Problem 1

The purpose of this recitation is to familiarize students with a variety of integer programming modeling techniques as described in the IP Formulation Guide and in the powerpoint tutorial on IP formulations.

We start with an integer program IP1 defined as follows:

\[
\begin{align*}
\text{max} & \quad 21x_1 + 32x_2 + 40x_3 + 49x_4 + 57x_5 + \\
& \quad +71x_6 + 82x_7 + 91x_8 + 100x_9 + 109x_{10} \\
\text{s.t.:} & \quad 2x_1 + 3x_2 + 4x_3 + 5x_4 + 6x_5 + \\
& \quad +7x_6 + 8x_7 + 9x_8 + 10x_9 + 11x_{10} \leq 900 \\
\forall i & = 1, \ldots, 3 \quad x_i \in \{0, 1\} \\
\forall i & = 4, \ldots, 10 \quad 0 \leq x_i \leq 100.
\end{align*}
\]

(IP1)

For each of the parts below, you are to add constraint(s) and possibly variables to ensure that the logical condition is satisfied by the integer program. Each part is independent; that is, no part depends on the parts preceding it. You do not need to repeat the integer programming objective or constraints given above. You may use the big M method for formulating constraint when it is appropriate.

(a) (4 points) Write a single linear constraint that is equivalent to the statement “If \( x_1 = 1 \), then \( x_2 = 0 \).”

(b) (4 points) Write a single linear constraint that is equivalent to the statement “\( x_2 = 1 \) or \( x_3 = 0 \),” but not both.

(c) (4 points) Add a binary variable \( w_1 \), and add two constraints that ensure that \( w_1 = 1 \) if \( x_5 + x_6 \geq 70 \), and \( w_1 = 0 \) if \( x_5 + x_6 \leq 69 \).

(d) (4 points) Add 3 binary variables \( w_2, w_3, \) and \( w_4 \) and at most 4 constraints so as to ensure that at least one of the following constraints is satisfied: (i) \( x_5 \leq 92 \), (ii) \( x_6 \leq 40 \), (iii) \( x_7 + x_8 \geq 74 \).

(e) (4 points) Add a single binary variable \( w_5 \) and two constraints to ensure that at least one of the following two constraints are satisfied (i) \( x_9 \leq 45 \), (ii) \( x_{10} \geq 22 \).

(f) (4 points) Add a single integer variable \( w_6 \) and a constraint that ensures that \( x_8 \) is divisible by 2 but not divisible by 4. (The remainder when dividing by 4 must be 2).

(g) (4 points) Add three binary variables \( w_7, w_8, \) and \( w_9 \) and two constraints that ensures that \( x_{10} = 13 \) or 39 or 88.
(h) (4 points) Add variable(s) and constraint(s) that model the cost of \( x_4 \) as \( f_4(x_4) \), which is defined as follows: If \( x_4 = 0 \), then \( f_4(x_4) = 0 \). If \( x_4 \geq 1 \), then \( f_4(x_4) = 250 + 49x_4 \).

(i) (8 points) Add variable(s) and constraint(s) that model the cost of \( x_5 \) as \( f_5(x_5) \), which is defined as follows: If \( 0 \leq x_5 \leq 10 \), then \( f_5(x_5) = 57x_5 \). If \( 11 \leq x_5 \leq 20 \), then \( f_5(x_5) = 570 \). If \( 21 \leq x_5 \leq 100 \), then \( f_5(x_5) = -480 + 50x_5 \).

Problem 2

As the leader of an oil-exploration drilling venture, you need to determine which 5 sites out of 10 to evaluate for drilling opportunities. The goal is to select 5 sites with the lowest overall cost. Label the sites \( S_1, S_2, \ldots, S_{10} \), and the exploration costs associated with each as \( c_1, c_2, \ldots, c_{10} \). Regional development restrictions are such that:

(i) Evaluating sites \( S_2 \) and \( S_7 \) will prevent you from evaluating either site \( S_6 \) or \( S_9 \).

(ii) Evaluating sites \( S_1 \) and \( S_3 \) will prevent you from also evaluating both sites \( S_5 \) and \( S_6 \).

(iii) Evaluating site \( S_3 \) or \( S_4 \) prevents you from evaluating site \( S_6 \).

(iv) Of the group \( S_3, S_6, S_7, S_8 \), at most two sites may be assessed.

Formulate an integer program to determine the minimum-cost exploration scheme that satisfies these restrictions. Try to develop a model in which the only variables are \( x_1, \ldots, x_{10} \), where \( x_j \) is 1 or 0 according as site \( j \) is evaluated or not. (For example, the constraint “Evaluating sites \( S_2 \) and \( S_7 \) will prevent you from exploring site \( S_6 \)” can be expressed as \( x_2 + x_6 + x_7 \leq 2 \) because the only binary solutions prohibited have \( x_2 = x_6 = x_7 = 1 \).)

Problem 3

Suppose you want to minimize or maximize a piecewise linear function of one variable, subject to linear constraints. This is a problem that can be solved by resorting to linear constraints only, possibly by adding extra variables. In this example, we consider the function with three pieces shown in Figure 1.

(a) Suppose we want to minimize \( f(x) \) shown in Figure 1. Assume that \( x \) is subject to a set of linear constraints that involve other variables \( A(x'|x) = b \), so that we cannot simply solve the problem by inspection because we do not know what values \( x \) will take. How can we formulate this problem in linear form? Do we need integer variables? (In your formulation, you can ignore the additional constraints \( A(x'|x) = b \).)

(b) Consider now the problem of maximizing \( f(x) \) of Figure 1, subject to a set of linear constraints that involve other variables. We cannot use the same approach of Part 2.A. Explain why and find an alternative way of formulating the problem, adding (binary or integer) variables as needed.
Figure 1: Piecewise linear function discussed in Problem 5.