

A Distributed Object Model for the Java System  
Wollrath, Riggs, and Waldo, USENIX 1996

Goals:

- transparent RPC for object methods
- avoid NFS-style explicit object handles
- automatic association of relevant server w/ object ref
- allow passing of object references as arguments
  - not just to object's home server (as in NFS)
  - even to other client hosts
- distributed GC, needed for remote refs

Situations in which one client might pass remote object ref to another?  
lots of modules: shopping cart, item db, checkout, front end

Are there other approaches?

- distributed shared memory, would allow direct access to object data
- move the object to caller

first a simple call/return

- `o = ???;`
- `o.fn("hello");`
- which server to send to?
- what object on server?
- what about "hello"?
- what does RPC message contain?
- how does RMI s/w on server gain control? thread...
- how does server find the real object?
- where does server-side dispatch fn come from?

what does a stub object look like?

- type?
- contents?
- where did it come from?

is there anything special about the server-side "real" object?

is "hello" sent as a remote object ref?

how about passing an object as an argument?

- `o1 = ???;`
- `o2 = ???;`
- `o1.fn(o2);`
- what must o2 look like in the RPC message?
  - server host, object ID
- what if o1's server already knows about o2?
  - must have a table mapping object ID to ptr to o2
- what if o1's server does not know about o2?
  - where does it get stub type, implementation?
  - can stub stuff be generated purely by client?

there are probably type IDs, so client can re-use stub code

- an object ID must contain type ID, or an RPC to fetch it
- clients and servers must have tables mapping type ID to stub code

when can a server free an object?

- only when no client has a live reference

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server must somehow learn when new client gets a reference  
and when client local ref count drops to zero  
so clients must send RPCs to server to note first/last knowledge  
what if C1 passes to C2, C1 sends de-ref RPC before C2 sends ref?

what if a client crashes?

will server ever be able to free the object?