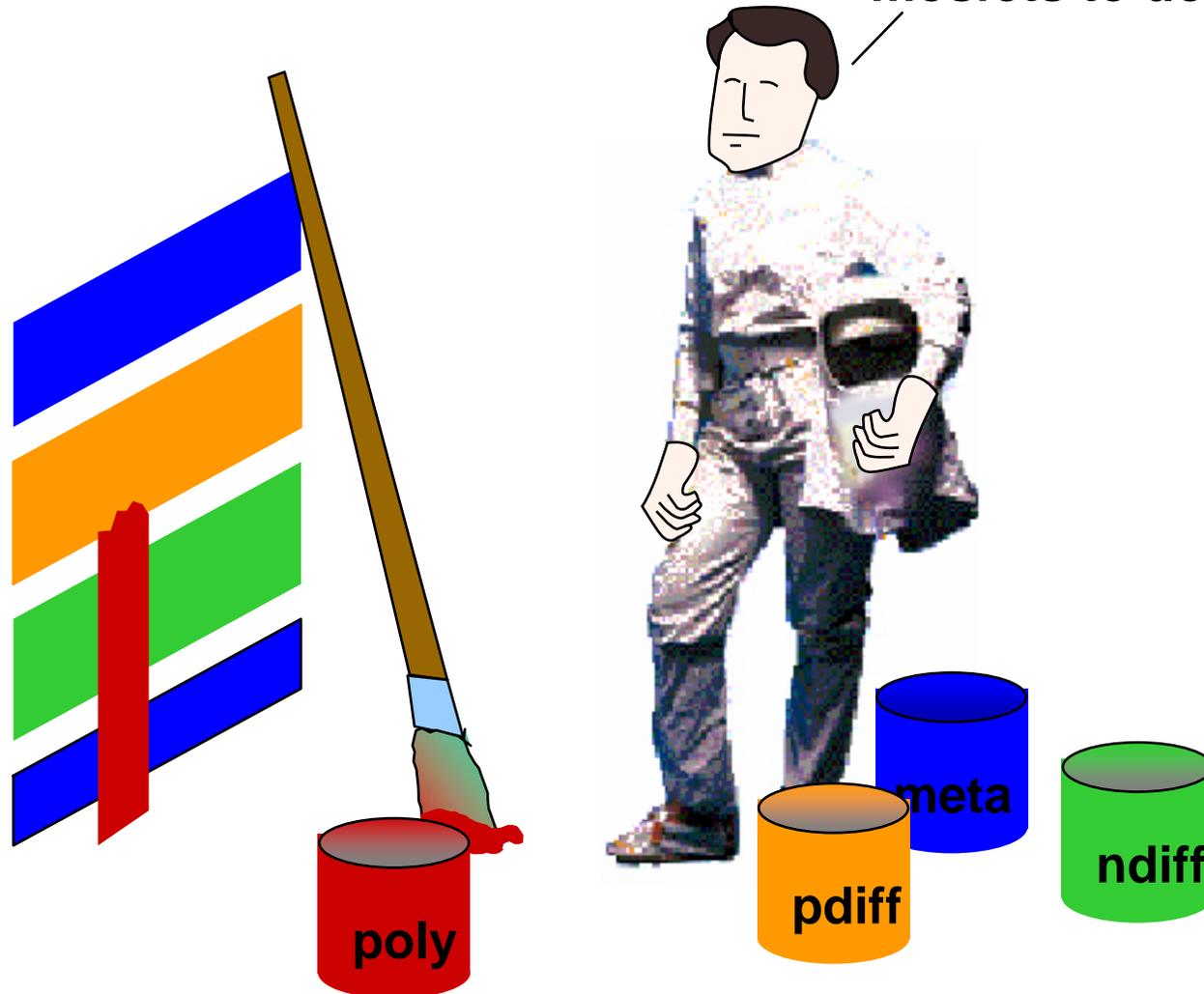


CMOS Technology and Logic Gates

Only 15,432,758 more mosfets to do...



Quality of Design

Quality of a hardware design primarily judged by:

- Price
- Performance
- Power and/or Energy

Other important metrics can include:

- Operating range
 - Temperature, voltage, background radiation
- Reliability
 - Mean-time between failures (MTBF)
- Form factor
 - Size, weight
- Flexibility
 - Tolerance to changes in specification

Need to understand implementation technology to understand tradeoffs among these attributes.

System-Level Impacts

Image removed due to copyright restrictions.

Digital IC in
Package

Chips do not exist in a vacuum, e.g.,
2/2.5G cell phone contains:

- RISC Application Processor (ARM)
- Digital Signal Processor
- SRAM/DRAM Chips
- Flash Memory Chips
- Analog Chips
 - E.g. headphone amplifier
- Radio Chips
- Power Management Subsystem
- Passive components
 - resistors, capacitors and inductors

Need to consider quality of a design
in context of target system.

- E.g., design alternative might have twice the performance but require 10x off-chip memory bandwidth.

[Buss, ISSCC 2002]

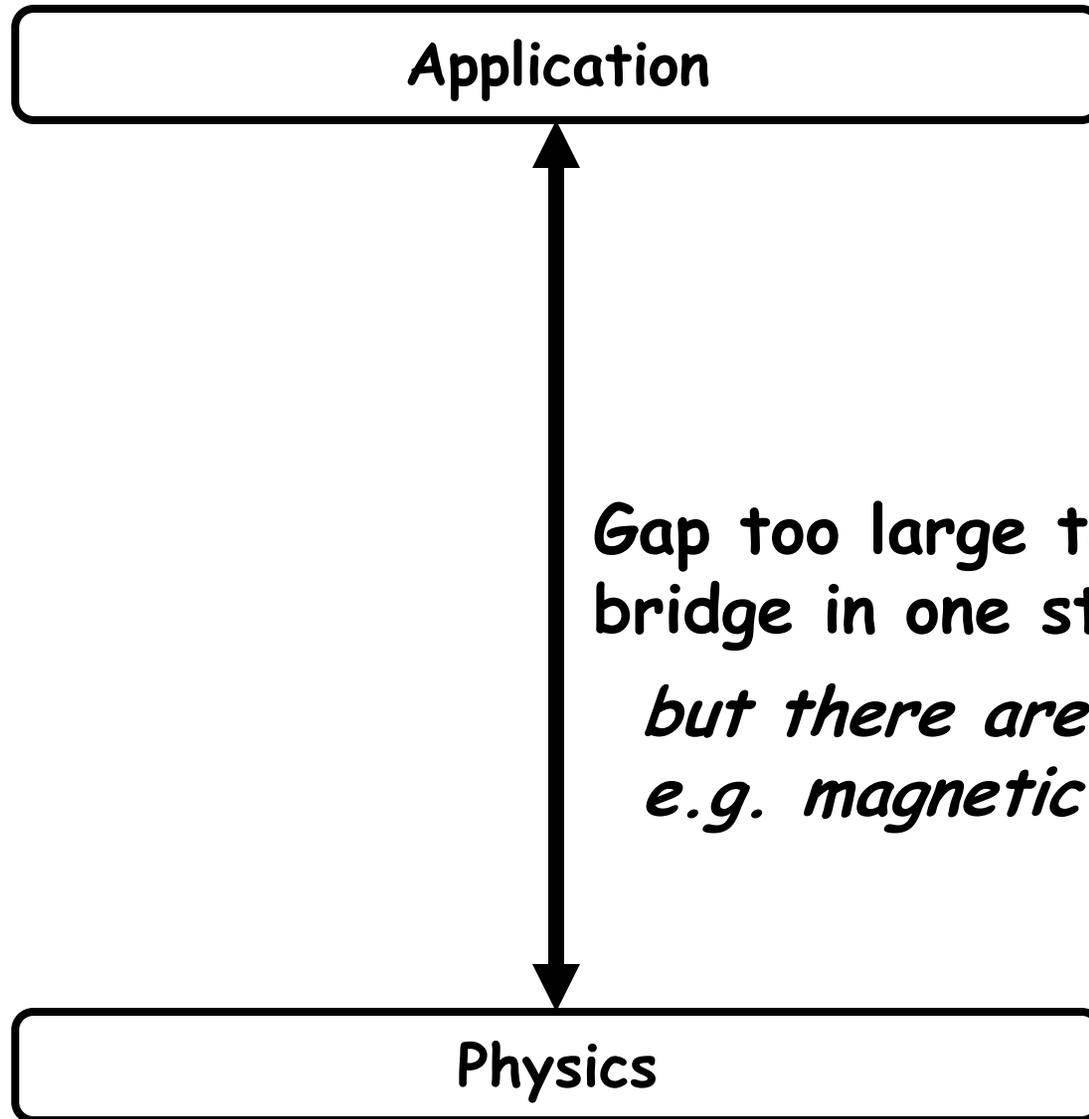
Digital Technology Generations

- Electromechanical Relays
- Vacuum Tubes
- Bipolar Transistors
- CMOS/FET Transistors
 - ~10,000nm gates originally, now down to 90nm in production
 - scaling will stop somewhere below 30nm (over 100 billion trans./chip)
- Future:
 - 3D CMOS (10 trillion transistors/system?)
 - Carbon Nanotubes?
 - Molecular Electronics?

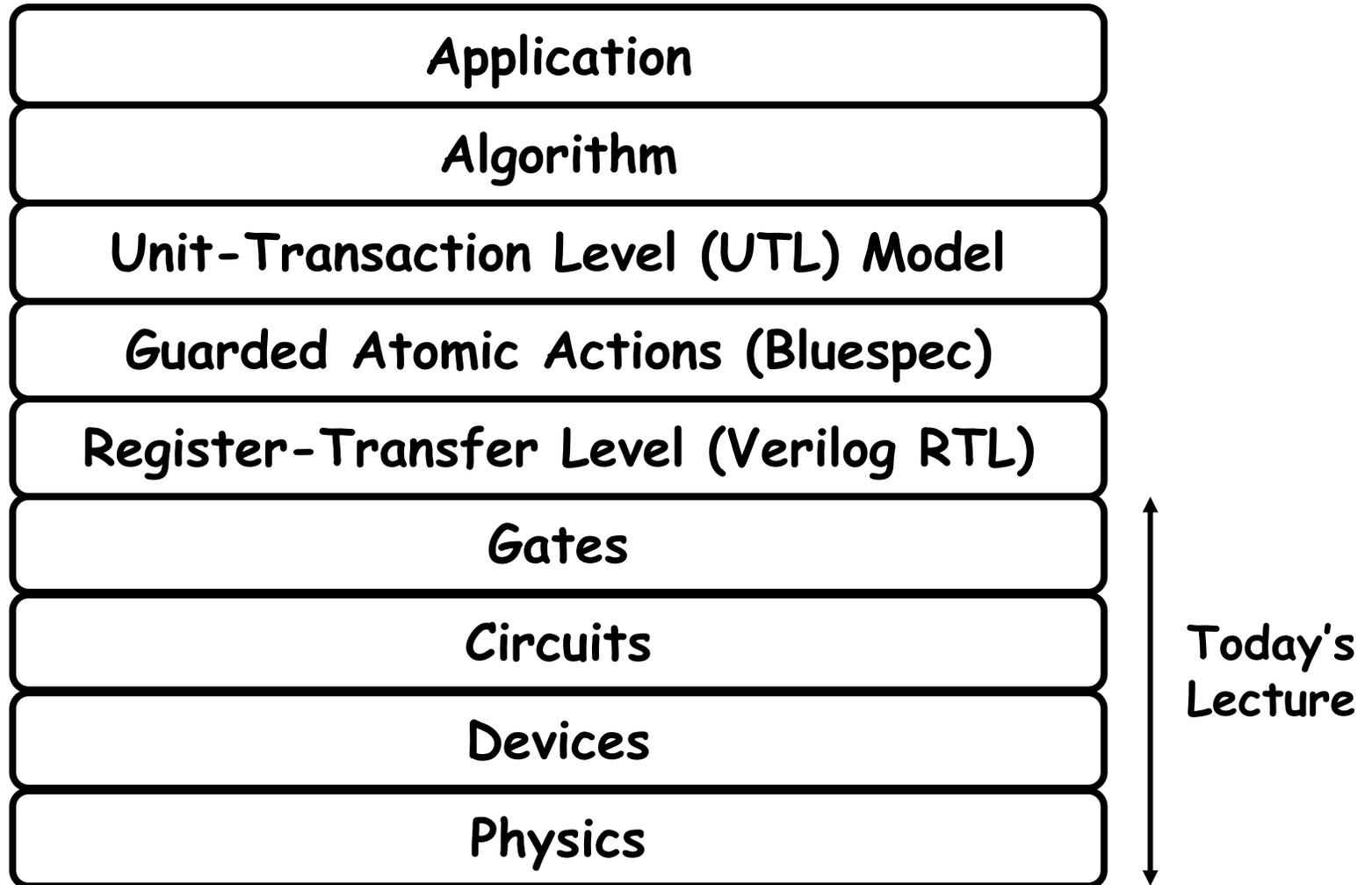
CMOS VLSI is *the* digital implementation technology of choice for the foreseeable future (next 10-20 years)

- Excellent energy versus delay characteristics
- High density of wires and transistors
- Monolithic manufacturing of devices and interconnect, cheap!

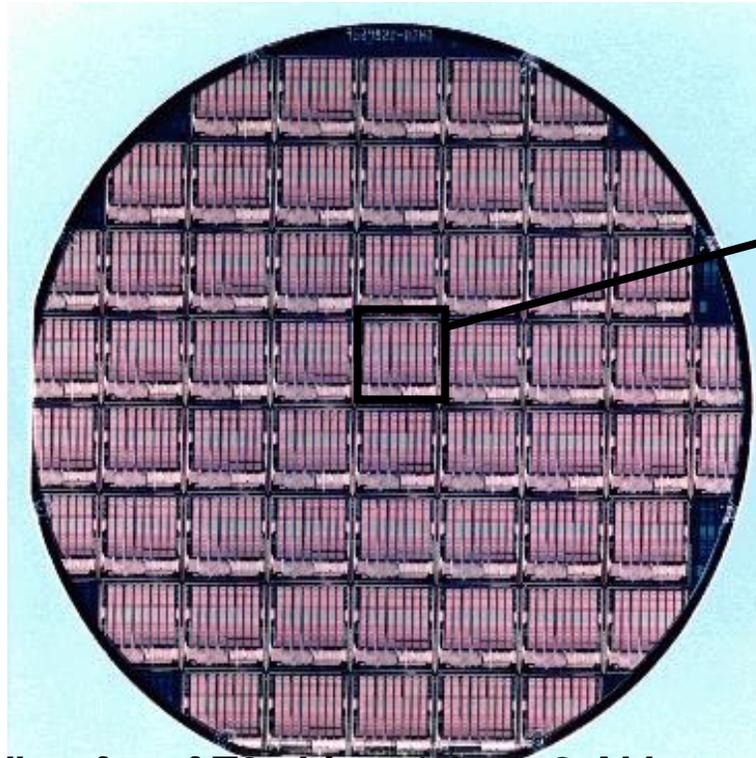
Abstraction Levels in Design



Hardware Design Abstraction Levels



CMOS Fabrication



One chip

[6" wafer of T0 chips, 1.0 μ m, 2 Al layers, 1995]

Starting wafer is pure silicon crystal.

Multiple process steps deposit new materials and etch existing layers using photolithography (light focused through masks).

Modern logic chips fabricated on 20cm (8") wafers, ~100s chips/wafer.

Wafer sawed into separate chips after fabrication.

Chips then placed into packages (*see packaging lecture later in course*)

Basic CMOS Fabrication Steps

Growing silicon dioxide to serve as an insulator between layers deposited on the surface of the silicon wafer.

Doping the silicon substrate with acceptor and donor atoms to create p- and n-type diffusions that form isolating PN junctions and one plate of the MOS capacitor.

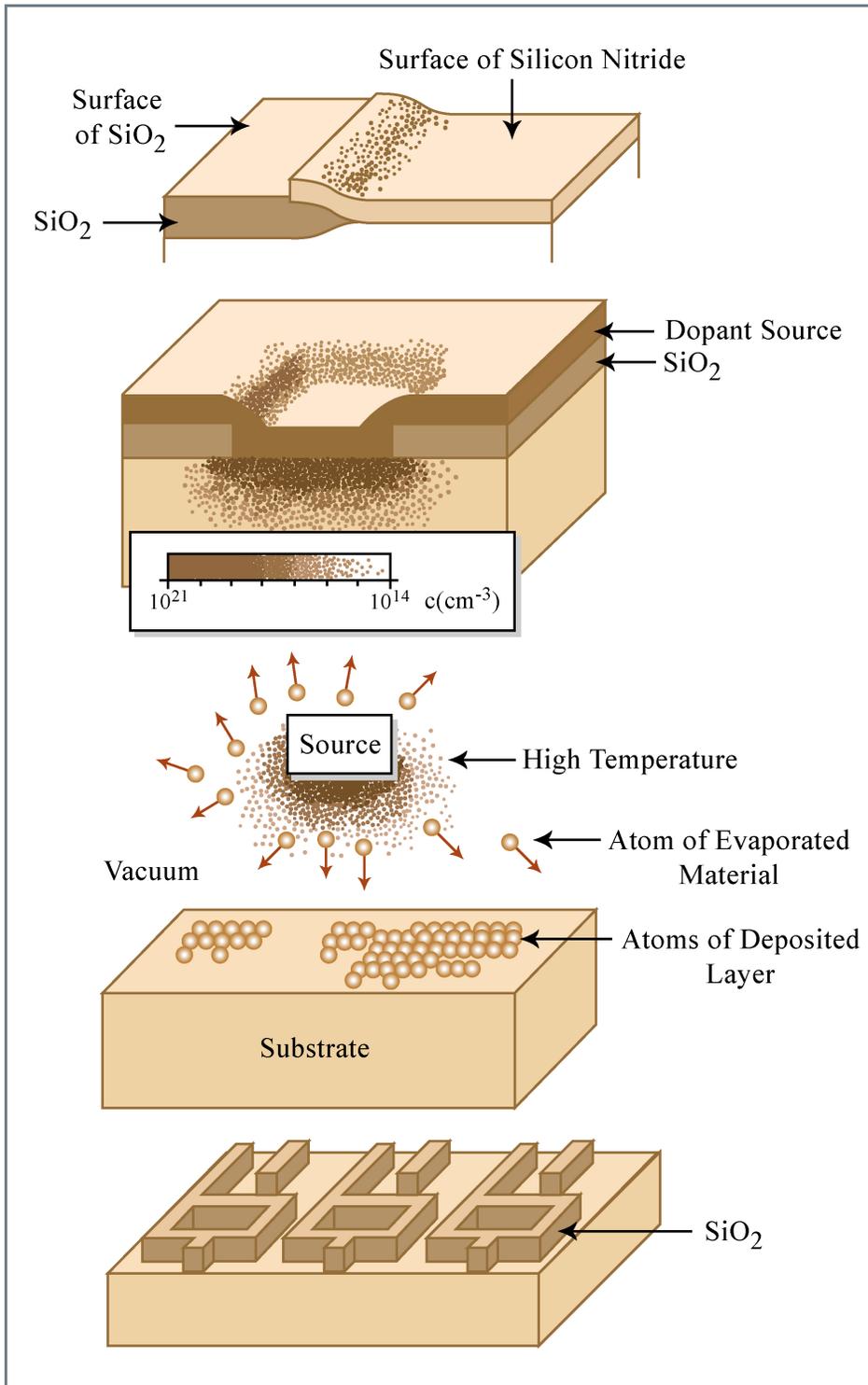
Depositing material on the wafer to create masks, wires and the other plate of the MOS capacitor.

Etching deposited materials to create the appropriate geometric patterns.

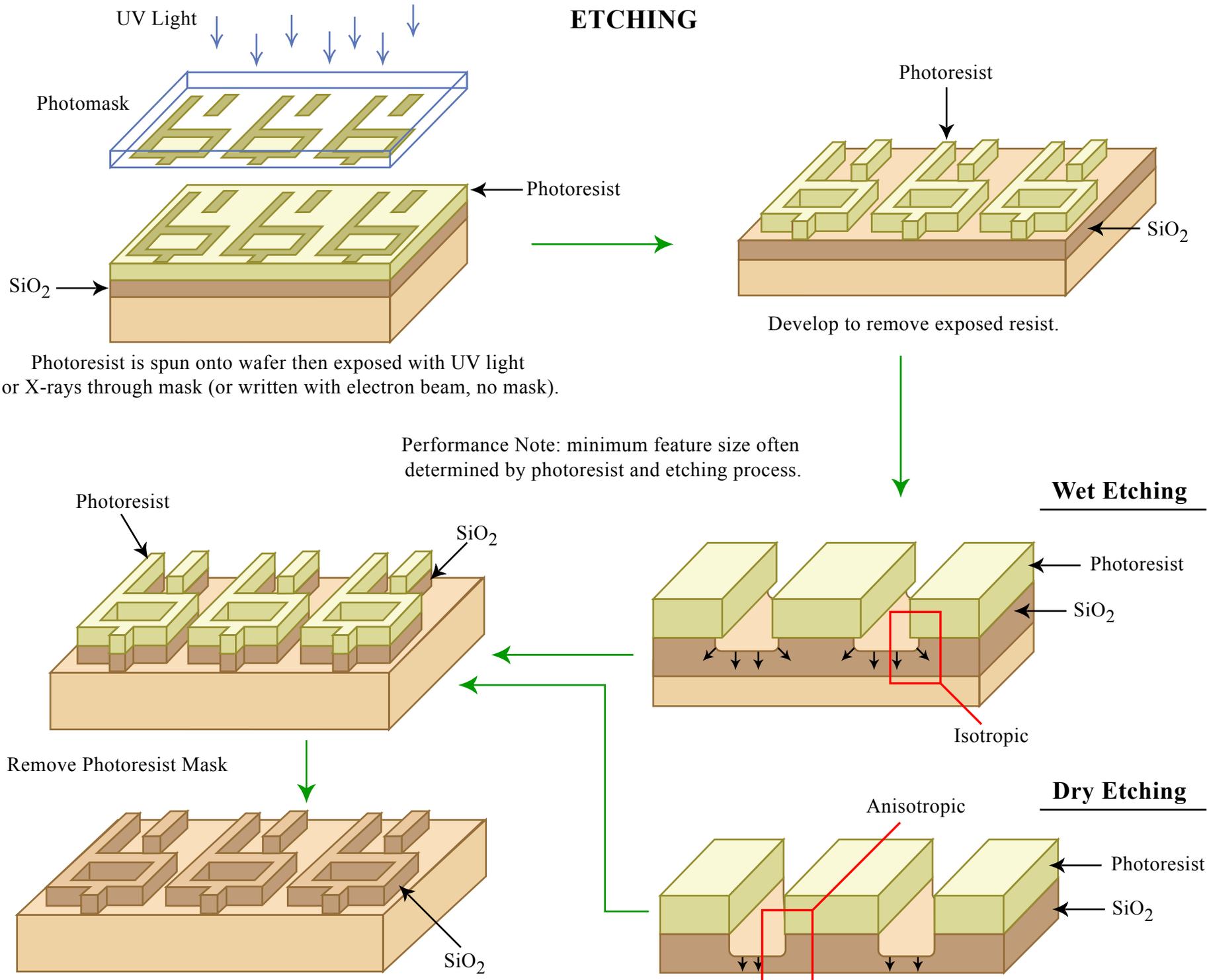
Figures by MIT OCW.

Adapted from Maly, W.

Atlas of IC Technologies: An Introduction to VLSI Processes.
(Ignore dimensions in figures)

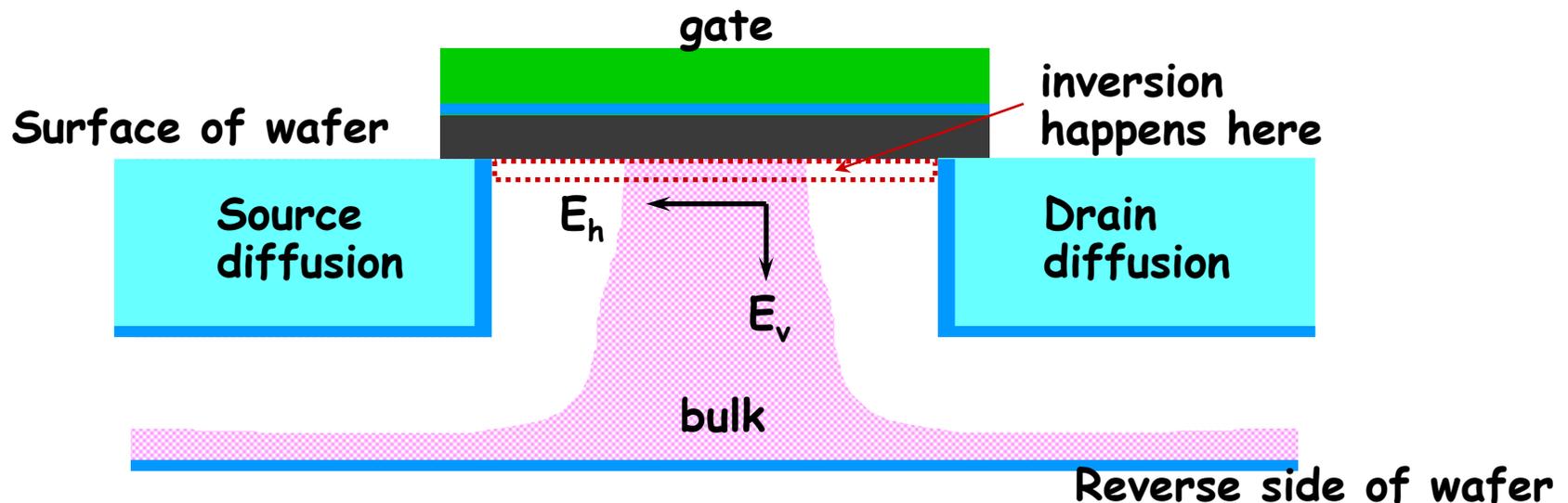


ETCHING



FET = Field-Effect Transistor

The four terminals of a fet (gate, source, drain and bulk) connect to conducting surfaces that generate a complicated set of electric fields in the channel region which depend on the relative voltages of each terminal.



INVERSION:

A sufficiently strong vertical field will attract enough electrons to the surface to create a conducting n-type channel between the source and drain.

CONDUCTION:

If a channel exists, a horizontal field will cause a drift current from the drain to the source.

Multiple Levels of Interconnect

IBM photomicrograph (Si has been removed!)

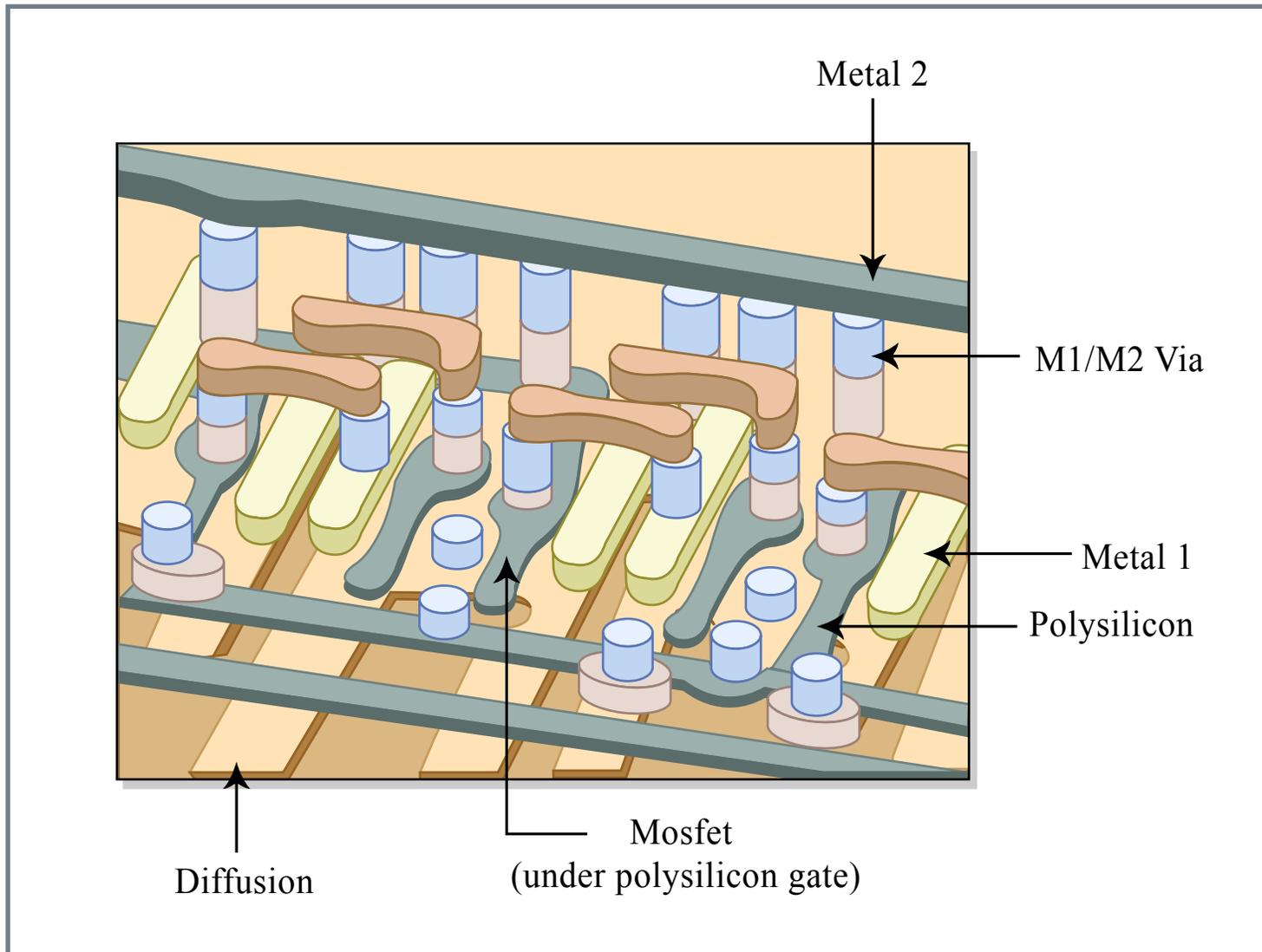
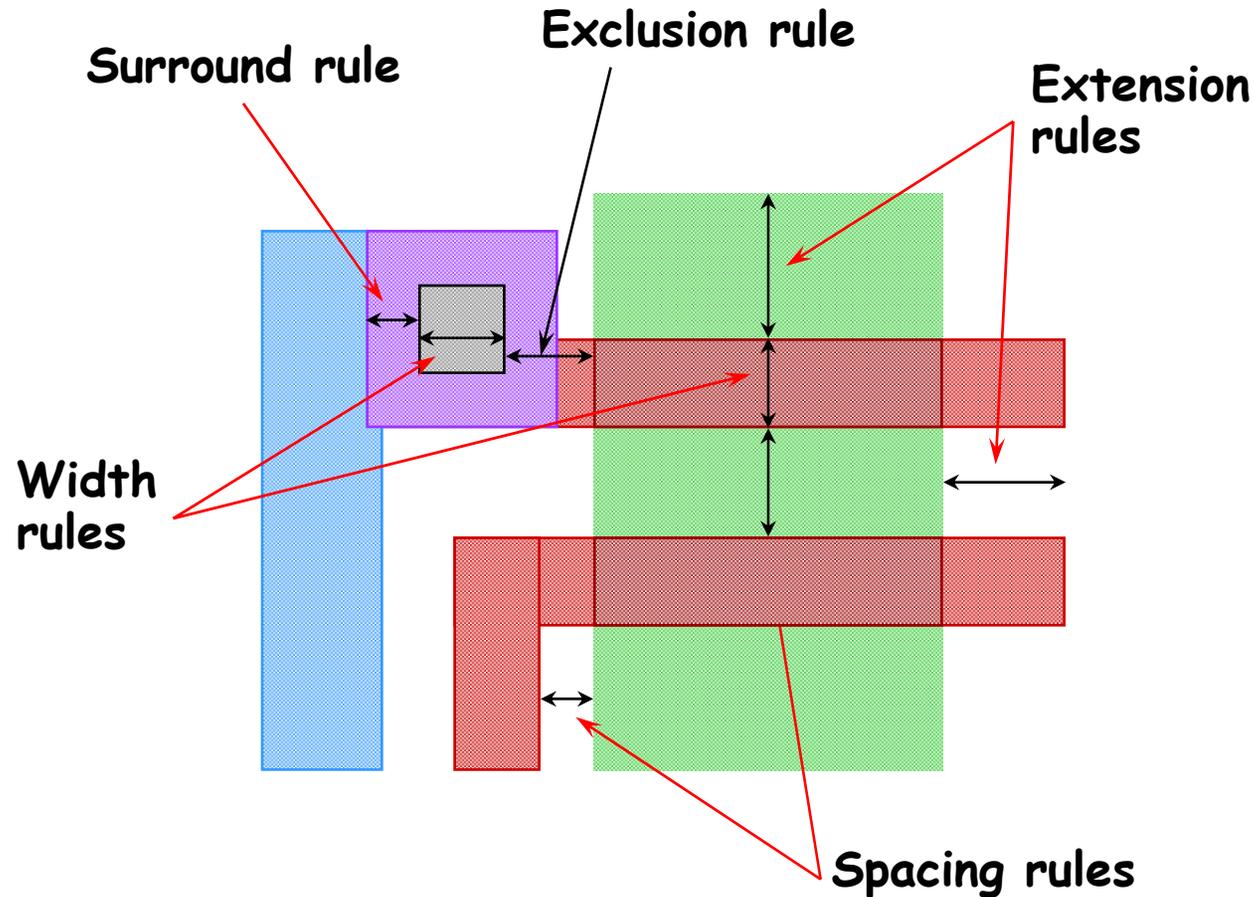


Figure by MIT OCW.
2/07/2005

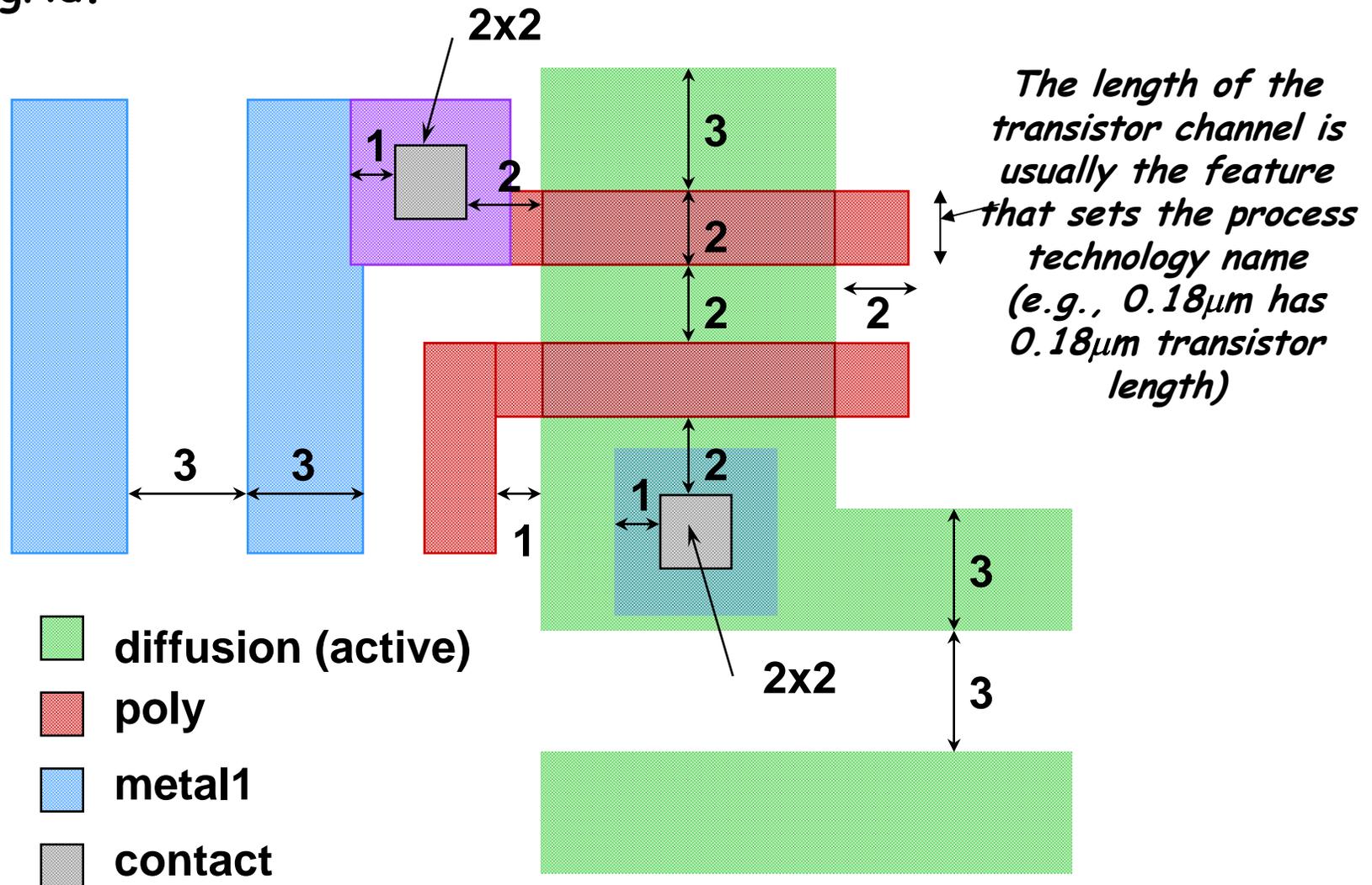
Design Rules



- Design rules are an abstraction of the fabrication process that specify various geometric constraints on how different masks can be drawn.
- Design rules can be absolute measurements (e.g. in nm) or scaled to an abstract unit, the *lambda*. Lambda-based designs are scaled to the appropriate absolute units depending on the manufacturing process finally used.

Lambda-based Design Rules

One lambda (λ) = one half of the "minimum" mask dimension.
Typically the length of a transistor channel is 2λ . Usually all edges must be "on grid", e.g., in the MOSIS scalable rules, all edges must be on a lambda grid.



Static CMOS Gates

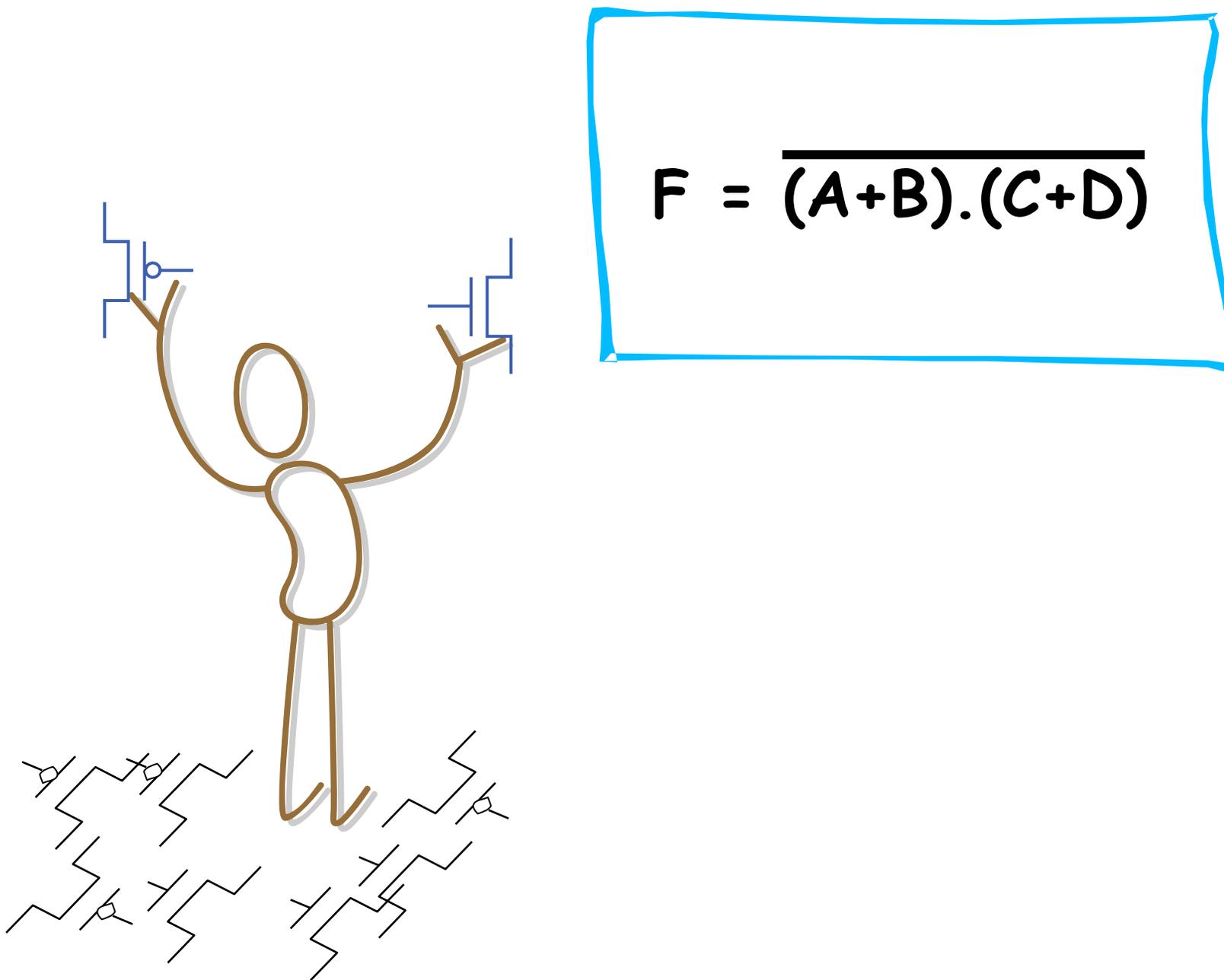
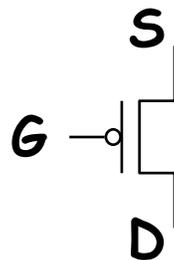


Figure by MIT OCW.

Simplified FET Model

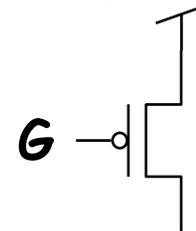
Binary logic values represented by voltages:

"High" = Supply Voltage, "Low" = Ground Voltage

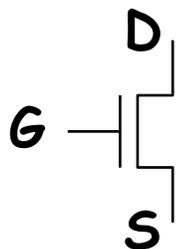


PFET connects
S and D when
 $G = \text{"low"} = 0V$

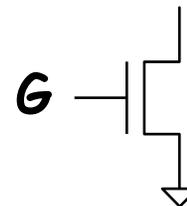
Supply Voltage = V_{DD}



PFET only good
at pulling up



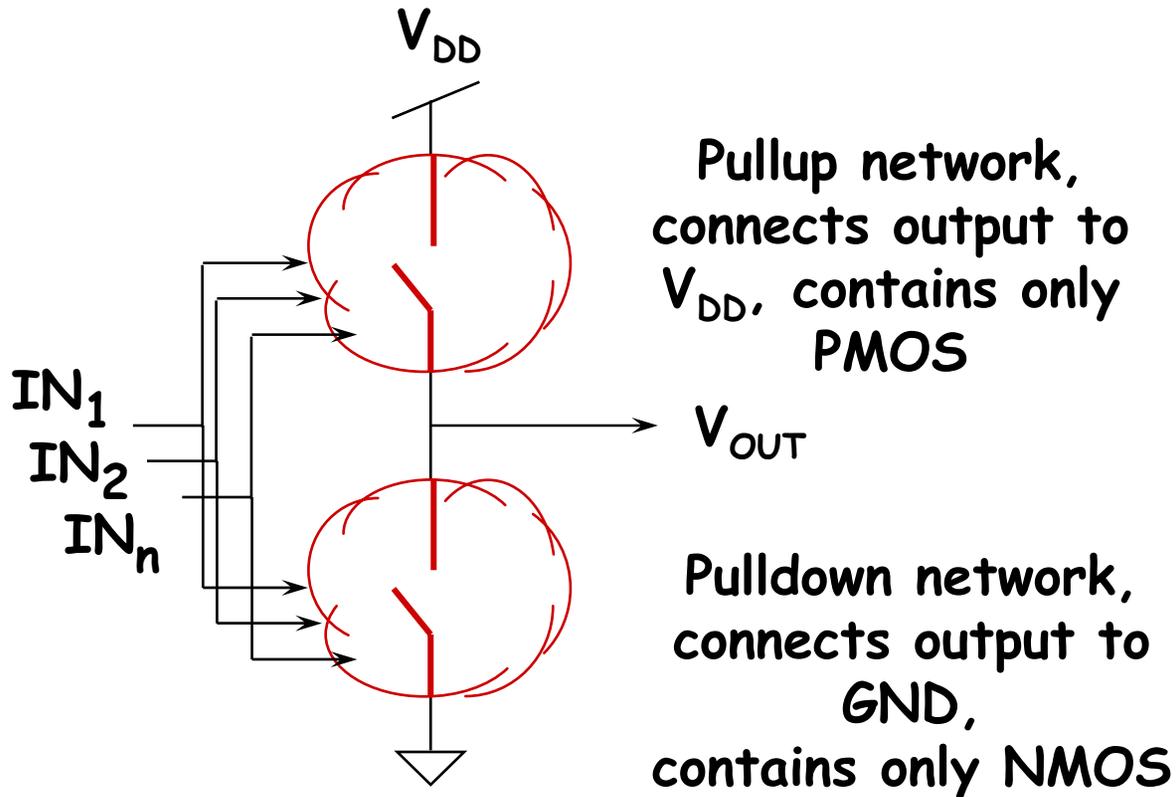
NFET connects
D and S when
 $G = \text{"high"} = V_{DD}$



NFET only good
at pulling down

Ground = GND = 0V

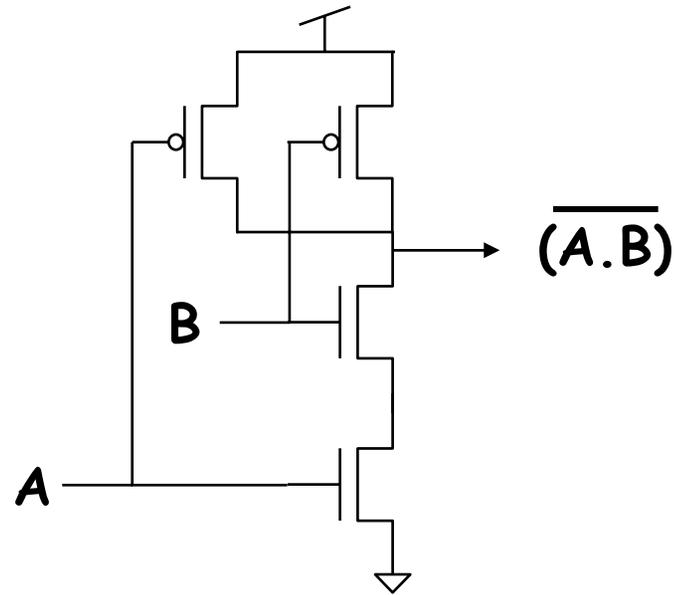
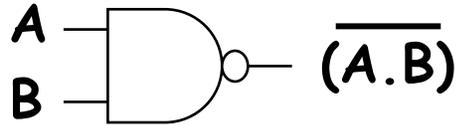
Generic Static CMOS Gate



For every set of input logic values, either pullup or pulldown network makes connection to V_{DD} or GND

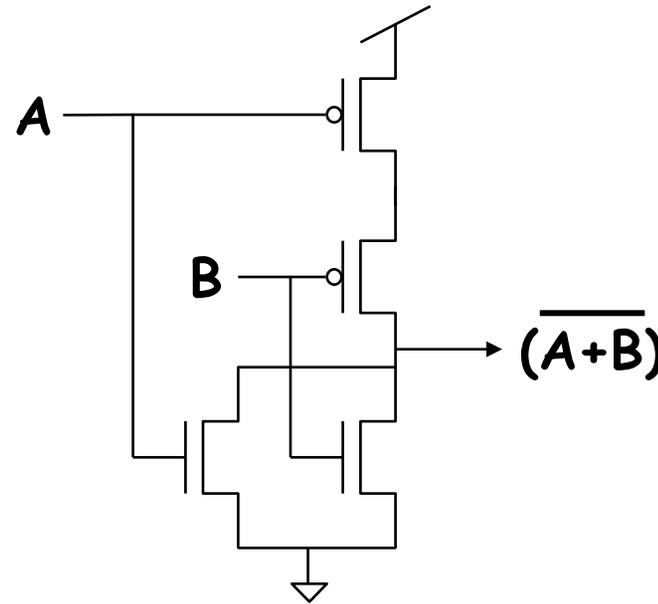
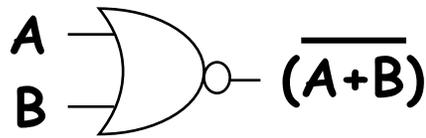
- If both connected, power rails would be shorted together
- If neither connected, output would float (tristate logic)

NAND Gate



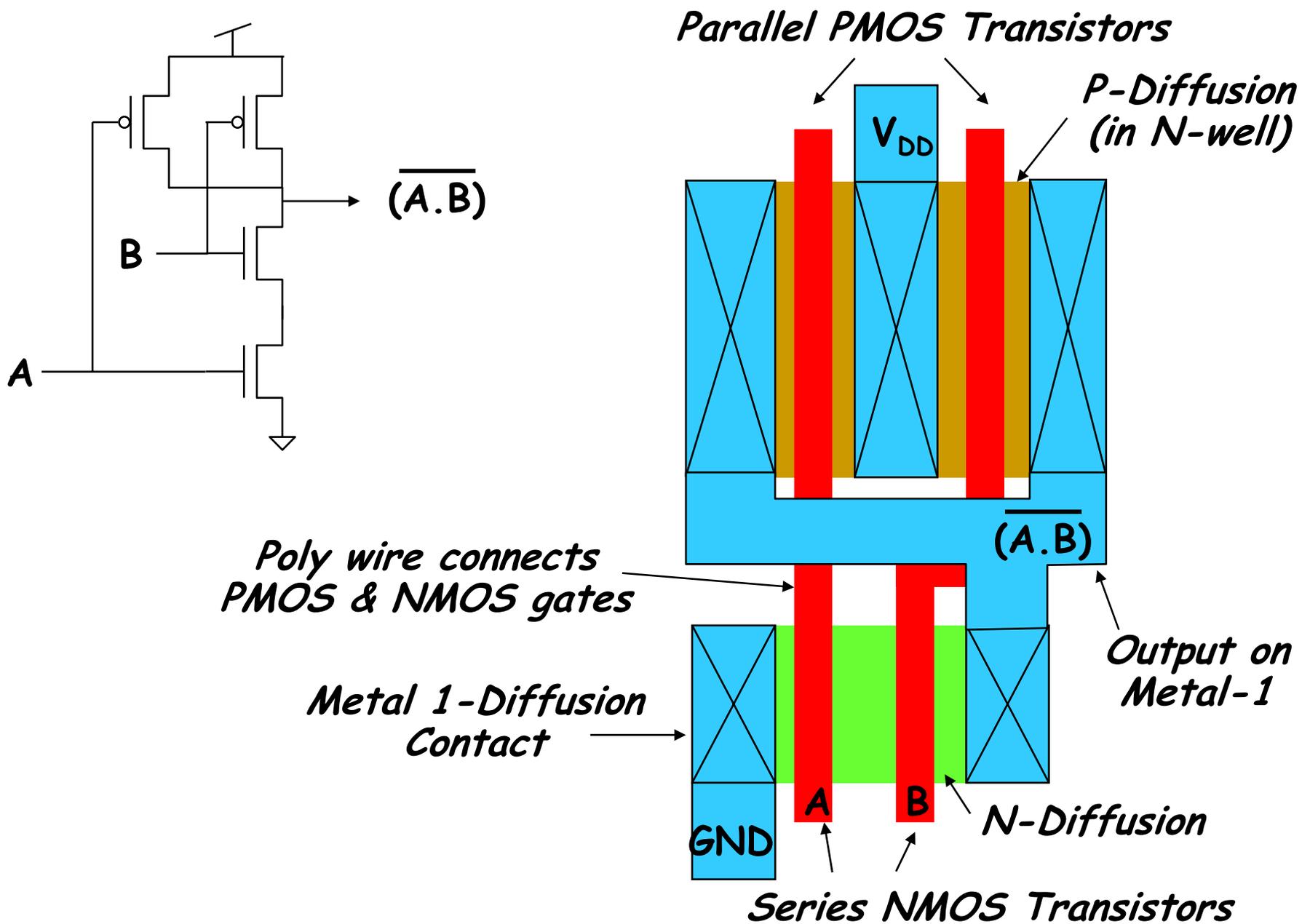
- When both A and B are high, output is low
- When either A or B is low, output is high

NOR Gate



- When both A and B are low, output is high
- When either A or B is high, output is low

NAND Gate Layout



Methodical Gate Building

Goal is to create a logic function $f(x_1, x_2, \dots)$

- must be inverting for single level of CMOS logic

Pull up network should connect output to V_{DD} when

$$f(x_1, x_2, \dots) = 1$$

Pull down network should connect output to GND

$$\text{when } \overline{f}(x_1, x_2, \dots) = 1$$

Because PMOS is conducting with low inputs, useful to write pullup as function of inverted inputs

$$p(\overline{x}_1, \overline{x}_2, \dots) = f(x_1, x_2, \dots)$$

Pullup is Dual of Pulldown Network

For NAND gate, $f = \overline{A \cdot B}$

Pulldown $f = A \cdot B$

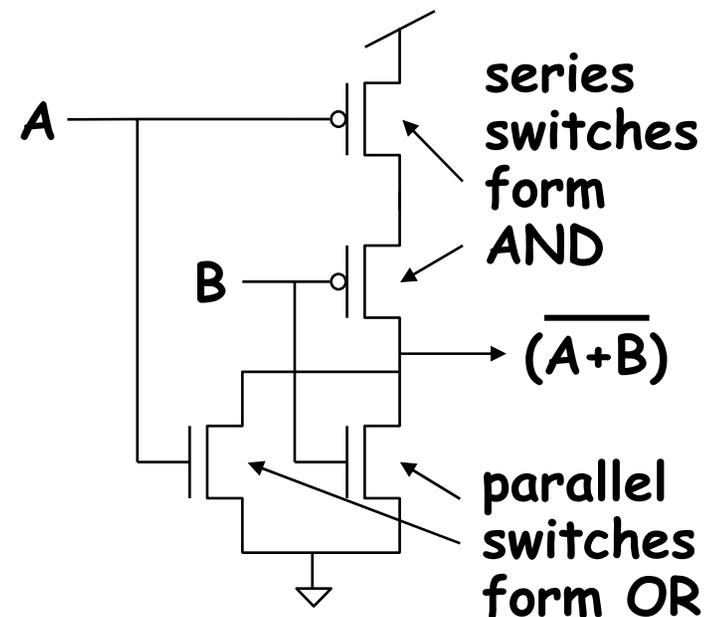
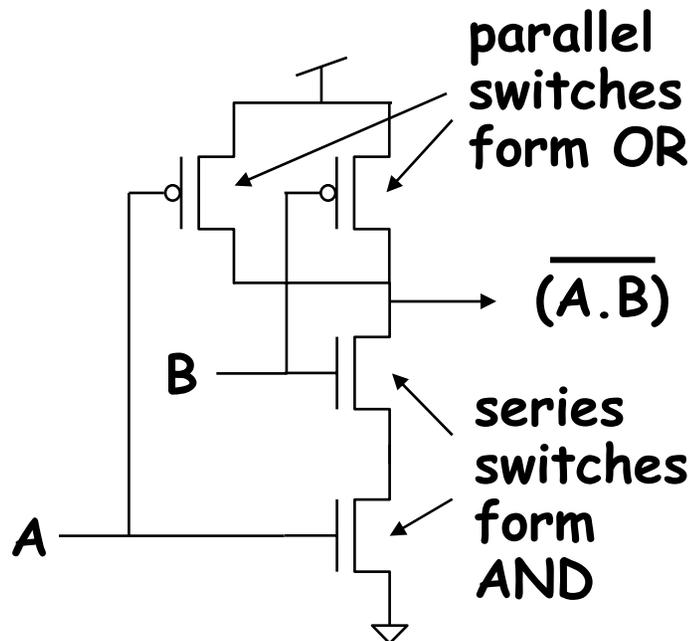
Pullup $p = f = \overline{A \cdot B}$
 $= \overline{A} + \overline{B}$

(De Morgan's Laws)

For NOR gate, $f = \overline{A + B}$

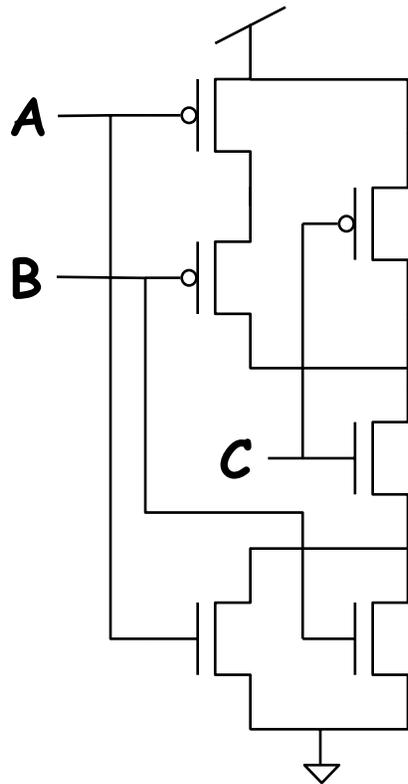
Pulldown $f = A + B$

Pullup $p = f = \overline{A + B}$
 $= \overline{A} \cdot \overline{B}$



More Complex Example

$$f = \overline{(A+B).C}$$



$$\text{pullup } p = \overline{(A+B).C}$$

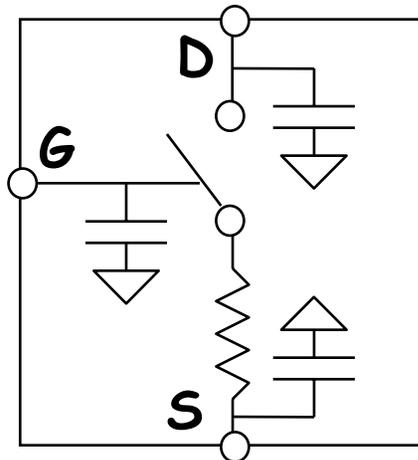
$$= \overline{(A+B)} + \overline{C}$$

$$= (\overline{A}.\overline{B}) + \overline{C}$$

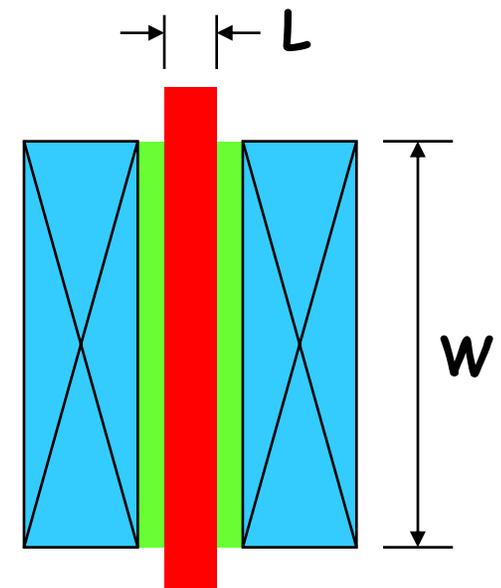
$$\text{pulldown } \overline{f} = (A+B).C$$

Transistor R and C

Simple Equivalent RC Model of Transistor

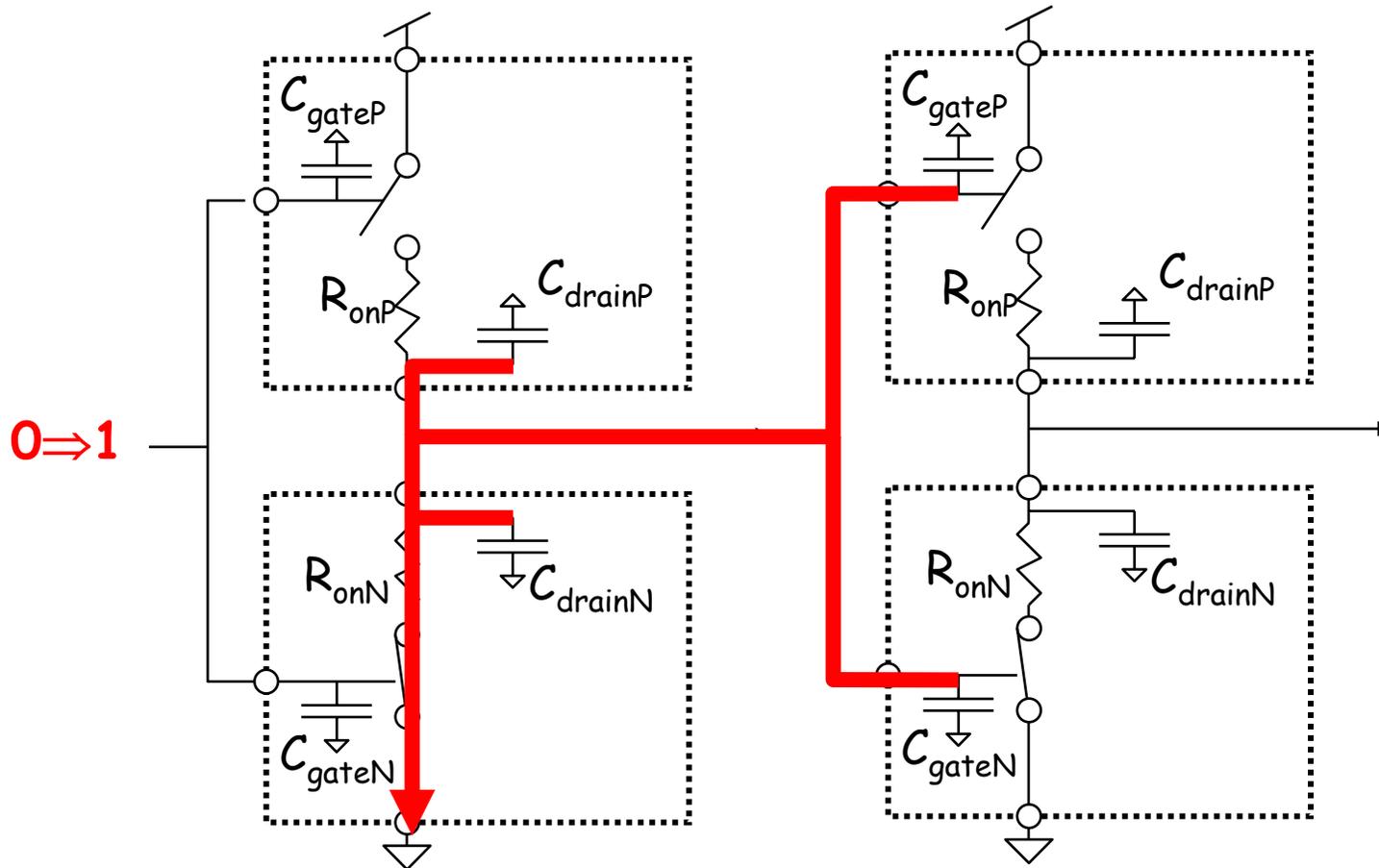


- Nearly all transistors in digital CMOS circuits have minimum L
 - but might use slightly longer L to cut leakage in parts of modern circuits
- Can scale transistor R and C parameters by width W
- Effective R scales linearly with $1/W$
 - $\sim 4\text{k}\Omega\mu\text{m}$ NMOS, $\sim 9\text{k}\Omega\mu\text{m}$ PMOS, in $0.25\mu\text{m}$ technology
- Gate capacitance scales linearly with W
 - $\sim 2\text{fF}/\mu\text{m}$
- Diffusion capacitance scales linearly with W
 - sum contributions from perimeter and area, $\sim 2\text{fF}/\mu\text{m}$



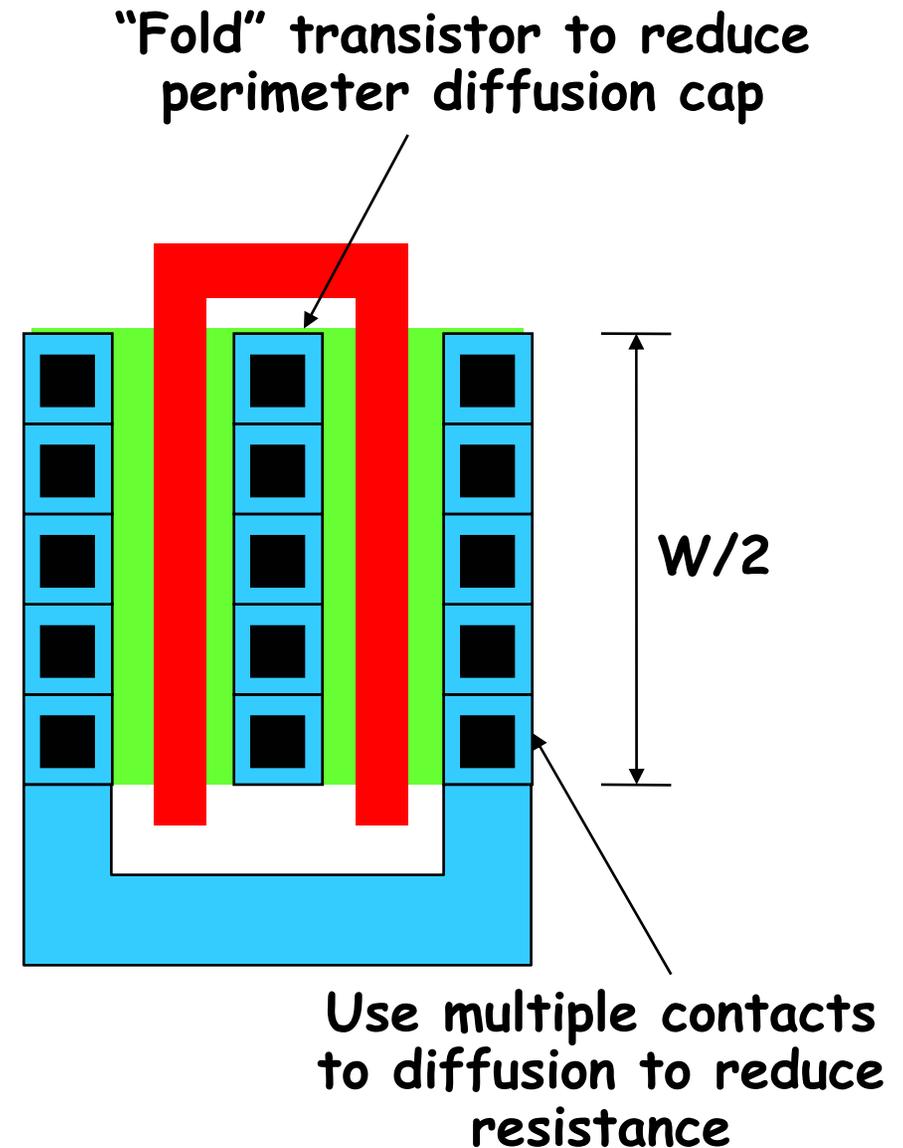
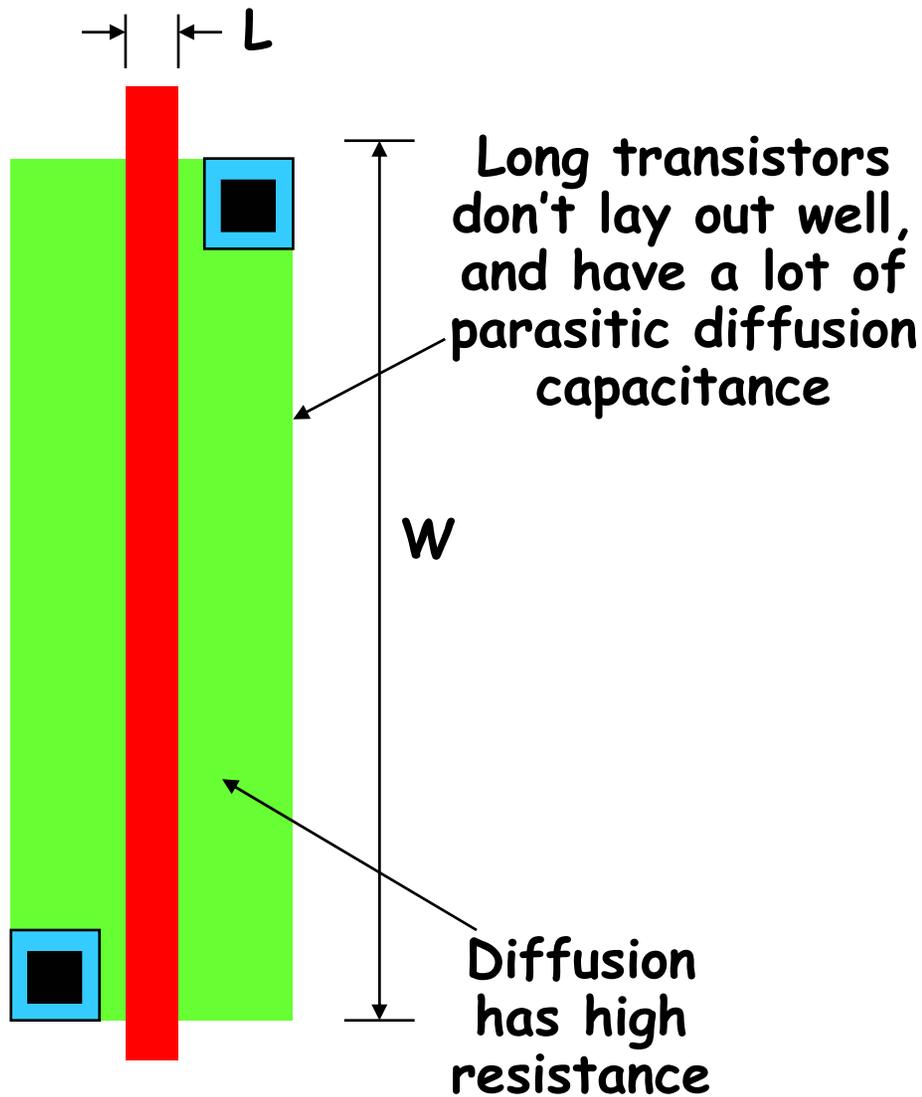
Transistor Delay

When one gate drives another, all capacitance on the node must be charged or discharged to change voltage to new state. Delay is proportional to driving resistance and connected capacitance.

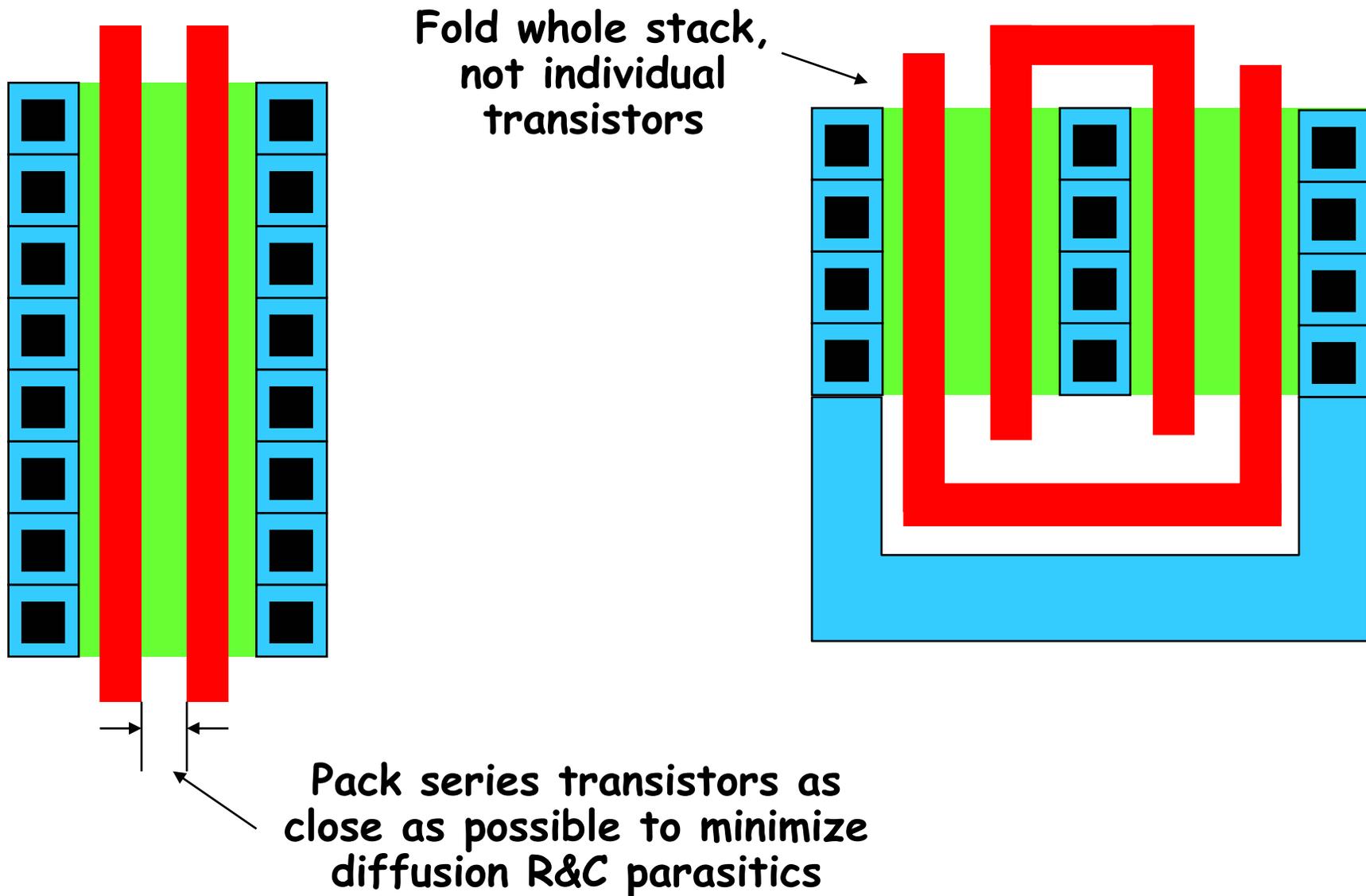


$$\text{Delay} \sim R_{onN}(C_{drainN} + C_{drainP} + C_{gateP} + C_{gateN})$$

Gate Layout Tricks



More Layout Tricks



Even More Complex Gates?

Can build arbitrarily complex logic function into one gate,
e.g.

?

$$F = \overline{(A+B) \cdot (C+D) \cdot E \cdot G + H \cdot (J+K)}$$

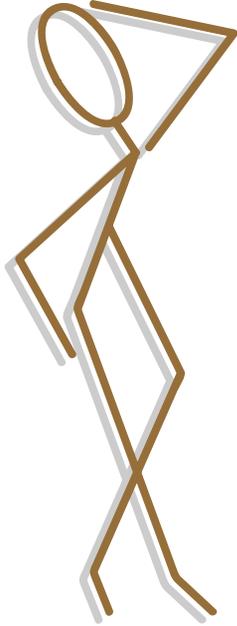


Figure by MIT OCW.

- But don't want to:
 - Usually less total delay using a few smaller logic gates rather than one large complex gate
 - Only want to design and characterize a small library of gates
- What's the best way to implement a given logic function?

Figure by MIT OCW.

Method of Logical Effort (Sutherland and Sproul)

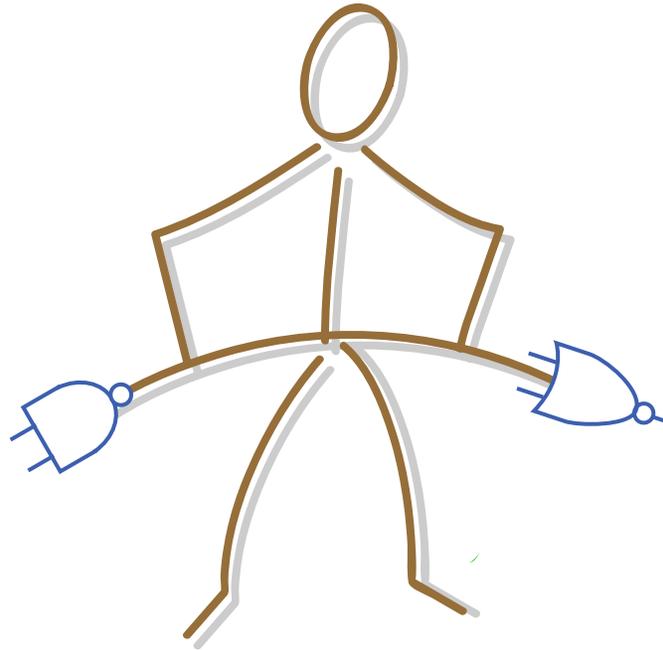


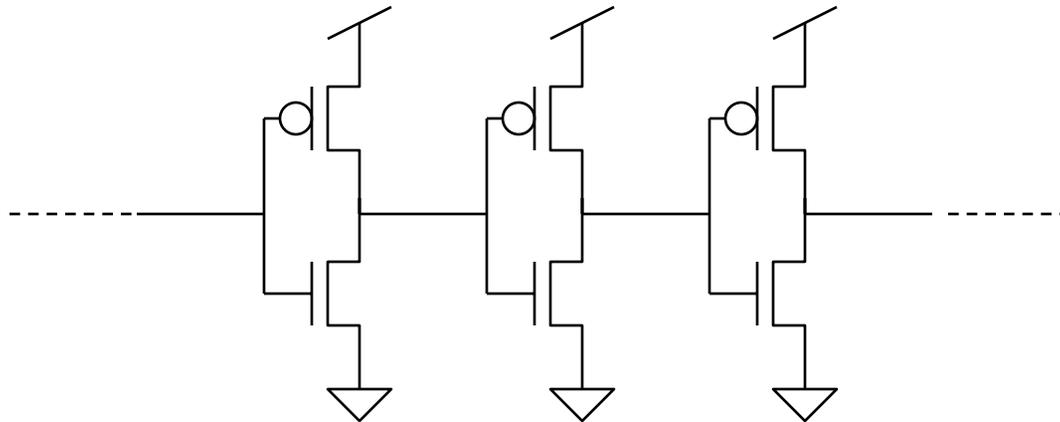
Figure by MIT OCW.

- Easy way to estimate delays in CMOS process.
- Indicates correct number of logic stages and transistor sizes.
- Based on simple RC approximations.
- Useful for back-of-the-envelope circuit design and to give insight into results of synthesis.

Technology Speed Parameter: τ

Characterize process speed with single delay parameter: τ

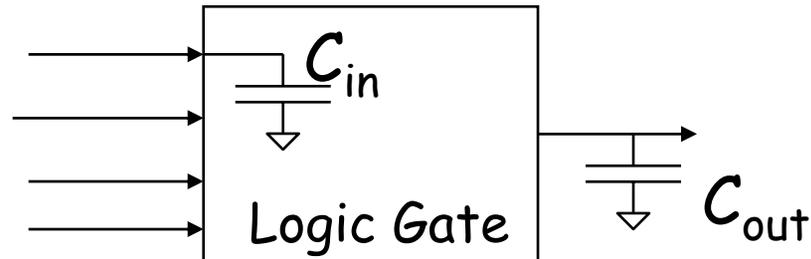
τ is delay of inverter driving same-sized inverter, with no parasitics other than gate



$\tau \sim 16\text{-}20\text{ps}$ for $0.25\mu\text{m}$ process

Gate Delay Components

Delay = Logical Effort \times Electrical Effort + Parasitic Delay



Logical Effort

Complexity of logic function (Invert, NAND, NOR, etc)

Define inverter has logical effort = 1

Depends only on topology not transistor sizing

Electrical Effort

Ratio of output capacitance to input capacitance C_{out}/C_{in}

Parasitic Delay

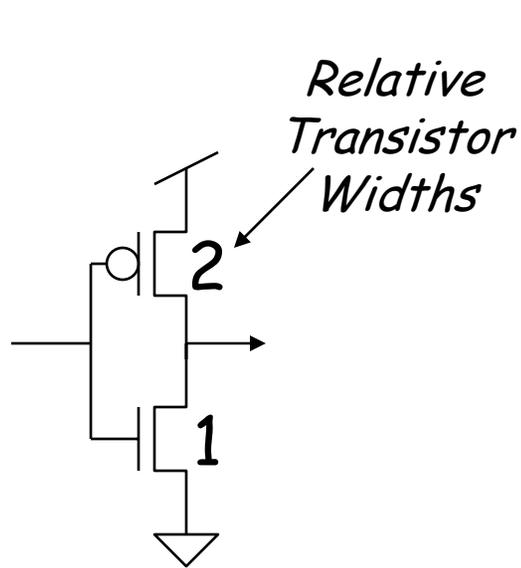
Intrinsic self-loading of gate

Independent of transistor sizes and output load

Logical Effort for Simple Gates

Define Logical Effort of Inverter = 1

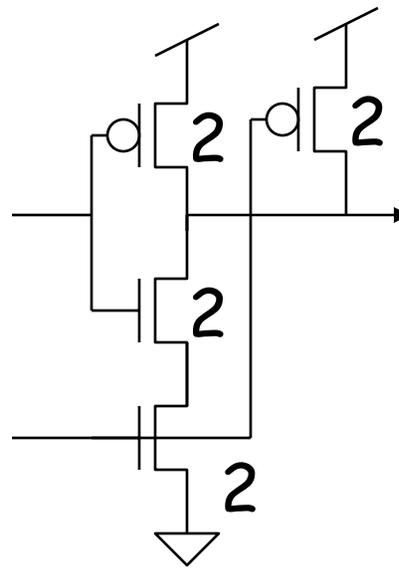
For other gates, Logical Effort is ratio of logic gate's input cap. to inverter's input cap., when gate sized to give same current drive as inverter



Inverter

Input Cap = 3 units

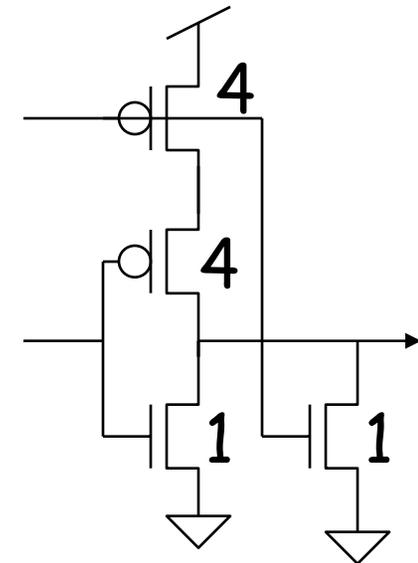
L.E.=1 (definition)



NAND

Input Cap = 4 units

L.E.=4/3

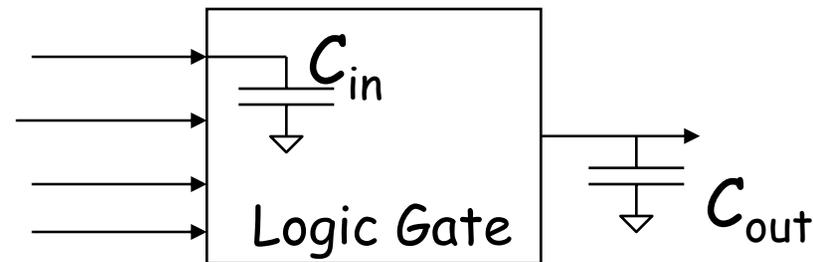


NOR

Input Cap = 5 units

L.E.=5/3

Electrical Effort



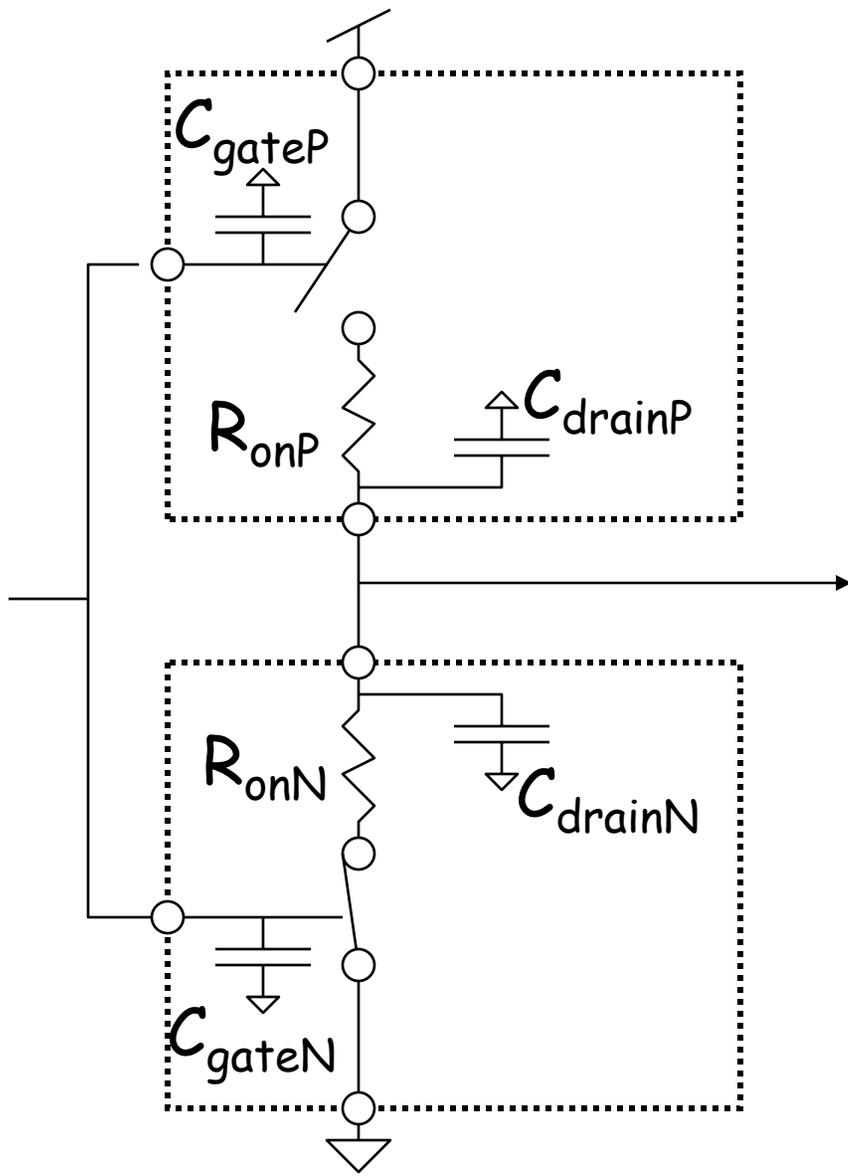
Ratio of output load capacitance over input capacitance:

$$\text{Electrical Effort} = C_{out}/C_{in}$$

Usually, transistors have minimum length

Input and output capacitances can be measured in units of transistor gate widths

Parasitic Delay



Main cause is drain diffusion capacitances.

These scale with transistor width so P.D. independent of transistor sizes.

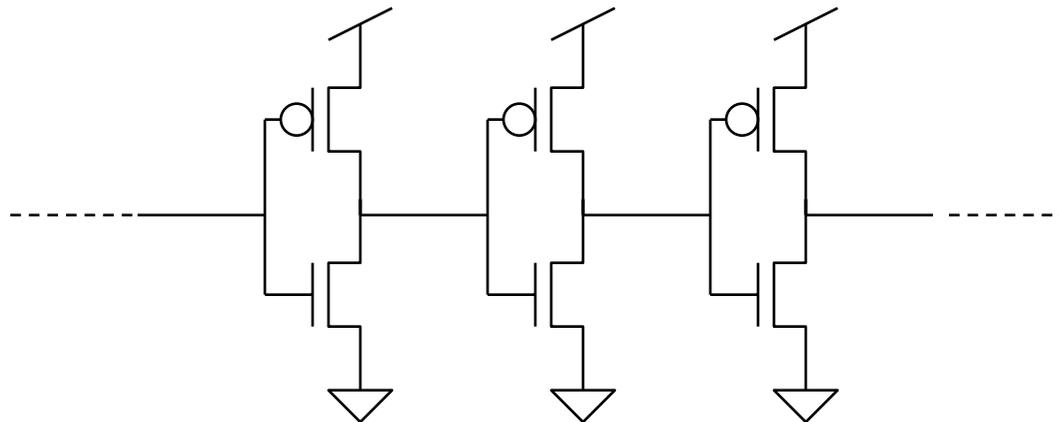
Useful approximation:

$$C_{gate} \approx C_{drain}$$

For inverter:

$$\text{Parasitic Delay} \approx 1.0 \tau$$

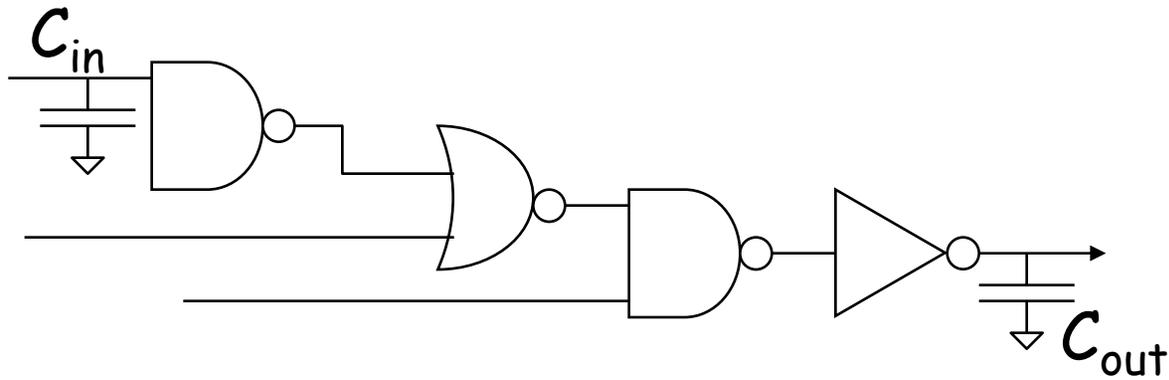
Inverter Chain Delay



For each stage:

$$\begin{aligned}\text{Delay} &= \text{Logical Effort} \times \text{Electrical Effort} + \text{Parasitic Delay} \\ &= 1.0 \text{ (definition)} \times 1.0 \text{ (in = out)} + 1.0 \text{ (drain } C) \\ &= 2.0 \text{ units}\end{aligned}$$

Optimizing Circuit Paths



Path logical effort, $G = \prod g_i$ ($g_i = \text{L.E. stage } i$)

Path electrical effort, $H = C_{\text{out}}/C_{\text{in}}$ ($h_i = \text{E.E. stage } i$)

Parasitic delay, $P = \sum p_i$ ($p_i = \text{P.D. stage } i$)

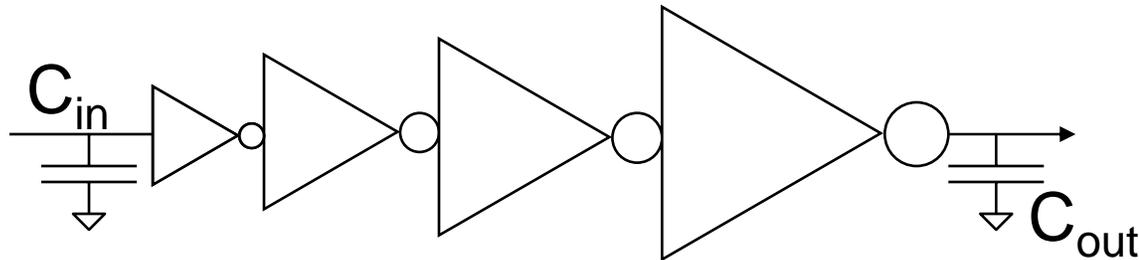
Path effort, $F = GH$

Minimum delay when each of N stages has equal effort

$$\text{Min. } D = NF^{1/N} + P$$

$$\text{i.e. } g_i h_i = F^{1/N}$$

Optimal Number of Stages



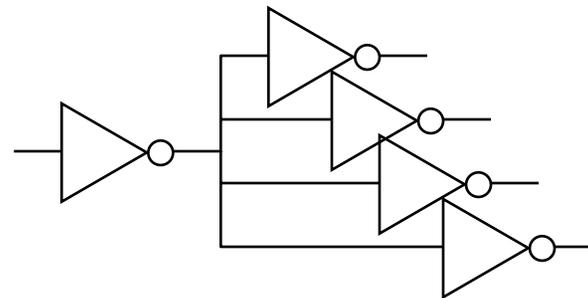
Minimum delay when:

stage effort = logical effort \times electrical effort $\approx 3.4-3.8$

- Some derivations have $e = 2.718..$ as best stage effort - this ignores parasitics
- Broad optimum, stage efforts of 2.4-6.0 within 15-20% of minimum

Fan-out-of-four (FO4) is convenient design size ($\sim 5\tau$)

FO4 delay: Delay of inverter driving four copies of itself



Using Logical Effort in Design

For given function, pick candidate gate topology

Determine optimal stage effort

- equal for all stages

Starting at last gate

- output load is known
- logical effort is known (from gate topology)
- calculate transistor size to give required stage effort
- gives output load for preceding stage
- lather, rinse, repeat...

Can modify stage efforts up or down to reduce area, power, or to fit fixed set of library cells

- optimal sizing has broad optimum

In 6.884, we'll just let synthesis tool handle gate sizing, but it's useful to know why the tool makes certain decisions.