Wires

Figure by MIT OCW.
Wires are an Old Problem

Cray-1, 1976

Cray-3, 1993

Cray-1 Wiring

Cray-3 wiring

Courtesy of Cray Company. Used with permission.
Interconnect Problems

- A lot of circuit designers are very worried about what's happening with wires in CMOS technology.
- Device technology has been scaling well, with gate performance increasing linearly with decreasing feature size.
- Wires scale differently, and long wires have been getting relatively slower over time.
- Wire delay is a function of wire resistance and capacitance.
Modern Interconnect Stack

IBM CMOS7 process
6 layers of copper wiring
1 layer of tungsten local interconnect

Images removed due to copyright restrictions.
Wire Resistance

\[
\text{resistance} = \frac{(\text{length} \times \text{resistivity})}{(\text{height} \times \text{width})}
\]

- bulk aluminum \(2.8 \times 10^{-8} \ \Omega\cdot m\)
- bulk copper \(1.7 \times 10^{-8} \ \Omega\cdot m\)
- bulk silver \(1.6 \times 10^{-8} \ \Omega\cdot m\)

- Height (Thickness) fixed in given manufacturing process
- Resistances quoted as \(\Omega/\text{square}\)
- TSMC 0.18\(\mu\)m 6 Aluminum metal layers
  - M1-5 0.08 \(\Omega/\text{square}\)  (0.5 \(\mu\)m x 1mm wire = 160 \(\Omega\))
  - M6 0.03 \(\Omega/\text{square}\)  (0.5 \(\mu\)m x 1mm wire = 60 \(\Omega\))
Local Interconnect

- Use contact material (tungsten) to provide extra layer of connectivity below metal 1
- Can also play same trick with silicided poly to connect gates to diffusion directly in RAMs
- Typically used to shrink memory cells or standard cells
- Contacts directly to poly gate or diffusion
Via Resistance

- Via resistance significant
- TSMC 0.18\(\mu\)m 6-Al

\[
\begin{array}{c|c}
\text{Layer} & \text{Resistance} \\
\hline
\text{Diff-M1} & 11.0 \, \Omega \\
\text{Poly-M1} & 10.4 \, \Omega \\
\text{M2-M1} & 4.5 \, \Omega \\
\text{M3-M1} & 9.5 \, \Omega \\
\text{M4-M1} & 15.0 \, \Omega \\
\text{M5-M1} & 19.6 \, \Omega \\
\text{M6-M1} & 21.8 \, \Omega \\
\end{array}
\]

- Vias made from Tungsten in Aluminum processes. Vias are Copper in Copper processes.

- Resistance of two via stacks at each end of M1 wire equivalent to about 0.1 mm wire (~20 \(\Omega\))
- Resistance of two via stacks at each end of M6 wire about the same as 1 mm narrow M6 wire (~60 \(\Omega\))!!!
- Use multiple vias in parallel to reduce effective contact resistance
- Copper processes have lower via resistance
• Capacitance depends on geometry of surrounding wires and relative permittivity, $\varepsilon_r$, of insulating dielectric
  - silicon dioxide, SiO$_2$, $\varepsilon_r = 3.9$
  - silicon flouride, SiOF, $\varepsilon_r = 3.1$
  - SiLK$^\text{TM}$ polymer, $\varepsilon_r = 2.6$

• Can have different materials between wires and between layers, and also different materials on higher layers
IBM Experimental 130nm Process

Figure by MIT OCW.
Capacitance Scaling

parallel plate capacitance $\propto \frac{\text{width}}{\text{spacing}} \times \text{length}$

- Capacitance/unit length $\sim$ constant with feature size scaling (width and spacing scale together)
  - Isolated wire sees approx. 100 fF/mm
  - With close neighbors about 160 fF/mm

- Need to use capacitance extractor to get accurate values
RC Delay Estimates

Penfield-Rubenstein model estimates:

\[ \text{Delay} = \sum_i \left( \sum_{j=1}^{j=i} R_j \right) C_i \]
Wire Delay Models

Wire has distributed $R$ and $C$ per unit length
- wire delay increases quadratically with length
- edge rate also degrades quadratically with length

Simple lumped $\Pi$ model gives reasonable approximation
- $R_w$ is lumped resistance of wire
- $C_w$ is lumped capacitance (put half at each end)

\[
\text{Delay} = R_{\text{driver}} \times \frac{C_w}{2} + (R_{\text{driver}} + R_w) \times \left(\frac{C_w}{2} + C_{\text{load}}\right)
\]
Wire Delay Example

• In 0.18\,\mu m TSMC, 5x minimum inverter with effective resistance of 3 \, k\Omega, driving FO4 load (25fF)

• Delay = R_{\text{driver}} \times C_{\text{load}} = 75 \, \text{ps}

• Now add 1mm M1 wire, 0.25\,\mu m wide
  − R_w = 320\,\Omega \text{ wire} + 22\,\Omega \text{ vias} = 344\,\Omega
  − C_w = 160\text{fF}

\[
\text{Delay} = R_{\text{driver}} \times \frac{C_w}{2} + (R_{\text{driver}} + R_w) \times \left( \frac{C_w}{2} + C_{\text{load}} \right) \\
= 3k\Omega \times 80fF + (3k\Omega + 344\Omega) \times (80fF + 25fF) \\
= 591\,\text{ps}
\]
Wire Delay Scaling, Local Wires

- For wire crossing same amount of circuitry
  - Resistance stays roughly constant
    - length decreases by same amount as width, height stays large and/or change material to copper
  - Capacitance decreases by scaling factor
    - cap/unit length constant, length decreases

- Wire delay tracks improvement in gate delay

[From Mark Horowitz, DAC 2000]
Wire Delay Scaling, Global Wires

- For wire crossing whole chip
  - Resistance grows linearly
  - Capacitance stays fixed
- Wire delay increases relative to gate delay

[From Mark Horowitz, DAC 2000]
Fewer Gates per Clock Cycle

- Processors in Intel 386 generation, around 50 FO4 gate delays per clock cycle

- Pentium-4 around 16 FO4 in normal clock, around 8 FO4 delays in fast ALU section

- Fastest 64-bit adder around 7 FO4 delays

- As measured in distance per clock cycle, wires are getting much slower
Process Technology Fixes

- **Reduce R**
  - use copper instead of aluminum, 40% reduction
  - provide upper layers with thicker metal for long range wires
  - provide more layers to improve density, makes wires shorter

- **Reduce C**
  - use low-k dielectric, >2x reduction possible
  - increase inter-layer spacing (limited effect, problems with via formation)
  - provide more layers to improve density, makes wires shorter
Layout Fixes

- **Reduce R**
  - make wires wider, increase in $C$ is less than increase in $C$ because of fringing fields
  - use parallel vias at contacts
  - floorplanning to keep wires short
  - careful routing to avoid unnecessary layer changes (vias)

- **Reduce $C$**
  - space wires further apart than minimum
  - avoid parallel wiring
Circuit Fixes - Repeaters

- Use repeaters
- Converts quadratic dependence into linear dependence on length (but watch the constants)
- Can determine optimal repeater sizing for minimum delay
Repeater Issues

- Repeater must connect to transistor layers
- Blocks other routes with vias that connect down
- Requires space on active layers for buffer transistors and power connections
- Repeaters often grouped in preallocated repeater boxes spread around chip
  - Repeater location might not give ideal spacing
Repeater Staging In and Out

- For minimum delay, must stage up to repeater size at start of wire, and stage up to load at end of wire
### Architectural Fixes: Pentium-4

<table>
<thead>
<tr>
<th></th>
<th>Pipeline Stages</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>TC Next IP</td>
</tr>
<tr>
<td>2</td>
<td>TC Fetch</td>
</tr>
<tr>
<td>3</td>
<td>Drive</td>
</tr>
<tr>
<td>4</td>
<td>Alloc</td>
</tr>
<tr>
<td>5</td>
<td>Rename</td>
</tr>
<tr>
<td>6</td>
<td>Queue</td>
</tr>
<tr>
<td>7</td>
<td>Schedule 1</td>
</tr>
<tr>
<td>8</td>
<td>Schedule 2</td>
</tr>
<tr>
<td>9</td>
<td>Schedule 3</td>
</tr>
<tr>
<td>10</td>
<td>Dispatch 1</td>
</tr>
<tr>
<td>11</td>
<td>Dispatch 2</td>
</tr>
<tr>
<td>12</td>
<td>Register File 1</td>
</tr>
<tr>
<td>13</td>
<td>Register File 2</td>
</tr>
<tr>
<td>14</td>
<td>Execute</td>
</tr>
<tr>
<td>15</td>
<td>Flags</td>
</tr>
<tr>
<td>16</td>
<td>Branch Check</td>
</tr>
<tr>
<td>17</td>
<td>Drive</td>
</tr>
</tbody>
</table>

Pipeline stages dedicated to driving signals across chip.
Scalable Layout

- Design such that critical paths are gate dominated not interconnect dominated on first target process.
- This helps ensure feature size shrinks give frequency improvements.
- If critical path contains too much interconnect RC, then shrink won’t see much frequency gain.
Coupling Capacitances

- Most capacitance is to neighboring wires
- If A switches, injects voltage noise on B
  - magnitude depends on capacitive divider formed:
    \[ \frac{C_{AB}}{C_{AB} + C_B} \]
- If A switches in opposite direction while B switches, coupling capacitance effectively doubles – Miller effect
- If A switches in same direction while B switches, coupling capacitance disappears
- These effects can lead to large variance in possible delay of B driver, possibly factor of 5 or 6 between best and worst case
Avoid placing simultaneously switching signals next to each other for long parallel runs.

Reroute signals which will be quiet during switching in between simultaneous switching signals.

Route signals close to power rails to provide capacitance ballast.

Tough problem to solve - moving one wire can introduce new problems.
- “Timing closure” causes many real-world schedule slips.
Electromigration

- The electrons from a DC current flow will tend to push metal atoms out of place (AC current flow OK)
- Main problem is power lines, but some signal lines have unidirectional current
- Manufacturers place a current density limit for metal to guarantee small resistance increase after ~10 years operation
- TSMC 0.18µm
  - 1mA/µm (metal wires 0.4µm thick)
  - 0.28 mA/via