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Agents

- An agent is an autonomous program.
- It executes code and can communicate with other agents.
- All the components in a pervasive computing application (whatever that is) usually called agents
- An agent may be a “proxy” for a device
- Devices, like camera or keyboards, are controlled by some proxy agent
- Agents may appear or disappear at any time
- There is some issue in how to start them
- There can be problems when they crash
- there may be replicates



A collection of agents

- ▶ Parallel or distributed programming
 - ▶ a bunch of communicating agents working to solve a problem
 - ▶ faster
 - ▶ two heads better than one
 - ▶ geographically distributed
 - ▶ everyone can't live together

Agent communication

- ★ Two main choices:
 - ★ (which was best used to be “religious battle”)
- ★ Shared memory (SM)
 - ★ agents load and store values
 - ★ start with a set of numbers
 - ★ remove two numbers, insert their sum
 - ★ done when only one value remains
 - ★ issues: synchronization, locks, etc.
- ★ Message-passing (MP)



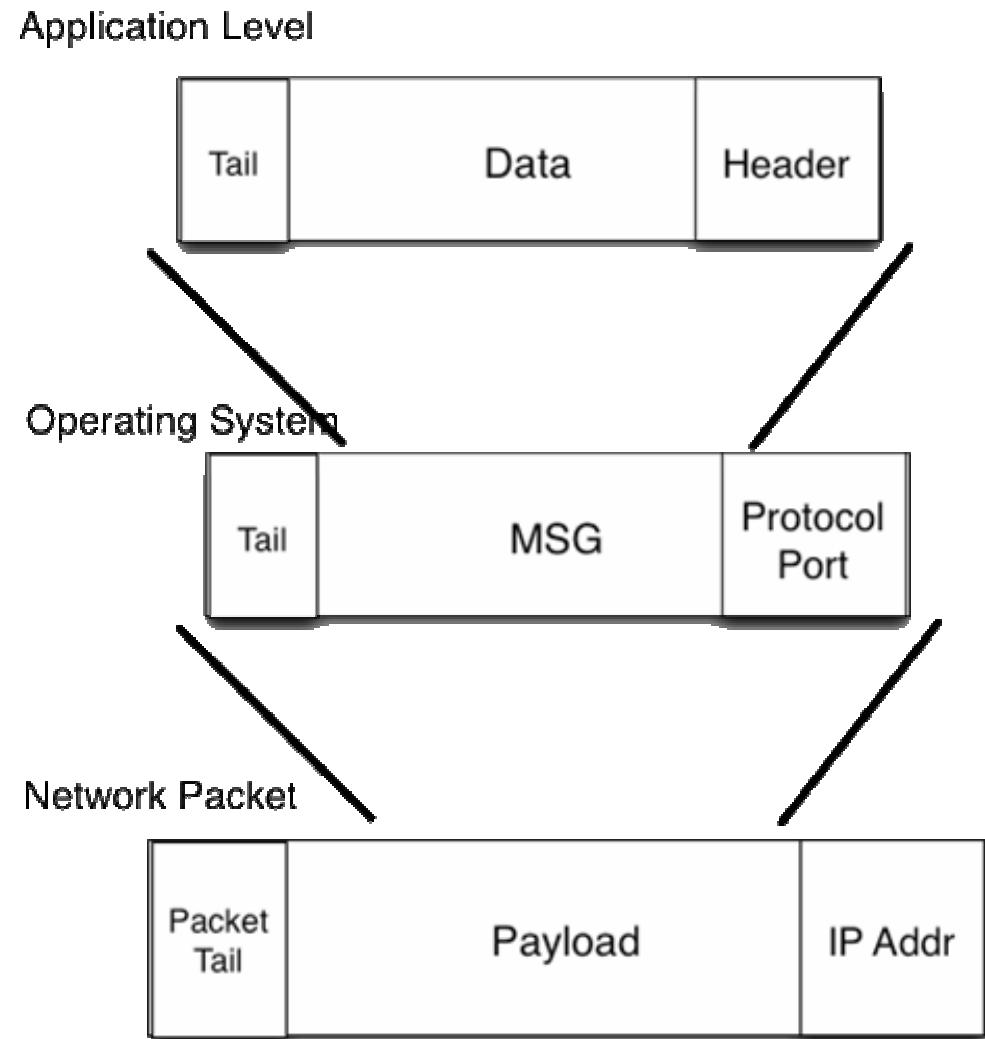
Agent communication

- Message-passing
 - two parts: destination, data
 - Agent Bob: Send(Alice, “Do you want to go out?”)
 - Agent Alice: Recv(from,msg)
 - from = Bob; msg = “do you want to go out?”
 - send(Bob, “No”)
- Issues:
 - Sender must know destination, recv need not
 - blocking or non-blocking
 - low performance, lots of copying of data
 - Note: MP can implement SM and vice-versa
 - MP on clusters, SM on multiprocessors



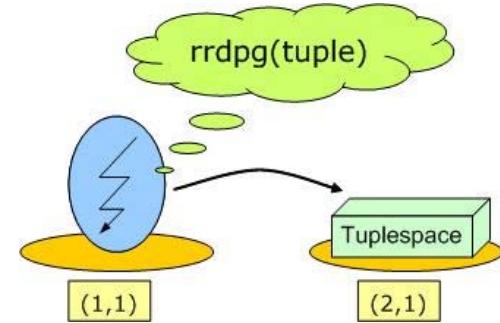
Message Passing via Sockets

- ▶ Sockets are general
 - ▶ Application can specify
 - ▶ port
 - ▶ protocol
 - ▶ other attributes
 - ▶ Message-Passing
 - ▶ library does all the specification
 - ▶ may reformat data

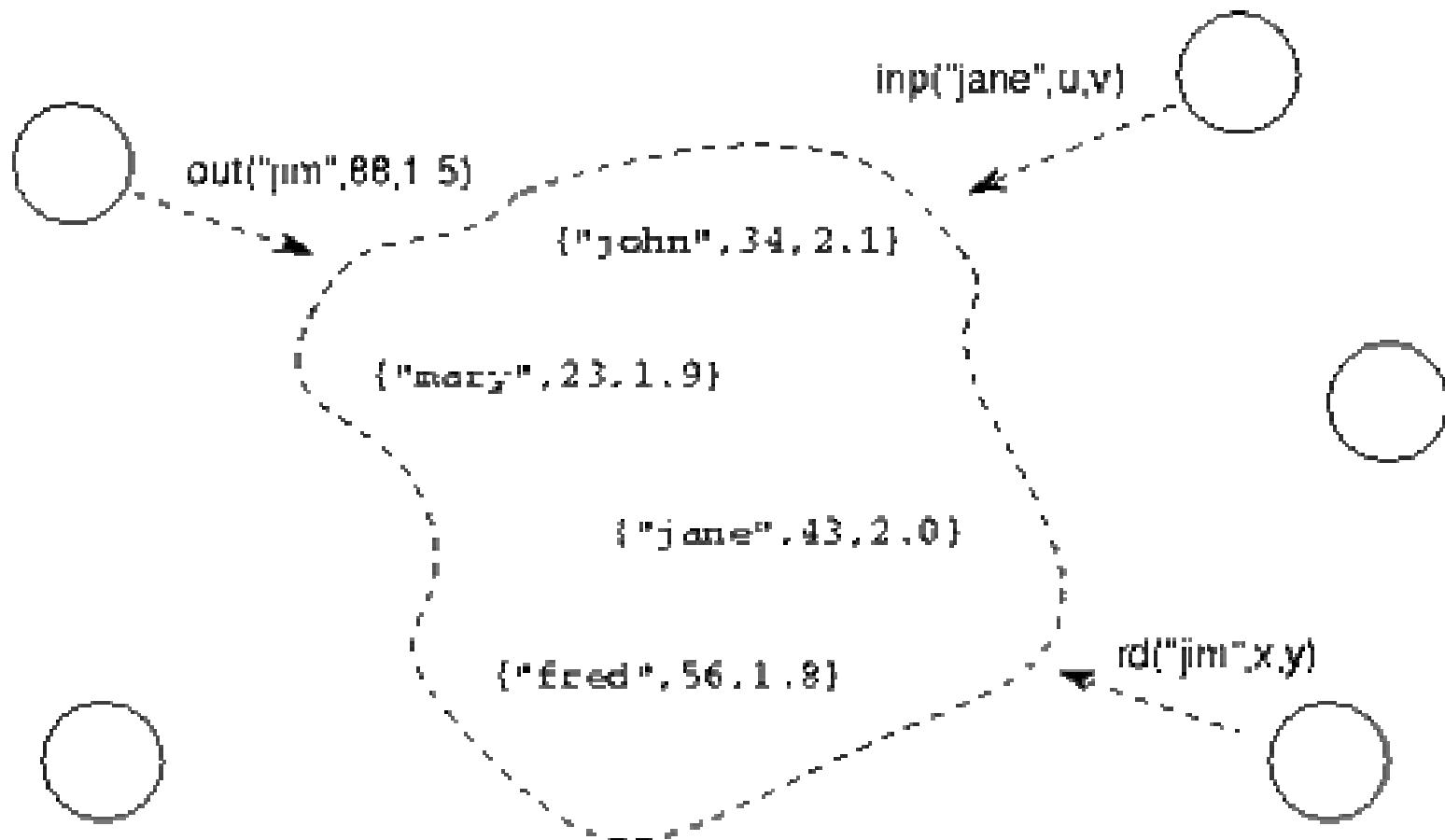


Tuple-space

- ★ A third communication mechanism!
 - ★ formed basis of Linda programming language
 - ★ tuple: ordered collection of typed elements
- ★ Basic Operations
 - ★ **out**: inserts a tuple, whose fields are either
 - ★ **actual**: a static value
 - ★ **formal**: a program variable
 - ★ **in**: extracts tuple, argument is template to match
 - ★ actuals match fields of equal type and value
 - ★ formals match fields of same type
 - ★ **rd**: same as in, but does not remove matched tuple



Tuple-space example



Linda programming example

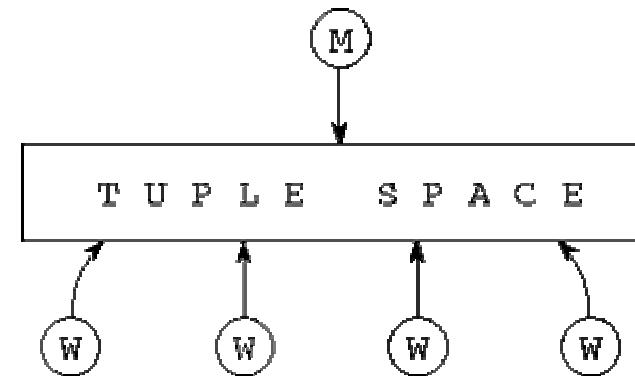
```
procedure manager
begin
  count = 0
  until end-of-file do
    read datum from file
    OUT("datum",datum)
    count = count+1
  enddo
  best = 0.0
  for i = 1 to count
    IN("score",value)
    if value > best then best = value
  endfor
  for i = 1 to numworkers
    OUT("datum","stop")
  endfor
end
```

```
procedure worker
begin
  IN("datum",datum)
  until datum = "stop" do
    value = compare(datum,target)
    OUT("score",value)
    IN("datum",datum)
  enddo
end
```



What is the big deal?

- Virtual shared memory
- tuples with [address,value]
- stores are inserts, loads are non-destructive reads
- Virtual message passing
- tuples with [dest, data]
- recv are destructive reads
- Even more, when matching on multiple fields
- Allows many types of implementations



Agent Interaction Choices

- ▶ Direct communication model
 - ▶ Jini
 - ▶ FIPA
- ▶ Indirect, Shared Data-space models
 - ▶ EventHeap (centralized)
 - ▶ MARS (fully distributed)
- ▶ Event-based publish/subscribe models
 - ▶ Siena
 - ▶ Jini Distributed Events
 - ▶ Selective subscription

Stanford's Event Heap

- ★ Based on Tuple Space paradigm
 - ★ tuple: arbitrary mix of typed fields
 - ★ mechanism for passing data & events
- ★ Extensions make it useful for agents
 - ★ many projects exist based on different extensions



Event Heap Extensions

- Extended Delivery Semantics:
 - Per-source ordering, always see events in order they are generated by the source
 - Total order: if tuple space is centralized, get this even if multiple sources
- Persistent Queries:
 - non-destructive read of those matching
 - also matches tuples inserted in future
- Event Notification:
 - like PQ, get notified of future matches
 - at most once semantics



Need more than
simple event heap

Suggested additions

- ★ Need “distributed, replicated or federated local instances
 - ★ (from paper by Storz, Friday, & Davies)
- ★ Multiple event heap instances -- but not easy of implement
- ★ View: processes that share a view have consistent ordering
- ★ Session identifiers
 - ★ non-destructive operation on per-session identifier basis
 - ★ can share, copy, or destroy id's for different semantics



More general issues

- ★ Lots and lots of middleware systems
 - ★ no winner (may never happen)
- ★ What gets communicated?
 - ★ services, events, XML records
- ★ The shared space is often a: BROKER
 - ★ The broker stores the tuples and does the matching



Big Issues

- Naming
 - This is a big, big deal.
 - e.g. how do you name a camera:
 - model brand, IP, DNS name, location, virtual space
 - via attributes (color, 740x1024), ownership?
 - Is there only one name for the agent?
- Matching
 - A big deal
 - Which attributes explicit, which implicit
 - Where to do the lookup?



Issues

- ▶ Addition information provided by broker
 - ▶ for services: how to interface them
 - ▶ filtering events
 - ▶ higher level events implemented at broker
 - ▶ based on multiple basic events
- ▶ Adaptivity
 - ▶ When to discard services, events
 - ▶ keep alive, heartbeats
 - ▶ Invoke new instance of service automatically
 - ▶ Fault tolerance

Issues

- Standards
 - XML, SOAP, WSDL
 - Proprietary Interfaces
- Middleware may be new Operating System
 - Whoever controls it will dominate
 - Not clear if there is or will be a winner
- Integration with web-services
 - Lightweight devices are different
 - May want stateful communication

