

Subject 24.242. Logic II. Answers to the last homework assignment

Recall that a *normal modal system* for the modal sentential calculus is a set of formulas Γ that meets the following conditions:

- (TC) Every tautological consequence of Γ is in Γ .
- (Nec) If ϕ is in Γ , so in $\Box\phi$.
- (K) All instances of the schema $(\Box(\phi \rightarrow \psi) \rightarrow (\Box\phi \rightarrow \Box\psi))$ are in Γ .

1. A binary relation R on a set W is *symmetric* iff, for every v and w in W , if Rvw then Rvw . Let KB be the smallest normal modal system that contains all instances of the schema

$$(B) \quad (\Diamond\Box\phi \rightarrow \phi)$$

Show that a sentence is in KB if and only if it's valid for the class of frames $\langle W, R, I \rangle$, with R symmetric.

First, we show that (B) is valid for the class of symmetric frames. Suppose that R is symmetric and that $\Diamond\Box\phi$ is true at the world w in the frame $\langle W, R, I \rangle$. Then there is a world v accessible from w in which $\Box\phi$ is true. So ϕ is true in every world accessible from v . In particular, ϕ is true in w , since, by symmetry, w is accessible from v . So $(\Diamond\Box\phi \rightarrow \phi)$ is valid in $\langle W, R, I \rangle$.

Let Γ be the set of sentences valid for the class of symmetric frames. Γ is a normal modal system that includes (B), and so Γ includes KB. We need to show that, if a sentence ϕ isn't in KB, it isn't in Γ . That is, we need to show that, if ϕ isn't in KB, then there is a symmetric frame in which there is a world in which ϕ is false. We know that the canonical frame for KB contains a world in which ϕ is false; so it will be enough to show that the canonical frame for KB is symmetric.

Suppose that w and v are worlds in the canonical frame for KB and that Rvw . We need to see that Rvw , that is, we need to see that, whenever $\Box\psi$ is in v , ψ is in w . Since $\Box\psi$ is true in v , $\Diamond\Box\psi$ is true in every world that has access to v ; in particular, $\Diamond\Box\psi$ is true in w . Since $(\Diamond\Box\psi \rightarrow \psi)$ is true in w , it follows that ψ is true in w , and so $\psi \in w$.

5. Prove de Jongh's theorem that all instances of schema

$$(4) \quad (\Box\phi \rightarrow \Box\Box\phi)$$

are elements of the smallest normal modal system that includes all instances of the schema:

$$(L) \quad (\Box(\Box\phi \rightarrow \phi) \rightarrow \Box\phi).$$

[Hint: The instance of schema (L) that you'll use is $(\Box(\Box(\phi \wedge \Box\phi) \rightarrow (\phi \wedge \Box\phi)) \rightarrow \Box(\phi \wedge \Box\phi))$.]

1. $((\phi \wedge \Box\phi) \rightarrow \phi)$ (TC)
2. $\Box((\phi \wedge \Box\phi) \rightarrow \phi)$ (Nec), 1
3. $(\Box((\phi \wedge \Box\phi) \rightarrow \phi) \rightarrow (\Box(\phi \wedge \Box\phi) \rightarrow \Box\phi))$ (K)
4. $(\Box(\phi \wedge \Box\phi) \rightarrow \Box\phi)$ (TC), 2, 3
5. $(\phi \rightarrow (\Box(\phi \wedge \Box\phi) \rightarrow (\phi \wedge \Box\phi)))$ (TC), 4
6. $\Box(\phi \rightarrow (\Box(\phi \wedge \Box\phi) \rightarrow (\phi \wedge \Box\phi)))$ (Nec), 5

7. $(\Box(\phi \rightarrow (\Box(\phi \wedge \Box\phi) \rightarrow (\phi \wedge \Box\phi))) \rightarrow (\Box\phi \rightarrow \Box(\Box(\phi \wedge \Box\phi) \rightarrow (\phi \wedge \Box\phi))))$ (K)
8. $(\Box\phi \rightarrow \Box(\Box(\phi \wedge \Box\phi) \rightarrow (\phi \wedge \Box\phi)))$ (TC), 6, 7
9. $(\Box(\Box(\phi \wedge \Box\phi) \rightarrow (\phi \wedge \Box\phi)) \rightarrow \Box(\phi \wedge \Box\phi))$ (L)
10. $(\Box\phi \rightarrow \Box(\phi \wedge \Box\phi))$ (TC), 8, 9
11. $((\phi \wedge \Box\phi) \rightarrow \Box\phi)$ (TC)
12. $\Box((\phi \wedge \Box\phi) \rightarrow \Box\phi)$ (Nec), 11
13. $(\Box((\phi \wedge \Box\phi) \rightarrow \Box\phi) \rightarrow (\Box(\phi \wedge \Box\phi) \rightarrow \Box\Box\phi))$ (K)
14. $(\Box(\phi \wedge \Box\phi) \rightarrow \Box\Box\phi)$ (TC), 12, 13
15. $(\Box\phi \rightarrow \Box\Box\phi)$ (TC), 10, 14