



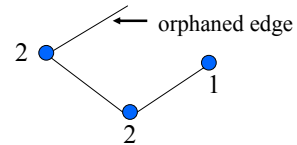
# Graphs: Handshaking, Connectivity & Trees



## Possible Graph?

Is there a graph with  
vertex degrees 2,2,1?

**NO!**



## An Impossible Graph

$$2+2+1 = \text{odd}$$

**The Handshaking Lemma:**

The sum of degrees must be even.

$$2|E| = \sum_{v \in V} \deg(v)$$



## The Handshaking Lemma

*Proof:* Each edge contributes 2  
to the sum on the right

$$2|E| = \sum_{v \in V} \deg(v)$$



## Paths & Simple Paths

Path:  $v_0, v_1, \dots, v_n$  for  $n \geq 0$

where  $v_i - v_{i+1}$  is an  
edge for all  $i$  s.t.  $0 \leq i < n$

*Simple* path: all  $v_i$  different



## Connectedness

**Vertices**  $v, w$  are **connected** iff  
there is a path starting at  $v$  and  
ending at  $w$ .

A **graph** is **connected** iff every two  
of its vertices are connected



## Paths & Simple Paths

*Lemma:* The shortest path connecting two vertices is simple!

*Simple* path: all  $v_i$  different

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## Cycles & Simple Cycles

Cycle:  $v_0, v_1, \dots, v_n, v_0$  for  $n \geq 1$

where  $v_i - v_{i+1}$  is an edge for all  $i$  s.t.  $0 \leq i < n$

and  $v_n - v_0$

*Simple* cycle: all  $v_i$  different and  $n \geq 2$

*same* cycle:  $v_2, v_1, v_0, v_n, v_{n-1}, \dots, v_3, v_2$

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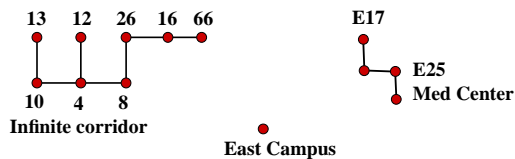
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## Connected Components

The more connected components, the more broken up the graph is.



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## Connected Components

The connected component of vertex  $v$  is

$\{w \mid v \text{ and } w \text{ are connected}\}$

Same as the block  $[v]$  of the *connectedness relation*

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## Cut Edges

An edge is a *cut edge* if removing it from the graph disconnects two vertices.

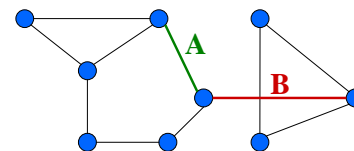
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## Cut Edges



*A* is not, *B* is

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## Cut Edges and Cycles

*Lemma:* An edge is a cut edge *iff* it is not part of a simple cycle.

*Proof:* Class problem.

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## Cut Edges

Fault-tolerant design:

In a tree, every edge is a cut edge  
(bad)

In a mesh, no edge is a cut edge  
(good)

Tradeoff edges for failure tolerance

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## $k$ -Connectedness

*Def:*  $k$ -connected *iff* remains connected when any  $k-1$  edges are deleted.

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## $k$ -Connectedness

*Example:*

$K_n$  is  $(n-1)$ -connected

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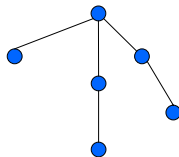
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## Trees

A *tree* is a connected graph with no cycles.



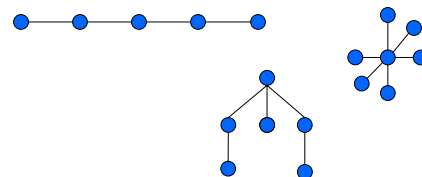
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## More Trees



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### Other Tree Definitions

*Lemma:* A tree is a connected graph with  $n$  vertices and  $n - 1$  edges.

*Lemma:* A tree is an edge-minimal connected graph on a set of vertices.

*Lemma:* A tree is a graph with a **unique** path between any 2 vertices.

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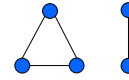
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### Be careful with these definitions

*A tree is a graph with  $n$  vertices and  $n - 1$  edges??*

**NO:**



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### Spanning Tree

*Def:* A subgraph that is a tree with all the vertices.

**Always exists:** any minimum-edge-size, connected subgraph on all the vertices

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# Problems 4 & 5

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