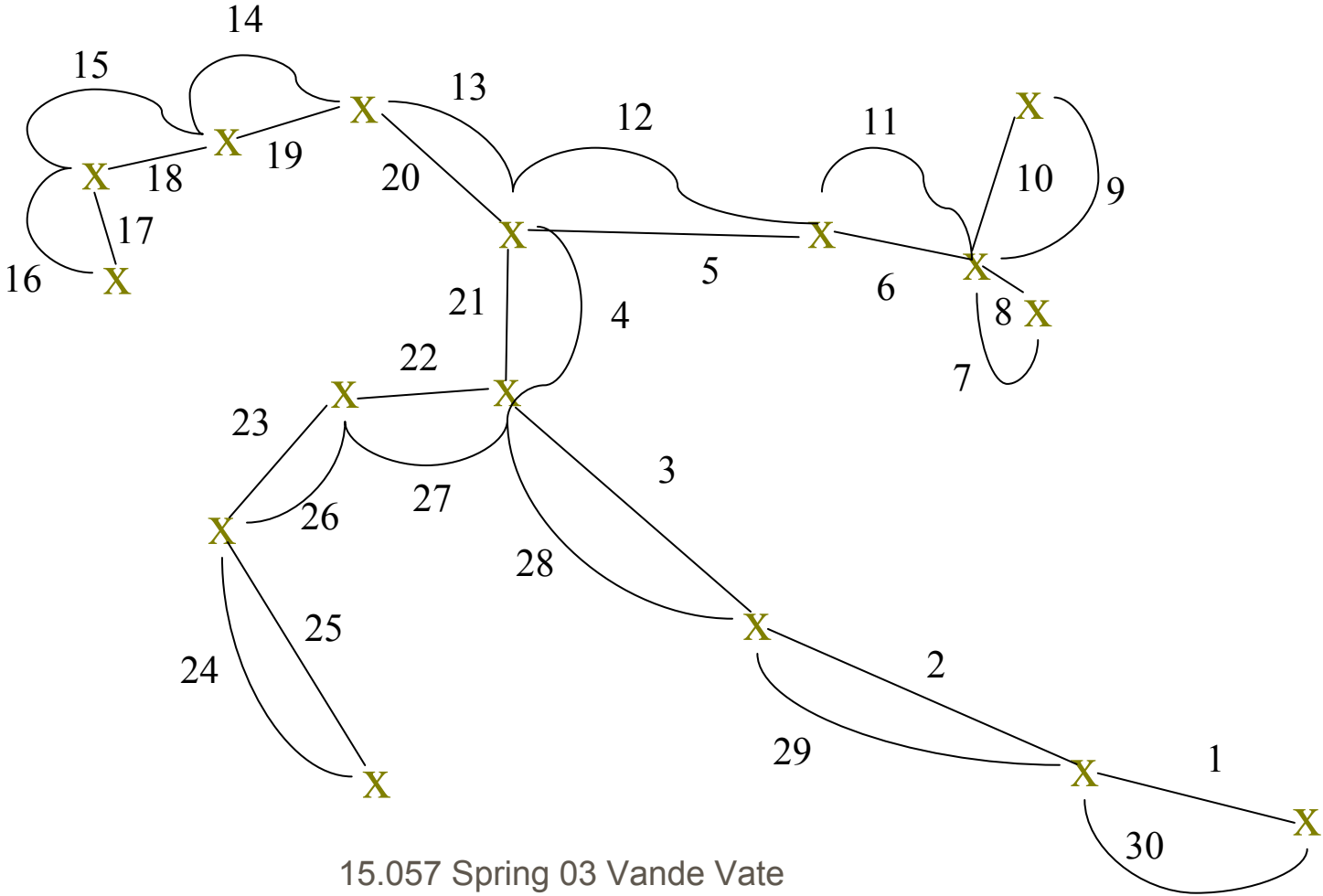
The Spanning Tree



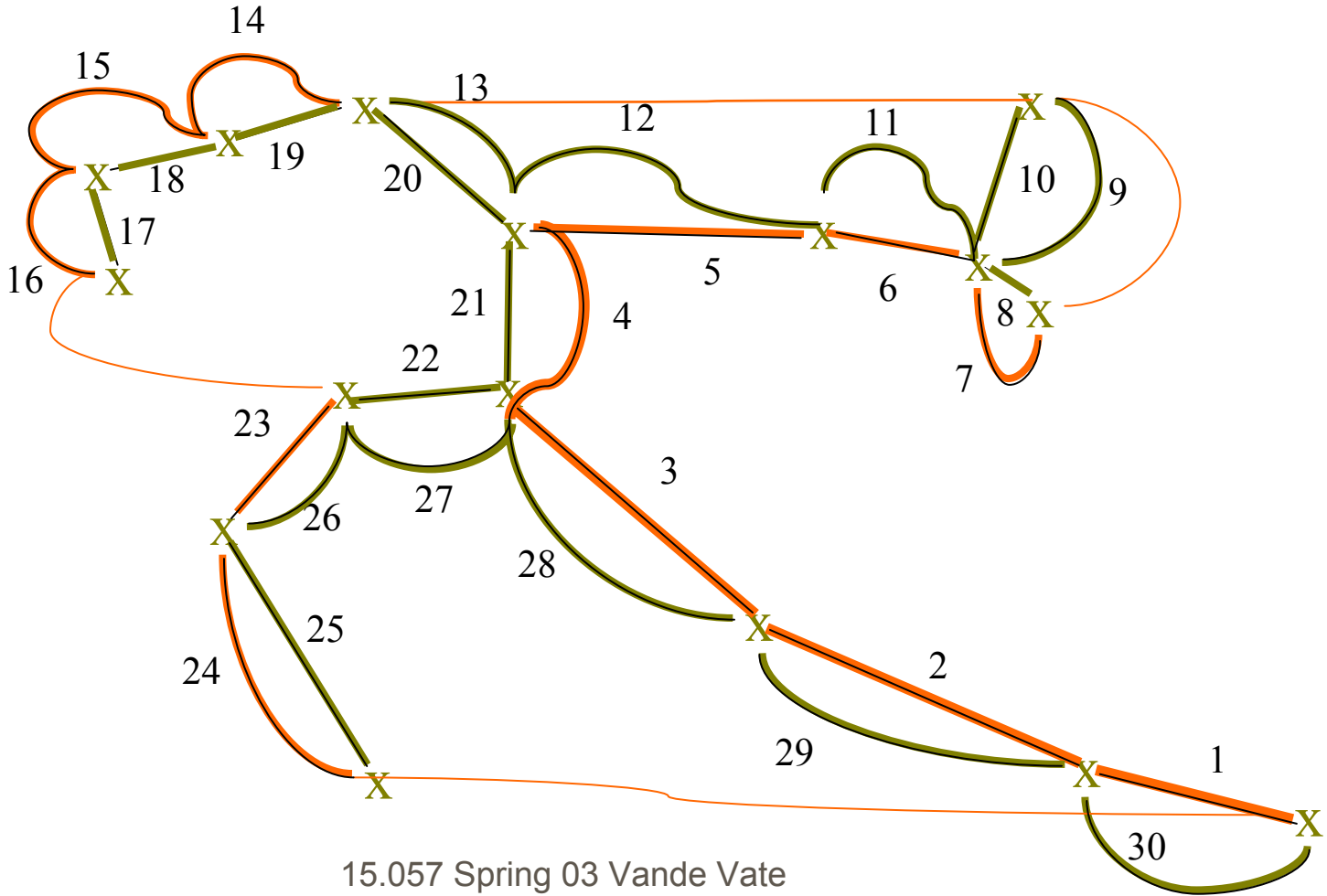
Double the Spanning Tree

- Duplicate each edge in the Spanning Tree
- The resulting graph is connected
- The degree at every node must be even
- That's an Eulerian Graph (you can start at a city, walk on each edge exactly once and return to where you started)
- It's no more than twice the length of the shortest TSP

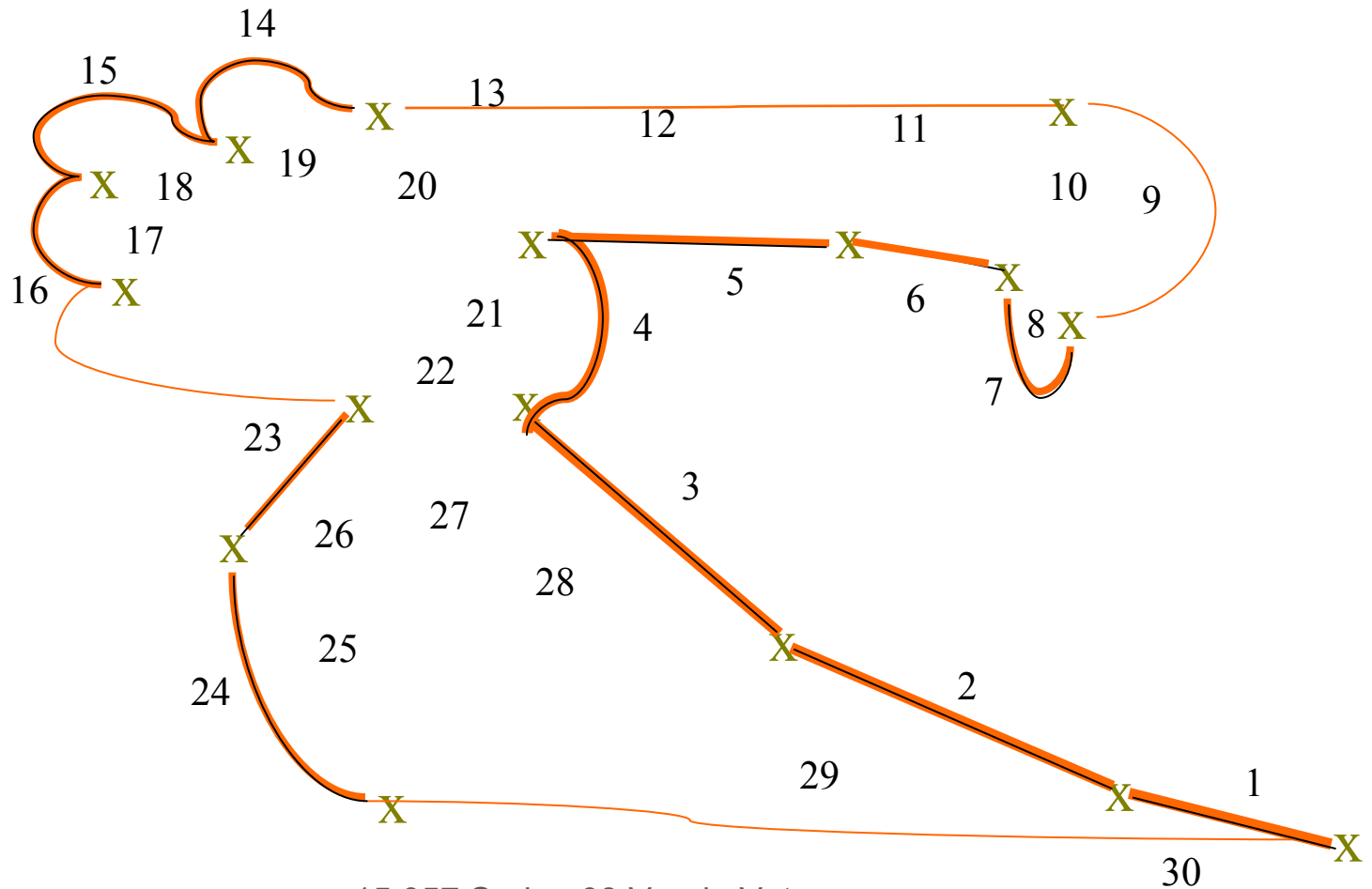
The Doubled Spanning Tree



Short Cut the Eulerian Tour



TSP Tour: at most 2X Optimum



A Novel Engineering Guarantee



- Aircraft loading
- Pallets of different weights and lengths to place on an aircraft.
- Want the center of gravity close to a specified point

The Heuristic



- Heuristic: Place the least dense remaining pallet as far as possible to the target center of gravity given the positions of the previously positioned pallets
- A sort of tempered greed

The Guarantee



- **Guarantee:** Produces a loading whose center of gravity is either within $\frac{1}{2}$ the length of the longest pallet from the target, or barring that, as close as possible
- **Comment:** Even an optimal procedure can't give a stronger guarantee across all loads.

Design Implications



- Design for the worst load
- Theoretically, the worst load is $\frac{1}{2}$ the longest pallet
- This approach guarantees it.
- The value of performance guarantees

Eliminate the Problem

- Toyota Sewing Systems
- Balancing workload in low-tech assembly systems
- Partitioning picking area among workers to improve throughput
- Optimization-based approach:
 - ▶ Detailed work-content models for each pick
 - ▶ Optimization routines to balance work-content among pickers
 - ▶ Team of Engineers to maintain the system

Self-Organizing Approach

- Bucket Brigades
- Water source = new work
- Fire = completed work
- Put fastest workers closer to the fire
- Run a bucket brigade
 - ▶ When a worker completes a job he walks back and takes over the next job from the previous worker...

Performance



■ Self Balancing

- ▶ The workers converge to a partition of the work that balances the workload
- ▶ Performance generally improves significantly

■ Bucket Brigades

- http://www.isye.gatech.edu/people/faculty/John_Bartholdi/bucket-brigades.html

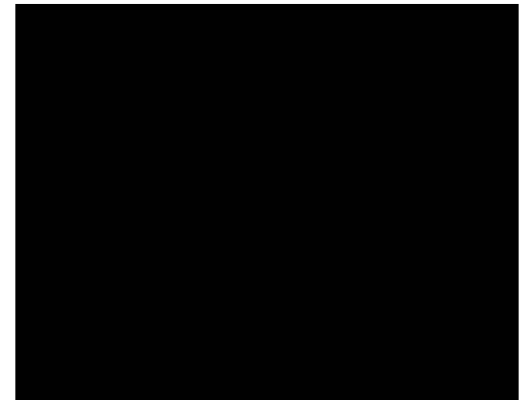
Something completely different



- Problems in low dimensions are relatively easy.
- Reduce problems in high dimensions to problems in low dimensions
- Based on fascinating mathematics

Spacefilling Curves

- There are no more points in the unit square than in the interval from 0 to 1!?



“Heuristic” Argument

- Each point (X,Y) on the map
 - ▶ $X = .165$
 - ▶ $Y = .975$
 - ▶ Space Filling Number - interleave digits
 - ▶ $\theta(X,Y) = .\mathbf{196755}$

So,...



■ Each pair of points

$$X = .165$$

$$Y = .975$$

maps to a unique point

$$\theta(X, Y) = .196755$$

How to Use this



- Map a higher dimensional problem down to a lower dimensional problem
- Solve the problem in the lower dimensional space
- Interpret the answer in the higher dimensional space

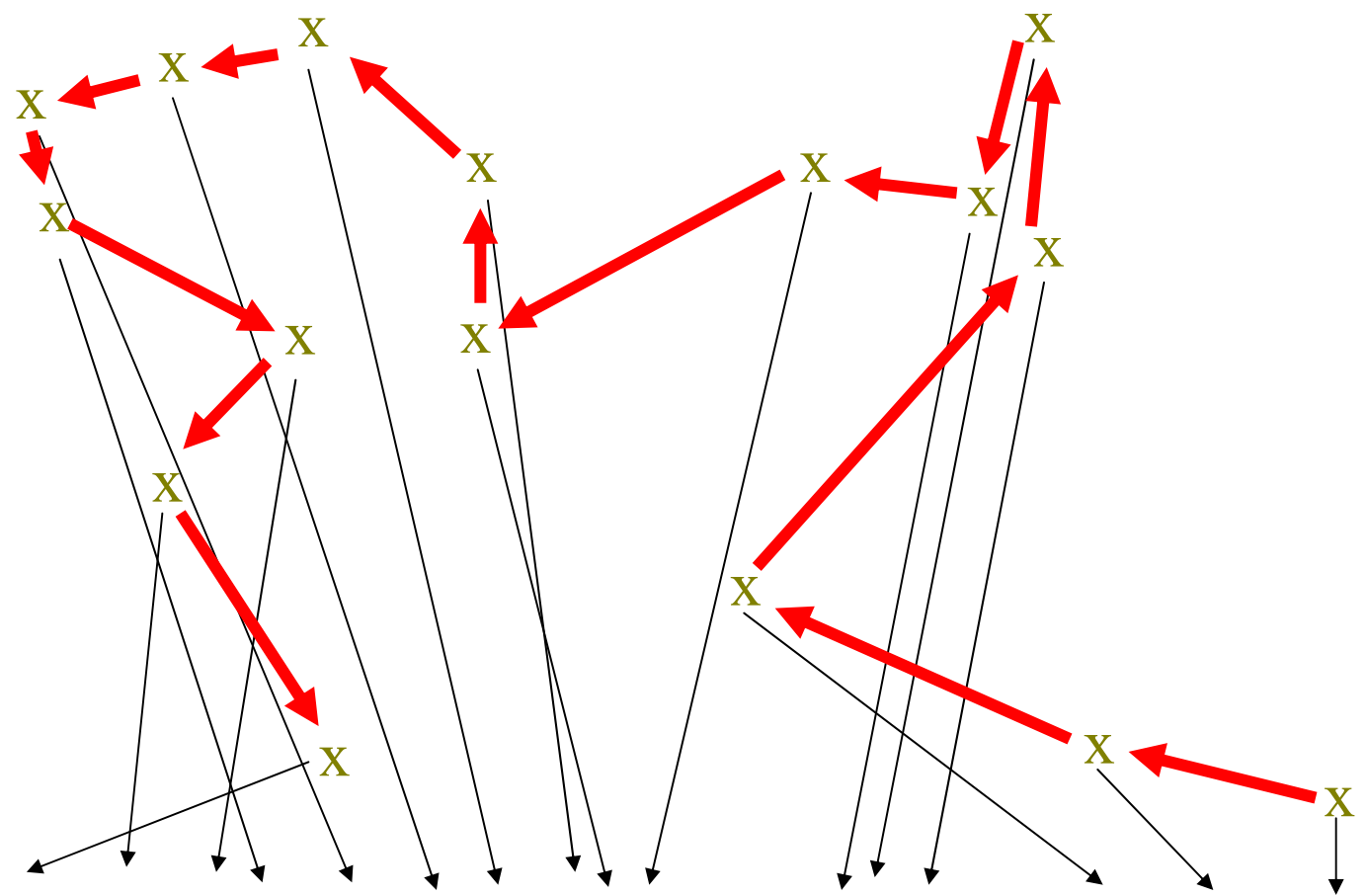
TSP Example

- A mapping of $\theta(X,Y)$ into the unit interval
- Think of this as the inverse mapping of the unit interval onto the square (our super tour)
- For a given customer $\theta(X,Y)$ is the fraction of the way along the super tour where it lies
- Visit the customers in the order of $\theta(X,Y)$ (short cut the super tour to visit our customers)



2D

1D



Applications

■ Used to

- ▶ route a plotter pen
- ▶ index data for GIS
- ▶ route drivers for Meals-On-Wheels

■ target warheads with missile defense systems

■ More on SpaceFilling Curves

<http://www.isye.gatech.edu/~jjb/>