

## Gödel's Completeness Theorem

Thm 1, good news: only need to know\* a few axioms & rules, to prove *all* validities. \**Theoretically only: having more rules is convenient.* 

Axioms & Inference Rules

Namely, starting from a few propositional & simple predicate validities, every valid assertion can be proved using just UG and *modus ponens* repeatedly!



## **Cannot Determine Validity**

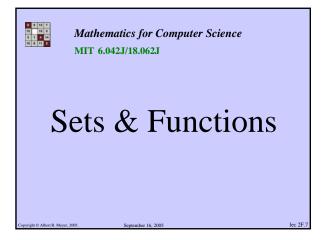
Thm 2, Bad News: there is *no procedure* that determines when quantified assertions are valid (in contrast to propositional formulas).

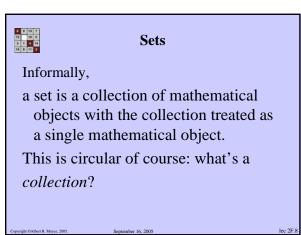
**Gödel's** *In***completeness Theorem for Arithmetic** Thm 3, **Worse News**: if we stick to domain,  $\mathbb{N}$ , with predicates x + y = z,  $x \cdot y = z$ , then **no proof system** can prove all the true assertions!

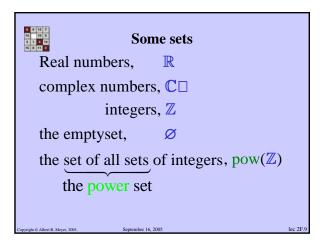


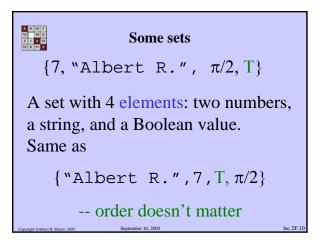
## **Three Profound Theorems**

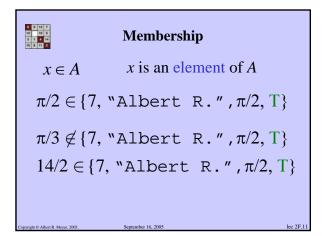
We won't prove these Theorems. Their proofs usually require half a term in an intro logic course after 6.042.

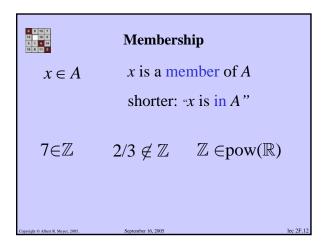


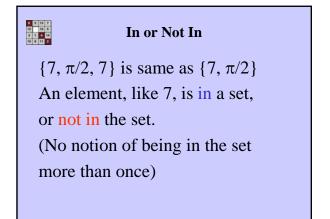


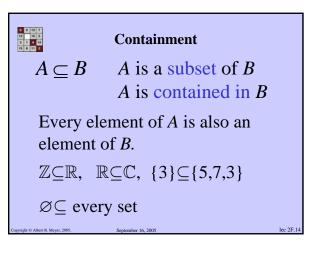


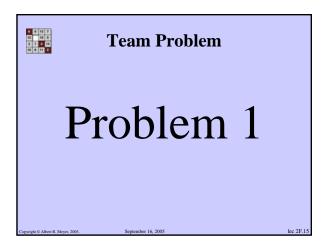


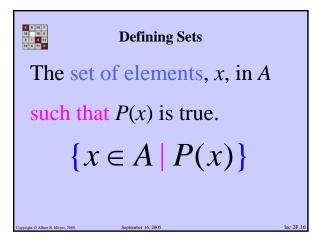


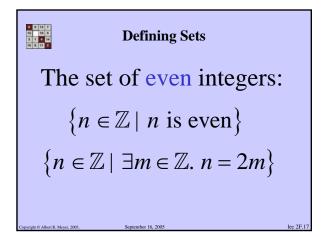


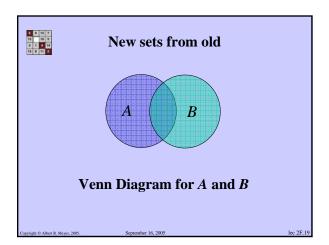


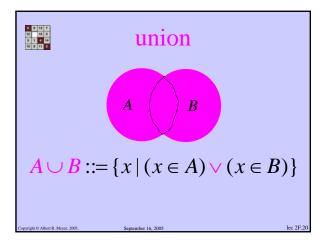


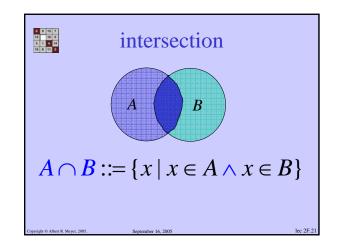


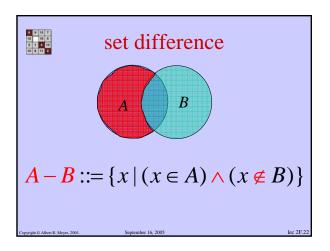


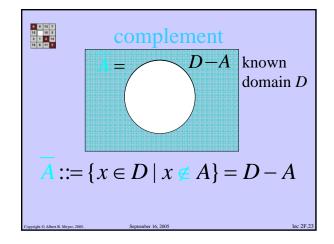


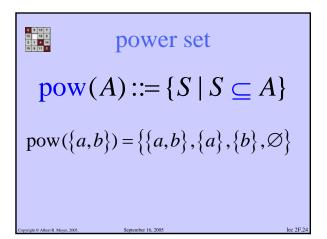


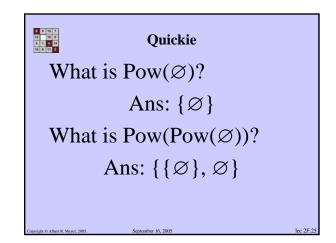




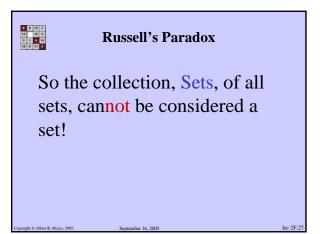


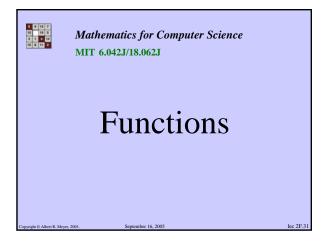






Russell's ParadoxLet 
$$W ::= \{S \in Sets \mid S \notin S\}$$
so  $S \in W \leftrightarrow S \notin S$ Let  $S$  be  $W$  and reach a  
contradiction: $W \in W \leftrightarrow W \notin W$ 





 $f: A \to B$ 

function, *f*, from set *A* to set *B* associates an element,  $f(a) \in B$ with an element  $a \in A$ . *Example: f* is the string-length function: f(``aabd'')=4

 $f: A \to B$ 

f is the string-length function. A, the domain of f, is the set of strings. B, the codomain of f, is  $\mathbb{N}$ 

$$g: \mathbb{R}^2 \to \mathbb{R}$$

$$g(x, y) = \frac{1}{x - y}$$
domain(g) = all pairs of reals
codomain(g) = all reals
But g is partial:
not defined on all pairs of reals

$$g': D \to \mathbb{R}$$

$$g'(x, y) = \frac{1}{x - y}$$
where  $D = \mathbb{R}^2 - \{(x, y) \mid y = x\}$ 
Then g' is total:
defined on all pairs in domain D.

