18.335 Midterm Exam Solutions: Spring 2019

Problem 1: (10 points)

The general formula for $\kappa(A)$, from the book, is the supremum of the condition number $||A|| \cdot ||x|| / ||Ax||$ for all *x*, i.e.

$$\kappa(A) = \|A\| \left(\sup_{x \neq 0} \frac{\|x\|}{\|Ax\|} \right) = \left(\sup_{x \neq 0} \frac{\|Ax\|}{\|x\|} \right) \left(\sup_{x \neq 0} \frac{\|x\|}{\|Ax\|} \right).$$

Since *A* (*n* columns) is a subset of the columns of $B(n' \ge n$ columns), then for every $x \in \mathbb{C}^n$ there is an $x' \in \mathbb{C}^{n'}$ such that Ax = Bx' — that is, x' is simply *x* padded with zeros for the extra columns of *B*. Furthermore, in any of our L_p norms we have ||x|| = ||x'||. So, if x_* is a vector where $\frac{||Ax||}{||x||}$ achieves its supremum, there is an x' such that $\frac{||Ax_*||}{||x_*||} = \frac{||Bx'||}{||x'||}$, and hence

$$\sup_{x \neq 0} \frac{\|Ax\|}{\|x\|} \le \sup_{x' \neq 0} \frac{\|Bx'\|}{\|x'\|}.$$

Similarly for $\frac{\|x\|}{\|Ax\|}$. Hence $\kappa(A) \leq \kappa(B)$.

Problem 2: (5+5 points)

For a diagonalizable $m \times m$ matrix $A = X\Lambda X^{-1}$, the matrix square root is

$$A^{\frac{1}{2}} = X\Lambda^{\frac{1}{2}}X^{-1} = X \begin{pmatrix} \sqrt{\lambda_1} & & \\ & \sqrt{\lambda_2} & & \\ & & \ddots & \\ & & & \sqrt{\lambda_m} \end{pmatrix} X^{-1}.$$

- (a) A may be nearly defective, in which case X is badly conditioned and multiplying by X or X^{-1} will be inaccurate. (Being exactly defective is exceedingly rare a set of measure zero among all matrices, so you might ignore this case, but you can't ignore the possibility of being nearly defective.)
- (b) One possible answer is that all her matrices are Hermitian (or anti-Hermitian/skew-Hermitian).

For the $X\Lambda^{\frac{1}{2}}X^{-1}$ formula to be accurate, you need *X* to be well-conditioned, and the best case for this is if *A* is normal ($AA^* = A^*A$), in which case *X* can be chosen unitary (condition number 1). The only cases where you can typically see that *A* is normal by inspection are the Hermitian or anti-Hermitian cases. (Another possibility would be diagonal matrices *A*, but you were told that the matrices were non-sparse.)

Problem 3: (10 points)

If one of the x_i values is sufficiently large and positive (\gtrsim 710 in double precision), then e^{x_i} will overflow and you will get +Inf. Alternatively, if *all* of the x_i values are sufficiently large in magnitude and negative (\lesssim -745 in double precision), then e^{x_i} will underflow to +0.0 and the log will give you –Inf. To start with, we want to avoid both of these cases.

8/10 points: A simple solution is to compute $X = \max_i x_i$, and then use the identity

$$f(x) = \log\left(\sum_{i=1}^{n} e^{x_i}\right) = \log\left(e^X \sum_{i=1}^{n} e^{x_i - X}\right) = X + \log\left(\sum_{i=1}^{n} e^{x_i - X}\right).$$

This solves the overflow problem, because $x_i - X \le 0$ and hence $e^{x_i - X}$ can only be small, not large. What about underflow? Without loss of generality, let's suppose that $X = x_1$. Then we have

$$f(x) = X + \log\left(1 + \sum_{i=2}^{n} e^{x_i - X}\right).$$

Notice that e^{x_i-X} in the sum may underflow to zero, but we will never get zero as the argument of the log because we have $1 + \cdots \ge 1$. So we won't get –Inf even if the x_i are large negative numbers.

10/10 points: However, there is still a subtle problem: if $\sum_{i=2}^{n} e^{x_i - X} \ll 1$, then in floating-point arithmetic we may get

$$X + \log\left(1 \oplus \sum_{i=2}^{n} e^{x_i - X}\right) = X + \log(1) = X,$$

so the contribution of the $\sum_{i=2}^{n} e^{x_i - X}$ is lost. Recall the Taylor expansion

$$\log(1+y) = y - \frac{y^2}{2} + \frac{y^3}{3} - \cdots,$$

so even if $0 < y \ll 1$, we are not supposed to get zero from the log. This can lead to an inaccurate result. For example, consider the case of n = 2 with $x_1 = 10^{-20} > x_2 = \log 10^{-20} \approx -46.0517$. Then the correct answer is

$$x_1 + \log(1 + e^{x_2}) = 10^{-20} + \log(1 + 10^{-20}) \approx 2 \times 10^{-20}$$

but in floating-point arithmetic we will get $x_1 \oplus \log(1 \oplus e^{x_2}) = x_1 \oplus \log(1) = x_1 \approx 10^{-20}$, which is off by a factor of 2! The solution is that we need to compute $\log 1p(y) = \log(1+y)$ accurately even for very small y, and fortunately most math libraries (including Julia's) provide a built-in "log1p" function that does just that. So, in summary, if we want an accurate result we really need to use a floating-point version of the expression:

$$f(x) = X + \log \ln(\sum' e^{x_i - X}),$$

where \sum' denotes the sum omitting a single term with $x_i = X = \max_j x_j$. If we want, we could implement this sum with pairwise summation or similar, for even more accuracy. If we didn't have a "log1p" function available, to accurately compute $\log_1 p(y) = \log_1(1+y)$, we could implement it ourselves using the Taylor series when |y| is sufficiently small (although it turns out that there are more clever ways to do it).

Problem 4: (10 points)

8/10: We can use the Hessenberg factorization $A = QHQ^*$, which can be computed in $\Theta(m^3)$ operations from class, and for which *H* is tridiagonal if *A* is Hermitian. Then

$$f(z) = \det(A - zI) = \det(QHQ^* - zI) = \det[Q(H - zI)Q^*] = \det(Q)\det(H - zI)\det(Q^*) = \det(H - zI)$$

by elementary properties of determinants. Since H - zI is tridiagonal, as mentioned in class we can find its LU factorization in $\Theta(m)$ operations, from which the determinant is simply the product of the diagonal entries of U. A little care is needed for the case where H - zI is nearly singular, though.

10/10: Since in neither the book nor in class did we explicitly study the LU decomposition of tridiagonal matrices — I only stated in passing that it was $\Theta(m)$ — and some care is needed in the singular case, to get full marks on this problem you need to do a bit more work to convince me of how you would compute det *H*. In particular, there are lots of ways to derive nice explicit formulas here. (Outside of an exam you would just google "determinant tridiagonal matrix," of course.) For example, if we write:

$$H = \begin{pmatrix} a_1 & b_1 & & & \\ b_1 & a_2 & \overline{b_2} & & & \\ & b_2 & a_3 & \ddots & & \\ & & \ddots & \ddots & \overline{b_{m-1}} \\ & & & b_{m-1} & a_m \end{pmatrix},$$

then each step of Gaussian elimination transforms the 2×2 diagonal block

$$\left(\begin{array}{cc} d_k & \overline{b_k} \\ b_k & a_{k+1} \end{array}\right) \longrightarrow \left(\begin{array}{cc} d_k & \overline{b_k} \\ 0 & a_{k+1} - \frac{b_k \overline{b_k}}{d_k} \end{array}\right)$$

so that the diagonal entries satisfy the recurrence relation

$$d_1 = a_1$$

 $d_{k+1} = a_{k+1} - \frac{|b_k|^2}{d_k}.$

and once it is reduced to upper-triangular form then the determinant is simply the product of the pivots $\prod d_k$. This recurrence may look slightly dangerous at first — what if $d_k = 0$? However, this division by zero goes away when you multiply the entries together — consider the term $d_k d_{k+1}$ — and after a little thought you can see that the the product

$$p_k = \prod_{i=1}^k d_k$$

satisfies a simpler recurrence (called the "continuant" in linear algebra):

$$p_0 = 1$$

$$p_1 = a_1$$

$$p_{k+1} = d_{k+1}p_k = p_k a_{k+1} - p_{k-1}|b_k|^2,$$

which has no possibility of division by zero, giving det $H = p_m$ in $\Theta(m)$ operations. Finally, get det(H - zI), we simply modify this recurrence to subtract *z* from the diagonals:

$$p_0 = 1$$

$$p_1 = a_1 - z$$

$$p_{k+1} = p_k(a_{k+1} - z) - p_{k-1}|b_k|^2.$$

This recurrence can also be derived in other ways, e.g. by cofactor formulas. For the case of real b_i (real-symmetric A and H), the same recurrence is given in equation (30.9) of the Trefethen & Bau textbook.

Another possible $\Theta(m)$ determinant algorithm is to do the QR factorization of H - zI, which can be accomplished in $\Theta(m)$ operations by Givens rotations as you showed in pset 3. Then the determinant is simply det *R* (since det Q = 1 for Givens rotations), which is the product of the diagonal entries of *R*.

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