## Solutions to In-Class Problems Week 3, Fri.

Problem 1. Given an unlimited supply of 3 cent and 5 cent stamps, what postages are possible? Prove it using Strong Induction. Hint: Try some examples! Which postage values between 1 and 25 cents can you construct from 3 cent and 5 cent stamps?
Solution. Let's use our examples to first try to guess the answer and then try to prove it. Let's begin filling in a table that shows the values of all possible combinations of 3 and 5 cent stamps. The column heading is the number of 5 cent stamps and the row heading is the number of 3 cent stamps.

|  | 0 | 1 | 2 | 3 | 4 | 5 | $\ldots$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 5 | 10 | 15 | 20 | 25 | $\cdots$ |
| 1 | 3 | 8 | 13 | 18 | 23 | $\ldots$ |  |
| 2 | 6 | 11 | 16 | 21 | $\ldots$ |  |  |
| 3 | 9 | 14 | 19 | 24 | $\cdots$ |  |  |
| 4 | 12 | 17 | 22 | $\cdots$ |  |  |  |
| 5 | 15 | 20 | $\cdots$ |  |  |  |  |
| $\ldots$ | $\cdots$ | $\cdots$ |  |  |  |  |  |

Looking at the table, a reasonable guess is that the possible postages are $0,3,5$, and 6 cents and every value of 8 or more cents. Let's try to prove this last part using strong induction.
Claim 1.1. For all $n \geq 8$, it is possible to produce $n$ cents of postage from $3 \nmid$ and $5 \nmid$ stamps.
Now let's preview the proof. The induction hypothesis will be

$$
\begin{equation*}
P(n) \quad::=\quad \text { if } n \geq 8, \text { then } n \not \subset \text { postage can be produced using } 3 \not \subset \text { and } 5 \not \subset \text { stamps } \tag{1}
\end{equation*}
$$

A proof by strong induction will have the same five-part structure as an ordinary induction proof. The base case, $P(0)$, won't be interesting because $P(n)$ is vacuously true for all $n<8$.

In the inductive step we have to show how to produce $n+1$ cents of postage, assuming the strong induction hypothesis that we know how to produce $k \notin$ of postage for all values of $k$ between 8 and $n$. A simple way to do this is to let $k=n-2$ and produce $k \not \subset$ of postage; then add a $3 \notin$ stamp to get $n+1$ cents.
But we have to be careful; there is a pitfall in this method. If $n+1$ is 8,9 or 10 , then we can not use the trick of creating $n+1$ cents of postage from $n-2$ cents and a 3 cent stamp. In these cases, $n-2$ is less than 8 . None of the strong induction assumptions help us make less than $8 \notin$ postage. Fortunately, making $n+1$ cents of postage in these three cases can be easily done directly.

[^0]Proof. The proof is by strong induction. The induction hypothesis, $P(n)$, is given by (1).
Base case: $n=0: P(0)$ is true vacuously.
Inductive step: In the inductive step, we assume that it is possible to produce postage worth $8,9, \ldots, n$ cents in order to prove that it is possible to produce postage worth $n+1$ cents.
There are four cases:

1. $n+1<8$ : So $P(n+1)$ holds vacuously.
2. $n+1=8: P(n+1)$ holds because we produce $8 \not \subset$ postage using one $3 \not \subset$ and one $5 \not \subset$ stamp.
3. $n+1=9$ : $P(n+1)$ holds by using three $3 \Varangle$ stamps.
4. $n+1=10: P(n+1)$ holds by using two $5 \not \subset$ stamps.
5. $n+1>10$ : We have $n \geq 10$, so $n-2 \geq 8$ and by strong induction we may assume we can produce exactly $n-2$ cents of postage. With an additional $3 \phi$ stamp we can therefore produce $n+1$ cents of postage.

So in every case, $P(0) \wedge P(1) \wedge \ldots P(n) \longrightarrow P(n+1)$. By strong induction, we have concluded that $P(n)$ is true for all $n \in \mathbb{N}$.

Problem 2. Use the Well-ordering Principle to prove that there is no solution over the positive integers to the equation:

$$
4 a^{3}+2 b^{3}=c^{3} .
$$

Solution. We use contradiction and the well-ordering principle. Let $S$ be the set of all positive integers, $a$, such that there exist positive integers, $b$, and, $c$, that satisfy the equation.
Assume for the purpose of obtaining a contradiction that $S$ is nonempty. Then $S$ contains a smallest element, $a_{0}$, by the well-ordering principle. By the definition of $S$, there exist corresponding positive integers, $b_{0}$, and, $c_{0}$, such that:

$$
4 a_{0}^{3}+2 b_{0}^{3}=c_{0}^{3}
$$

The left side of this equation is even, so $c_{0}^{3}$ is even, and therefore $c_{0}$ is also even. Thus, there exists an integer, $c_{1}$, such that $c_{0}=2 c_{1}$. Substituting into the preceding equation and then dividing both sides by 2 gives:

$$
2 a_{0}^{3}+b_{0}^{3}=4 c_{1}^{3}
$$

Now $b_{0}^{3}$ must be even, so $b_{0}$ is even. Thus, there exists an integer, $b_{1}$, such that $b_{0}=2 b_{1}$. Substituting into the preceding equation and dividing both sides by 2 again gives:

$$
a_{0}^{3}+4 b_{1}^{3}=2 c_{1}^{3}
$$

From this equation, we know that $a_{0}^{3}$ is even, so $a_{0}$ is also even. Thus, there exists an integer, $a_{1}$, such that $a_{0}=2 a_{1}$. Substituting into the previous equation one last time and dividing by 2 one last time gives:

$$
4 a_{1}^{3}+2 b_{1}^{3}=c_{1}^{3}
$$

Evidently, $a=a_{1}, b=b_{1}$, and $c=c_{1}$ is another solution to the original equation, and so $a_{1}$ is an element of $S$. But this is a contradiction, because $a_{1}<a_{0}$ and $a_{0}$ was defined to be the smallest element of $S$. Therefore, our assumption was wrong, and the original equation has no solutions over the positive integers.
This argument is quite similar to the proof that $\sqrt{2}$ is irrational. In fact, looking back, we implicitly relied on the Well-ordering Principle in that proof when we claimed that a rational number could be written as a fraction in lowest terms. We've been using the Well-ordering Principle on the sly from early on!


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