

## Homework 5 additional problems

1. *Heuristic suboptimal solution for Boolean LP.* This exercise builds on exercises 4.15 and 5.13 in *Convex Optimization*, which involve the Boolean LP

$$\begin{aligned} & \text{minimize} && c^T x \\ & \text{subject to} && Ax \preceq b \\ & && x_i \in \{0, 1\}, \quad i = 1, \dots, n, \end{aligned}$$

with optimal value  $p^*$ . Let  $x^{\text{rlx}}$  be a solution of the LP relaxation

$$\begin{aligned} & \text{minimize} && c^T x \\ & \text{subject to} && Ax \preceq b \\ & && 0 \preceq x \preceq \mathbf{1}, \end{aligned}$$

so  $L = c^T x^{\text{rlx}}$  is a lower bound on  $p^*$ . The relaxed solution  $x^{\text{rlx}}$  can also be used to guess a Boolean point  $\hat{x}$ , by rounding its entries, based on a threshold  $t \in [0, 1]$ :

$$\hat{x}_i = \begin{cases} 1 & x_i^{\text{rlx}} \geq t \\ 0 & \text{otherwise,} \end{cases}$$

for  $i = 1, \dots, n$ . Evidently  $\hat{x}$  is Boolean (*i.e.*, has entries in  $\{0, 1\}$ ). If it is feasible for the Boolean LP, *i.e.*, if  $A\hat{x} \preceq b$ , then it can be considered a guess at a good, if not optimal, point for the Boolean LP. Its objective value,  $U = c^T \hat{x}$ , is an upper bound on  $p^*$ . If  $U$  and  $L$  are close, then  $\hat{x}$  is nearly optimal; specifically,  $\hat{x}$  cannot be more than  $(U - L)$ -suboptimal for the Boolean LP.

This rounding need not work; indeed, it can happen that for all threshold values,  $\hat{x}$  is infeasible. But for some problem instances, it can work well.

Of course, there are many variations on this simple scheme for (possibly) constructing a feasible, good point from  $x^{\text{rlx}}$ .

Finally, we get to the problem. Generate problem data using

```
rand('state',0);
n=100;
m=300;
A=rand(m,n);
b=A*ones(n,1)/2;
c=-rand(n,1);
```

You can think of  $x_i$  as a job we either accept or decline, and  $-c_i$  as the (positive) revenue we generate if we accept job  $i$ . We can think of  $Ax \preceq b$  as a set of limits on

$m$  resources.  $A_{ij}$ , which is positive, is the amount of resource  $i$  consumed if we accept job  $j$ ;  $b_i$ , which is positive, is the amount of resource  $i$  available.

Find a solution of the relaxed LP and examine its entries. Note the associated lower bound  $L$ . Carry out threshold rounding for (say) 100 values of  $t$ , uniformly spaced over  $[0, 1]$ . For each value of  $t$ , note the objective value  $c^T \hat{x}$  and the maximum constraint violation  $\max_i (A\hat{x} - b)_i$ . Plot the objective value and the maximum violation versus  $t$ . Be sure to indicate on the plot the values of  $t$  for which  $\hat{x}$  is feasible, and those for which it is not.

Find a value of  $t$  for which  $\hat{x}$  is feasible, and gives minimum objective value, and note the associated upper bound  $U$ . Give the gap  $U - L$  between the upper bound on  $p^*$  and the lower bound on  $p^*$ . If you define vectors `obj` and `maxviol`, you can find the upper bound as `U=min(obj(find(maxviol<=0)))`.

2. *Three measures of the spread of a group of numbers.* For  $x \in \mathbf{R}^n$ , we define three functions that measure the spread or width of the set of its elements (or coefficients). The first function is the *spread*, defined as

$$\phi_{\text{sprd}}(x) = \max_{i=1,\dots,n} x_i - \min_{i=1,\dots,n} x_i.$$

This is the width of the smallest interval that contains all the elements of  $x$ .

The second function is the *standard deviation*, defined as

$$\phi_{\text{stdev}}(x) = \left( \frac{1}{n} \sum_{i=1}^n x_i^2 - \left( \frac{1}{n} \sum_{i=1}^n x_i \right)^2 \right)^{1/2}.$$

This is the statistical standard deviation of a random variable that takes the values  $x_1, \dots, x_n$ , each with probability  $1/n$ .

The third function is the average absolute deviation from the median of the values:

$$\phi_{\text{aamd}}(x) = (1/n) \sum_{i=1}^n |x_i - \text{med}(x)|,$$

where  $\text{med}(x)$  denotes the median of the components of  $x$ , defined as follows. If  $n = 2k - 1$  is odd, then the median is defined as the value of middle entry when the components are sorted, *i.e.*,  $\text{med}(x) = x_{[k]}$ , the  $k$ th largest element among the values  $x_1, \dots, x_n$ . If  $n = 2k$  is even, we define the median as the average of the two middle values, *i.e.*,  $\text{med}(x) = (x_{[k]} + x_{[k+1]})/2$ .

Each of these functions measures the spread of the values of the entries of  $x$ ; for example, each function is zero if and only if all components of  $x$  are equal, and each function is unaffected if a constant is added to each component of  $x$ .

Which of these three functions is convex? For each one, either show that it is convex, or give a counterexample showing it is not convex. By a counterexample, we mean a specific  $x$  and  $y$  such that Jensen's inequality fails, *i.e.*,  $\phi((x+y)/2) > (\phi(x) + \phi(y))/2$ .

3. *Minimax rational fit to the exponential.* (See exercise 6.9 of *Convex Optimization*.) We consider the specific problem instance with data

$$t_i = -3 + 6(i - 1)/(k - 1), \quad y_i = e^{t_i}, \quad i = 1, \dots, k,$$

where  $k = 201$ . (In other words, the data are obtained by uniformly sampling the exponential function over the interval  $[-3, 3]$ .) Find a function of the form

$$f(t) = \frac{a_0 + a_1 t + a_2 t^2}{1 + b_1 t + b_2 t^2}$$

that minimizes  $\max_{i=1, \dots, k} |f(t_i) - y_i|$ . (We require that  $1 + b_1 t_i + b_2 t_i^2 > 0$  for  $i = 1, \dots, k$ .)

Find optimal values of  $a_0, a_1, a_2, b_1, b_2$ , and give the optimal objective value, computed to an accuracy of 0.001. Plot the data and the optimal rational function fit on the same plot. On a different plot, give the fitting error, *i.e.*,  $f(t_i) - y_i$ .

*Hint.* You can use `strcmp(cvx_status, 'Solved')`, after `cvx_end`, to check if a feasibility problem is feasible.

4. *Complex least-norm problem.* We consider the complex least  $\ell_p$ -norm problem

$$\begin{aligned} & \text{minimize} && \|x\|_p \\ & \text{subject to} && Ax = b, \end{aligned}$$

where  $A \in \mathbf{C}^{m \times n}$ ,  $b \in \mathbf{C}^m$ , and the variable is  $x \in \mathbf{C}^n$ . Here  $\|\cdot\|_p$  denotes the  $\ell_p$ -norm on  $\mathbf{C}^n$ , defined as

$$\|x\|_p = \left( \sum_{i=1}^n |x_i|^p \right)^{1/p}$$

for  $p \geq 1$ , and  $\|x\|_\infty = \max_{i=1, \dots, n} |x_i|$ . We assume  $A$  is full rank, and  $m < n$ .

- Formulate the complex least  $\ell_2$ -norm problem as a least  $\ell_2$ -norm problem with real problem data and variable. *Hint.* Use  $z = (\Re x, \Im x) \in \mathbf{R}^{2n}$  as the variable.
- Formulate the complex least  $\ell_\infty$ -norm problem as an SOCP.
- Solve a random instance of both problems with  $m = 30$  and  $n = 100$ . To generate the matrix  $A$ , you can use the Matlab command `A = randn(m,n) + i*randn(m,n)`. Similarly, use `b = randn(m,1) + i*randn(m,1)` to generate the vector  $b$ . Use the Matlab command `scatter` to plot the optimal solutions of the two problems on the complex plane, and comment (briefly) on what you observe. You can solve the problems using the CVX functions `norm(x,2)` and `norm(x,inf)`, which are overloaded to handle complex arguments. To utilize this feature, you will need to declare variables to be `complex` in the `variable` statement. (In particular, you do not have to manually form or solve the SOCP from part (b).)

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